

Refining models expressed in UML sequence diagrams

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- STAIRS -Steps to Analyze Sequence Diagrams with Refinement Semantics

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Motivation

Make use of classical refinement theory in a practical UML setting

- From theory to practice, and not the other way around
- We aim to explain how classical theory of refinement can be used to refine specifications expressed with the help
- Sequence diagrams can be used to explain other kinds of UML diagrams
- By defining refinement for sequence diagrams we implicitly define refinement for the UML as a whole



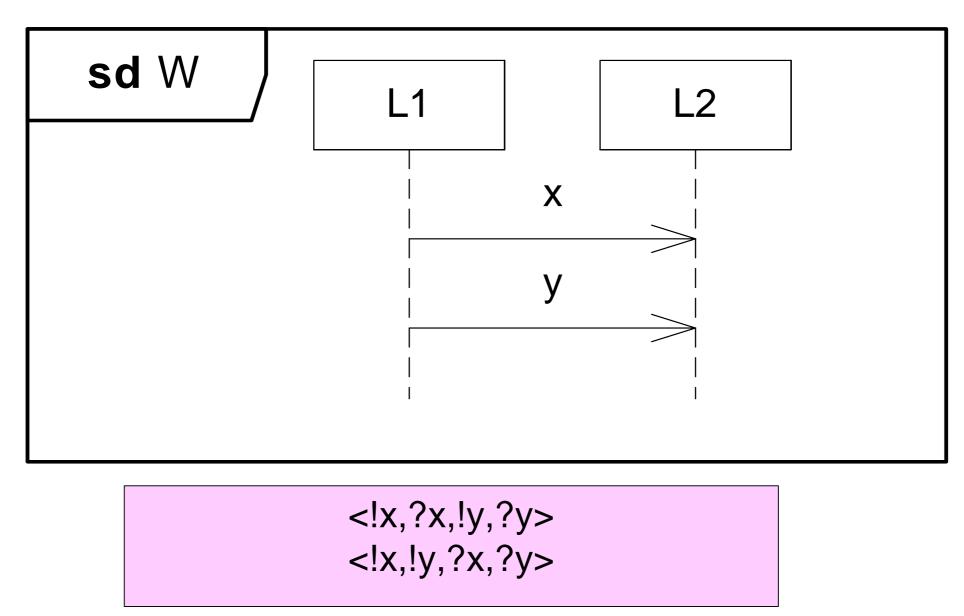
- Traces are used to represent system runs mathematically
- In the literature there are many different kinds of traces
- INF 5150 traces are sequences of events

<e1, e2, e3, e4, e4, e1, e2, e5,>

- Events are instantaneous
- The number of events in a trace may be finite
 - may be caused by: termination, deadlock, infinite waiting, system crash
- The number of events in a trace may be infinite
 - May be cause by: non-termination, livelock, non-termination by purpose



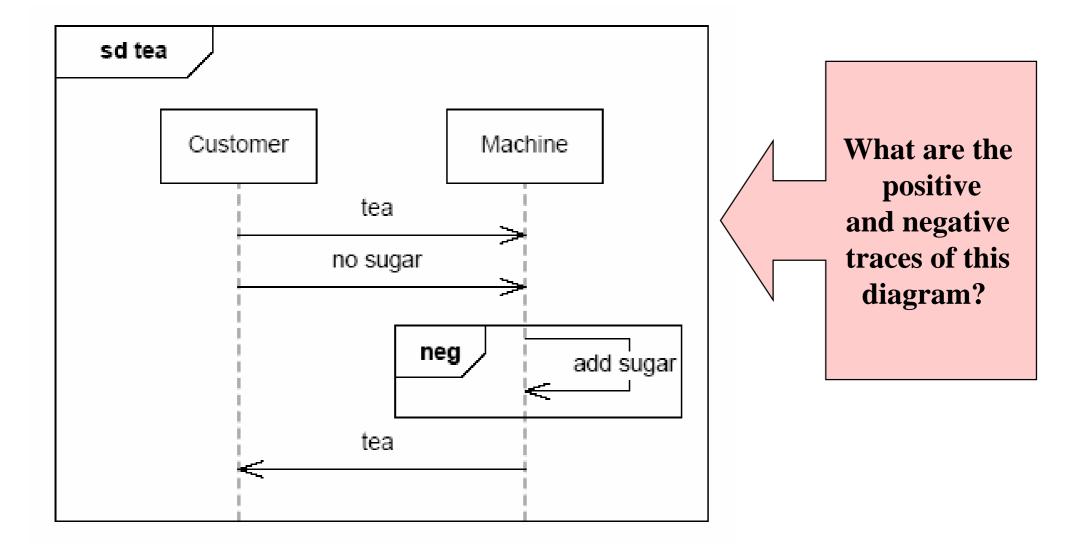
Weak sequencing



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Semantics of sequence diagrams (without xalt)

The formal semantics of a sequence diagram (without xalt) d is a pair (P,N) of sets of traces such that

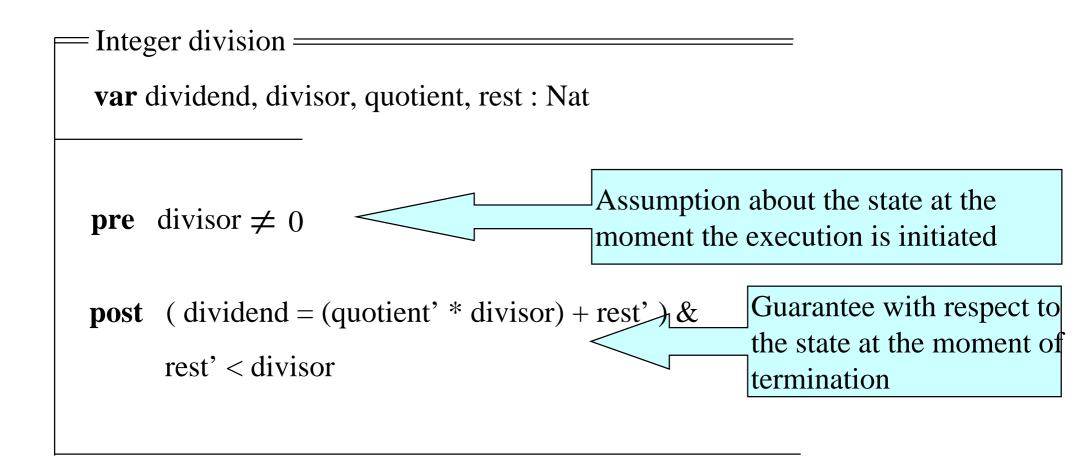
- P contains exactly the positive traces of d
- N contains exactly the negative traces of d
- For any diagram d, we use [[d]] to denote its semantics
- The same trace may be both positive and negative with respect to the same diagram
 - But if this is the case, you have probably not specified what you intended to specify

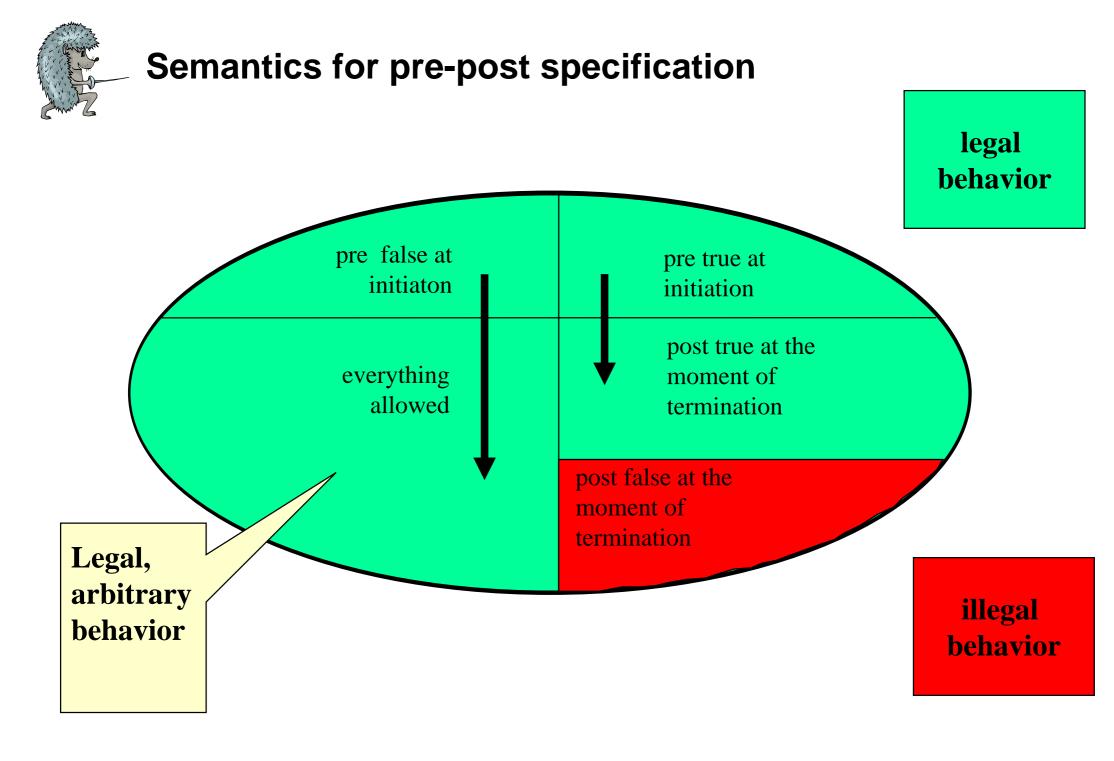
 A trace that is neither positive nor negative for a sequence diagram d is inconclusive with respect to d



Pre-post specifications

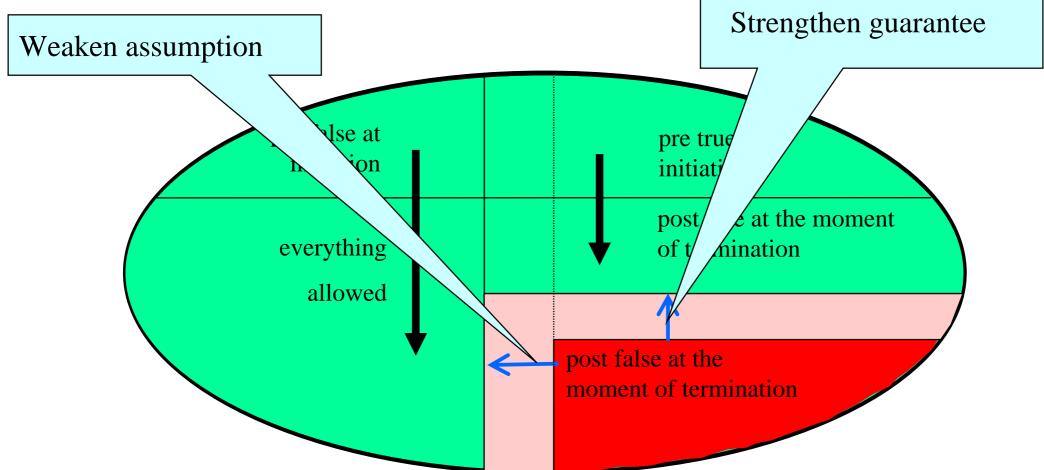
Pre-post specifications are based on the assumption-guarantee paradigm





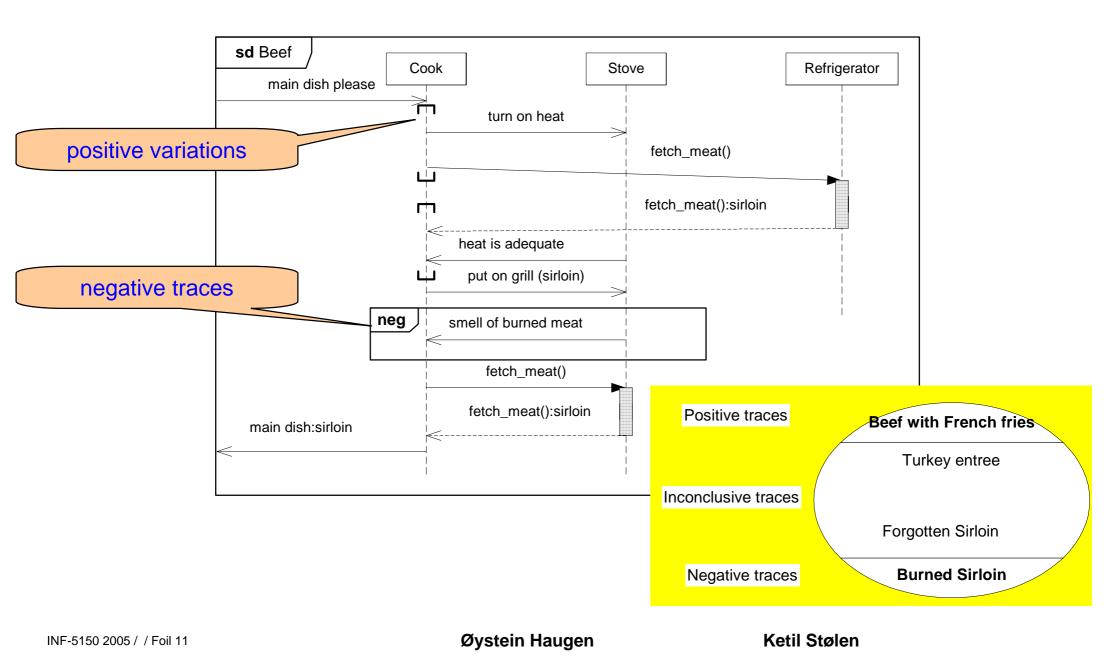


Property refinement for pre-post specifications





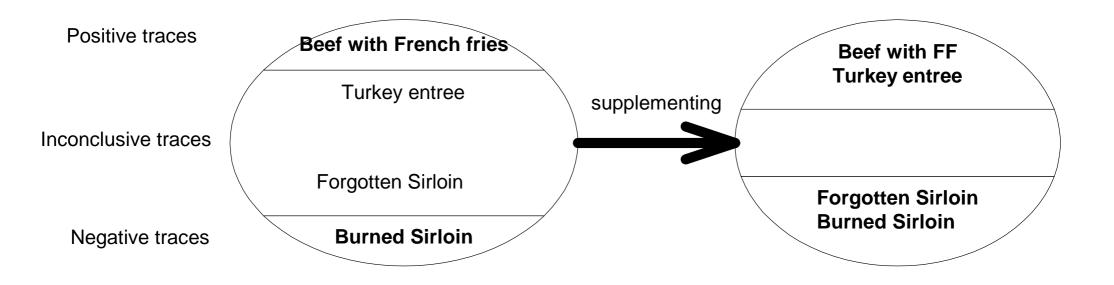
Property refinement of sequence diagrams





Supplementing

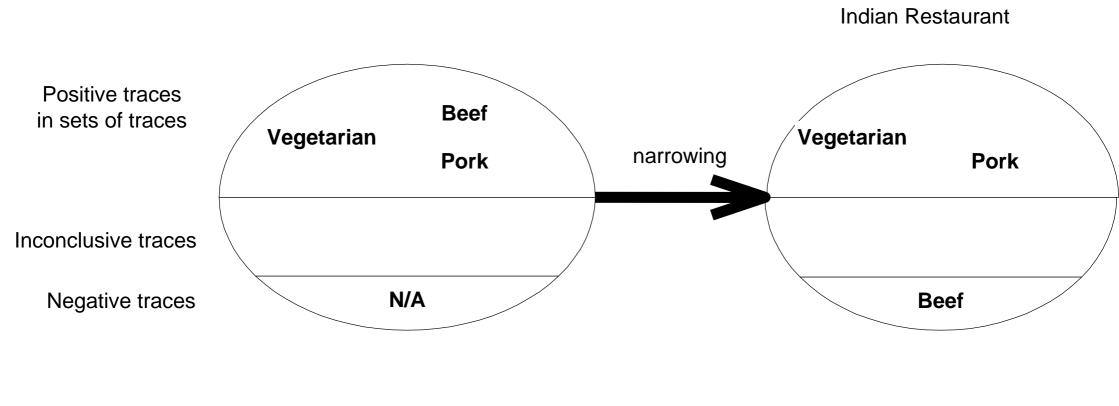
- To supplement means reducing the set of inconclusive traces by reclassifying inconclusive traces as either positive or negative
- The already positive traces remain positive
- The already negative traces remain negative





Narrowing

- To narrow means reducing the set of positive traces by reclassifying positive traces as negative
- The inconclusive traces remain inconclusive
- The already negative traces remain negative





INDIRECT DEFINITION: Property refinement in STAIRS

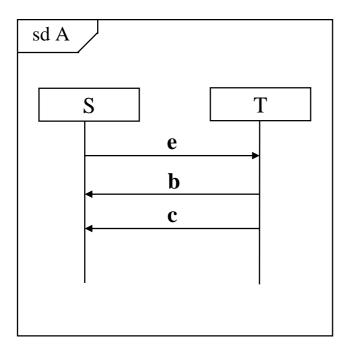
A sequence diagram B is a property refinement of a sequence diagram A if

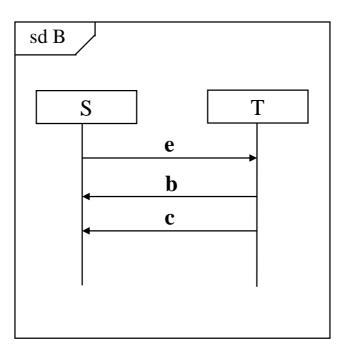
- A and B are semantically identical
- B can be obtained from A by supplementing
- B can be obtained from A by narrowing
- B can be obtained from A by a finite number of steps

A -> C1 -> C2 -> ->Cn->B

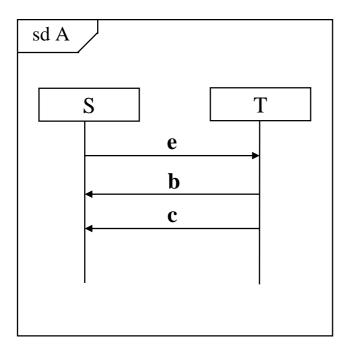
each of which is either a supplementing or a narrowing

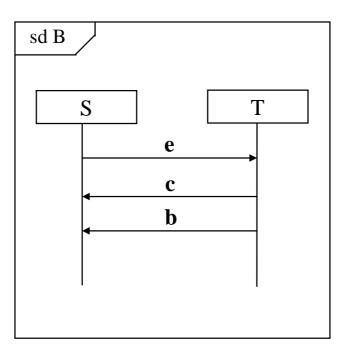




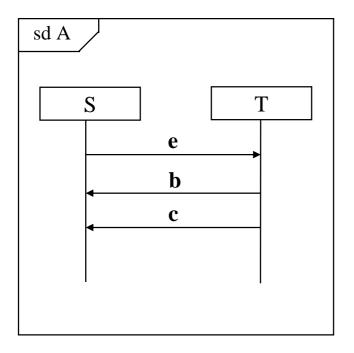


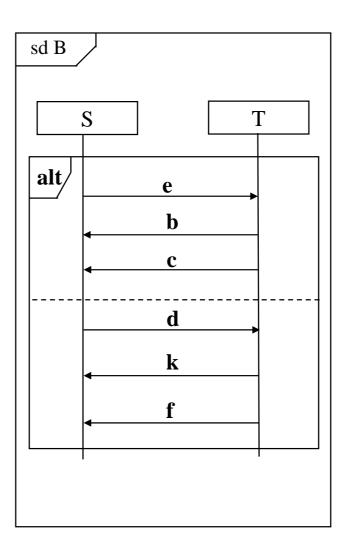




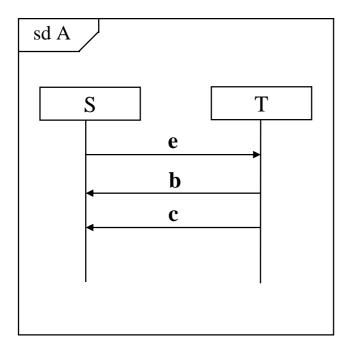


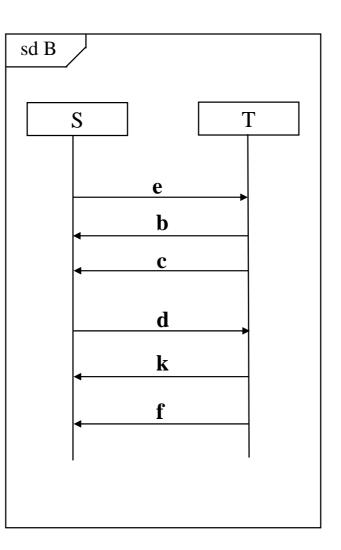




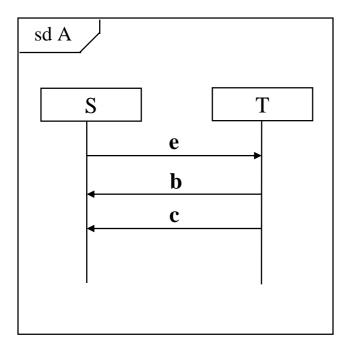


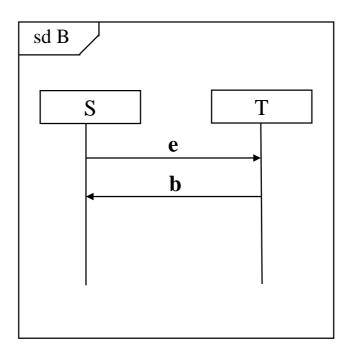










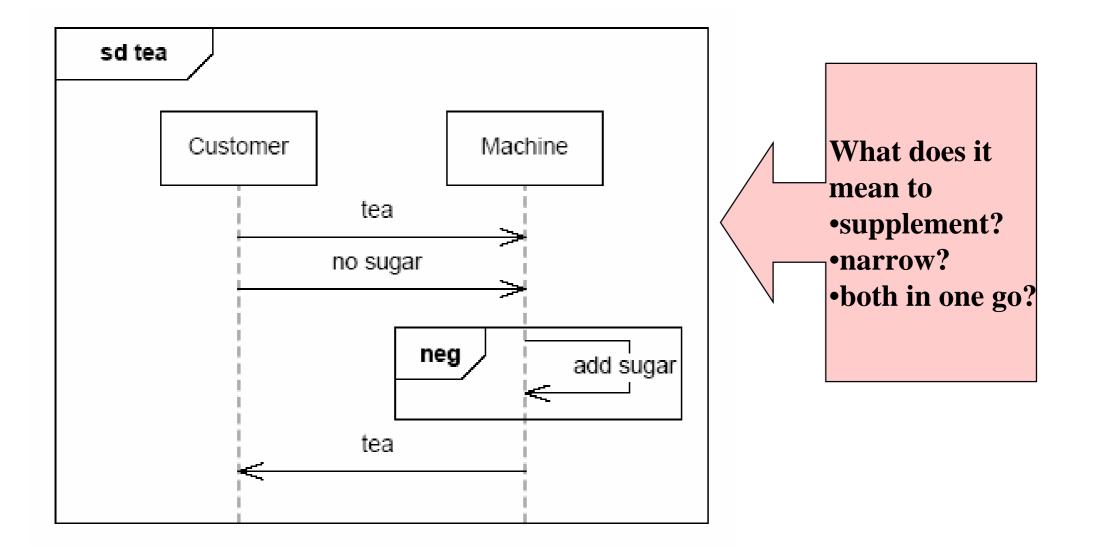




DIRECT DEFINITION: Property refinement in STAIRS

- A sequence diagram B is a property refinement of a sequence diagram A if
 - every trace classified as negative by A is also classified as negative by B
 - every trace classified as positive by A is classified as either positive or negative by B







Distinguishing two different types of nondeterminism

- Classical distinction between mandatory and potential nondeterminism
- Most specification languages cannot distinguish mandatory and potential non-determinism



The need for both alt and xalt

 Potential non-determinism captured by alt allows abstraction and inessential non-determinism

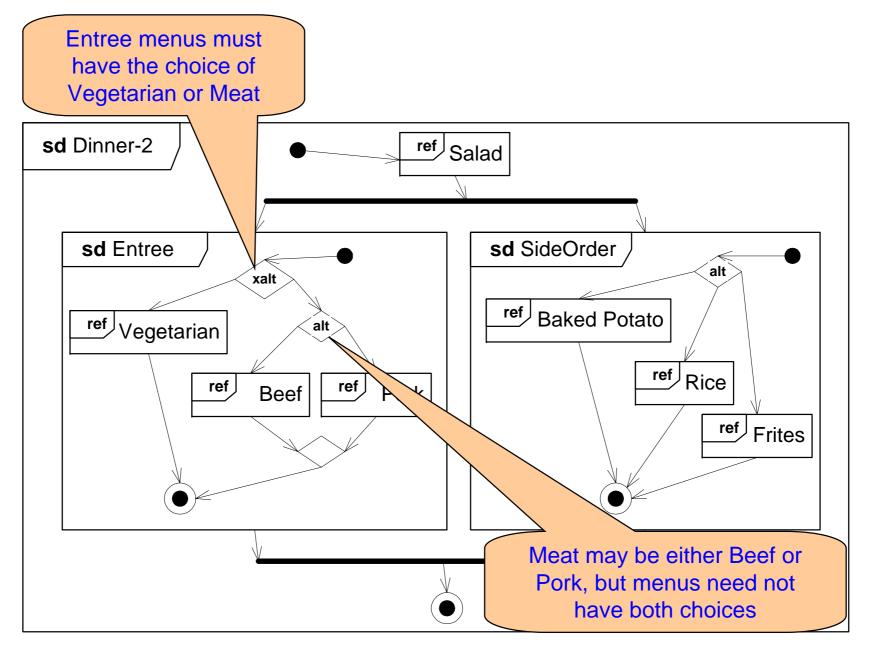
- Under-specification
- Non-critical design decisions may be postponed

Mandatory non-determinism captured by xalt characterizes nondeterminism that must be reflected in every correct implementation

- Makes it possible to specify games
- Important in relation to security
- Also helpful as a means of abstraction



Restaurant example with both alt and xalt



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General STAIRS-semantics for sequence diagrams

 The semantics of a sequence diagram is a set of pairs of sets of traces

- {(P1,N1), (P2,N2),..., (Pn,Nn)}
- where for each index j
 - Nj is a set of negative traces
 - Pj is a set of positive traces
- We refer to these pairs as ``interaction obligations''
- Each interaction obligation must be fulfilled by a correct implementation



The semantics of alt and xalt

$$\llbracket d_1 \text{ alt } d_2 \rrbracket \stackrel{\text{def}}{=} \{ (p_1 \cup p_2, n_1 \cup n_2) \mid (p_1, n_1) \in \llbracket d_1 \rrbracket \land (p_2, n_2) \in \llbracket d_2 \rrbracket \}$$

$\llbracket d_1 \operatorname{xalt} d_2 \rrbracket \stackrel{\text{def}}{=} \llbracket d_1 \rrbracket \cup \llbracket d_2 \rrbracket$



Property refinement of an interaction obligation

$(p_1, n_1) \rightsquigarrow (p_2, n_2)$ iff $n_1 \subseteq n_2 \land p_1 \subseteq p_2 \cup n_2$

An interaction obligation (p2,n2) is a property refinement of an Interaction obligation (p1,n1) if and only if

n1 is a subset of or equal to n2, and p1 is a subset of or equal to the union of p2 and n2



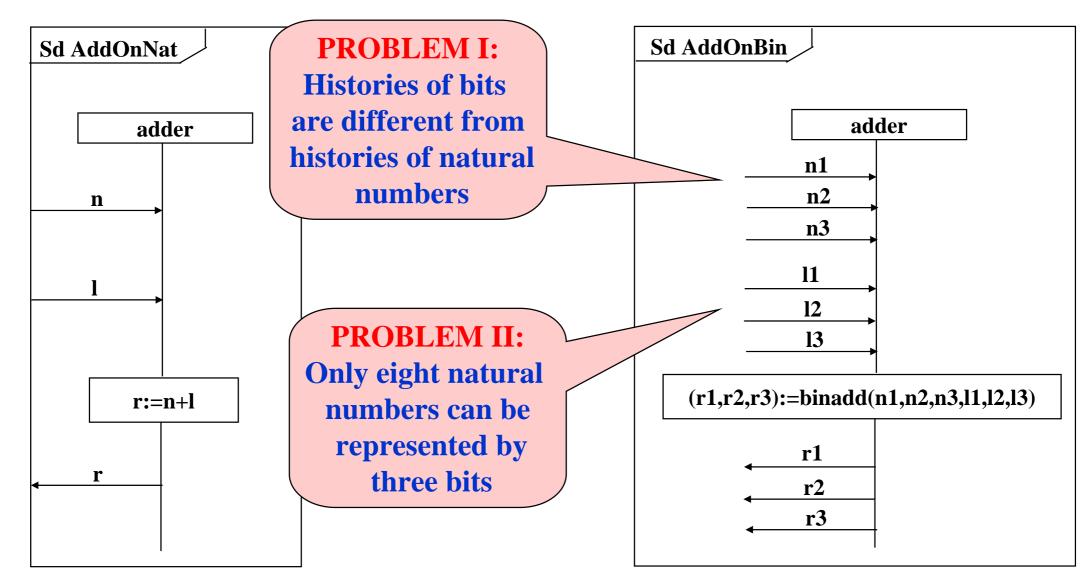
A sequence diagram d' is a refinement of a sequence diagram d, written $d \rightsquigarrow d'$, iff

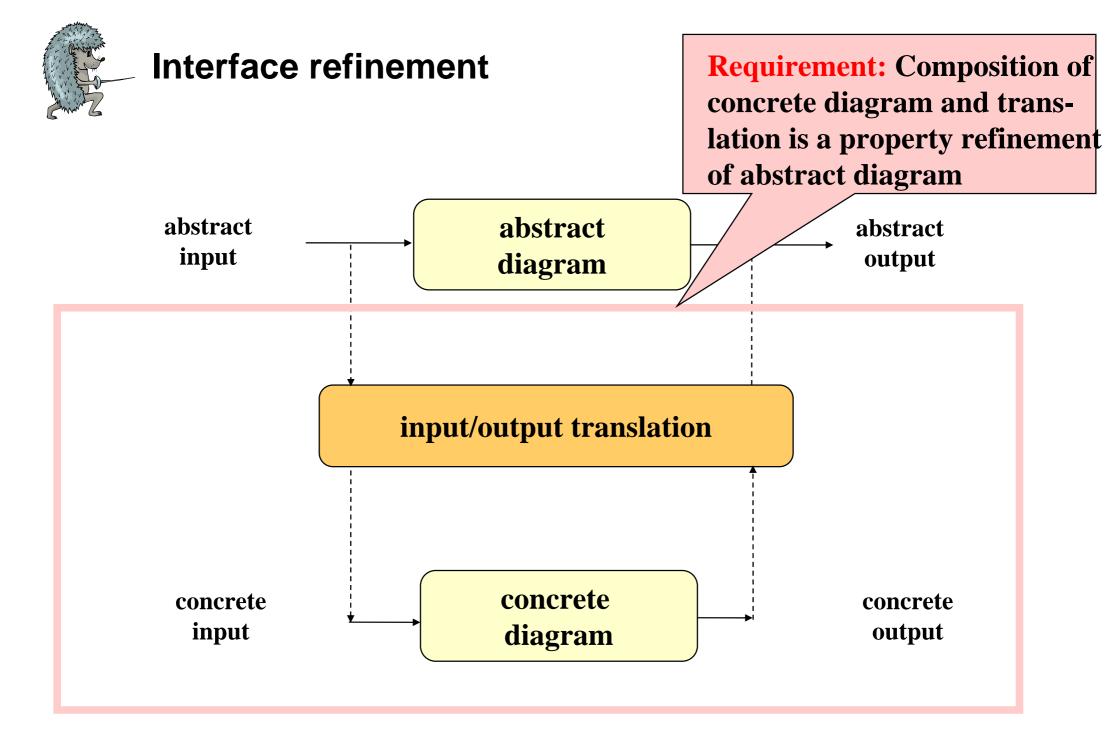
$$\forall o \in \llbracket d \rrbracket : \exists o' \in \llbracket d' \rrbracket : o \rightsquigarrow o'$$

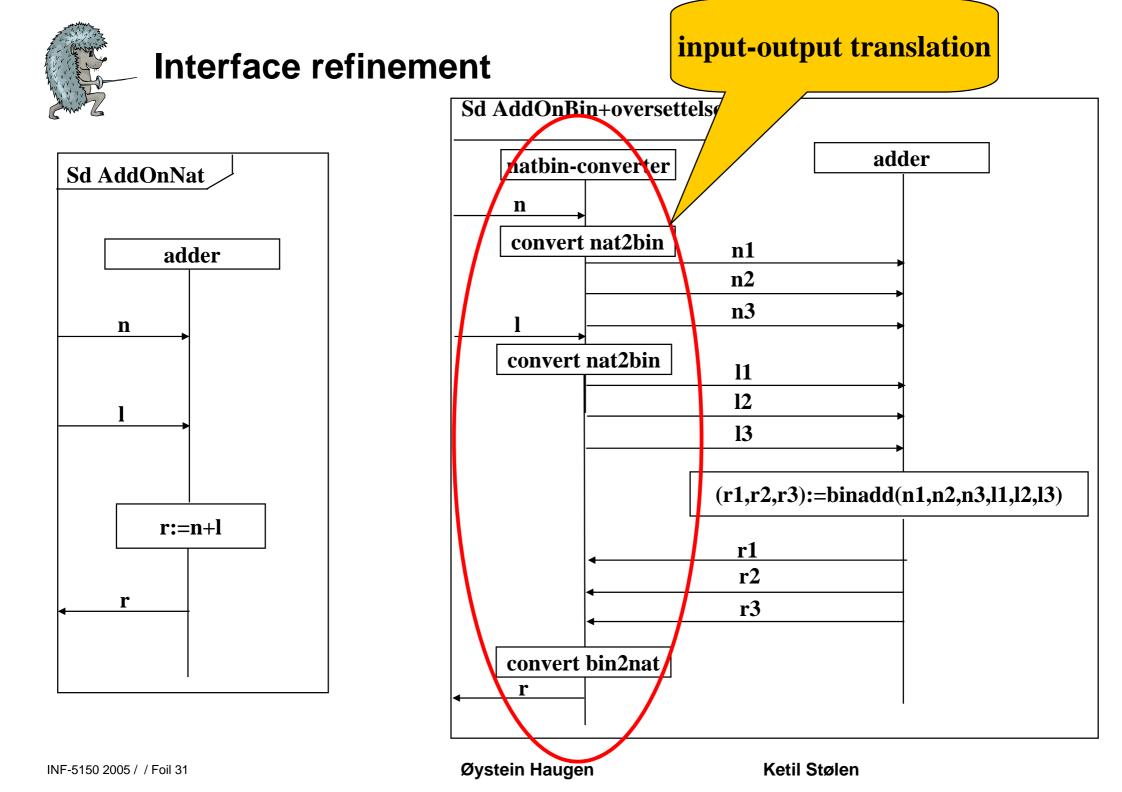
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Property refinement is not enough

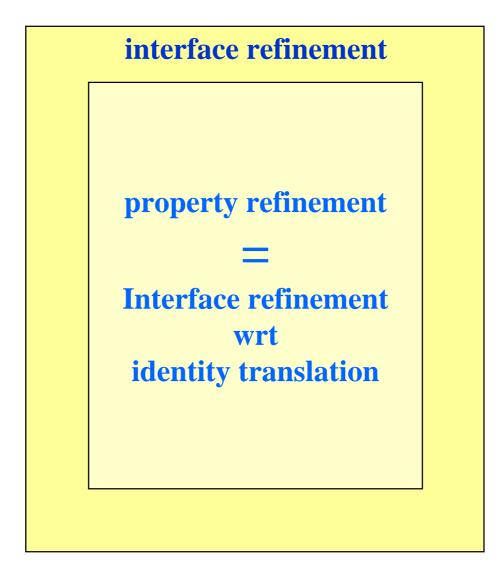








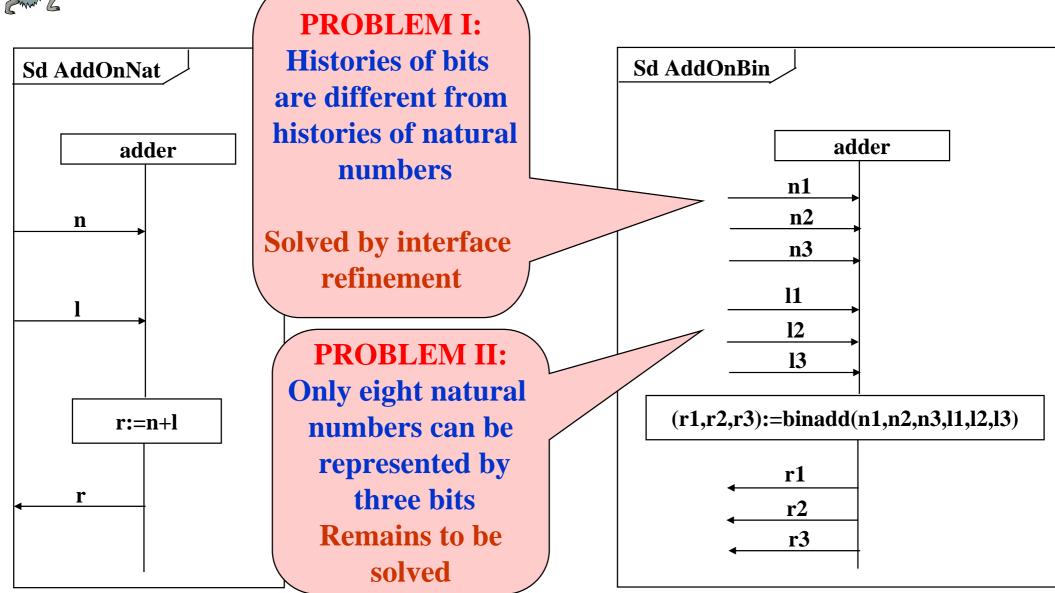
Interface refinement generalizes property refinement



In STAIRS interface refinement is a special case of detailing

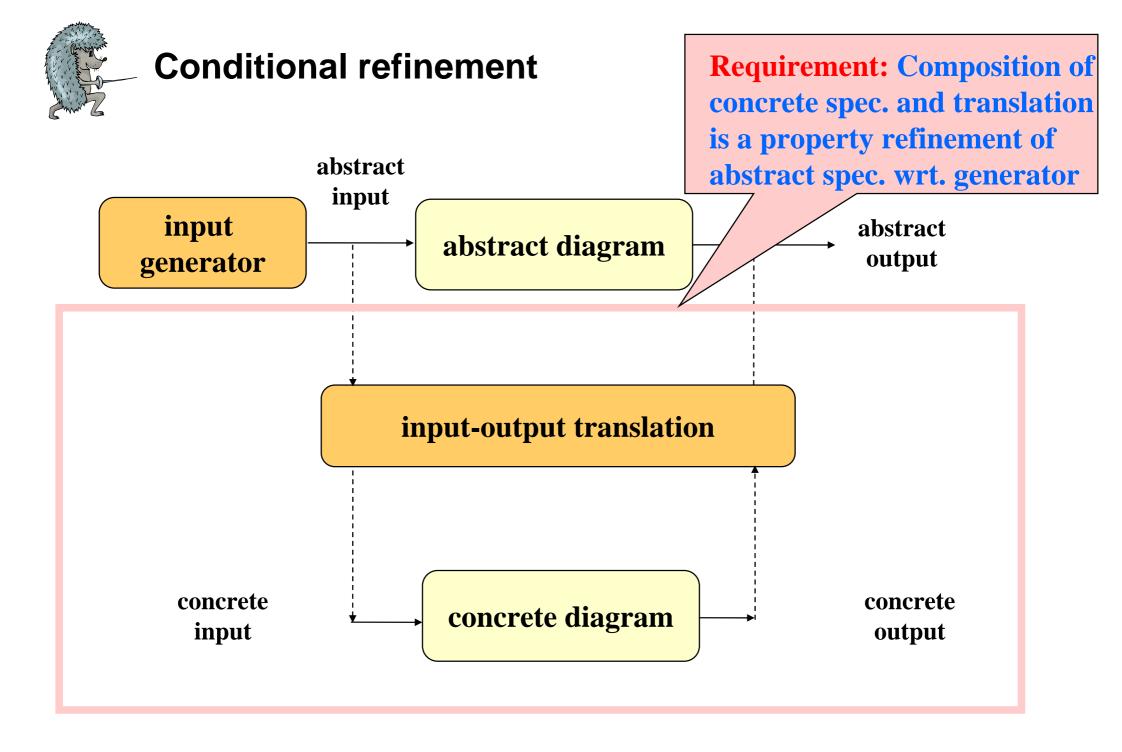


Interface refinement is not enough



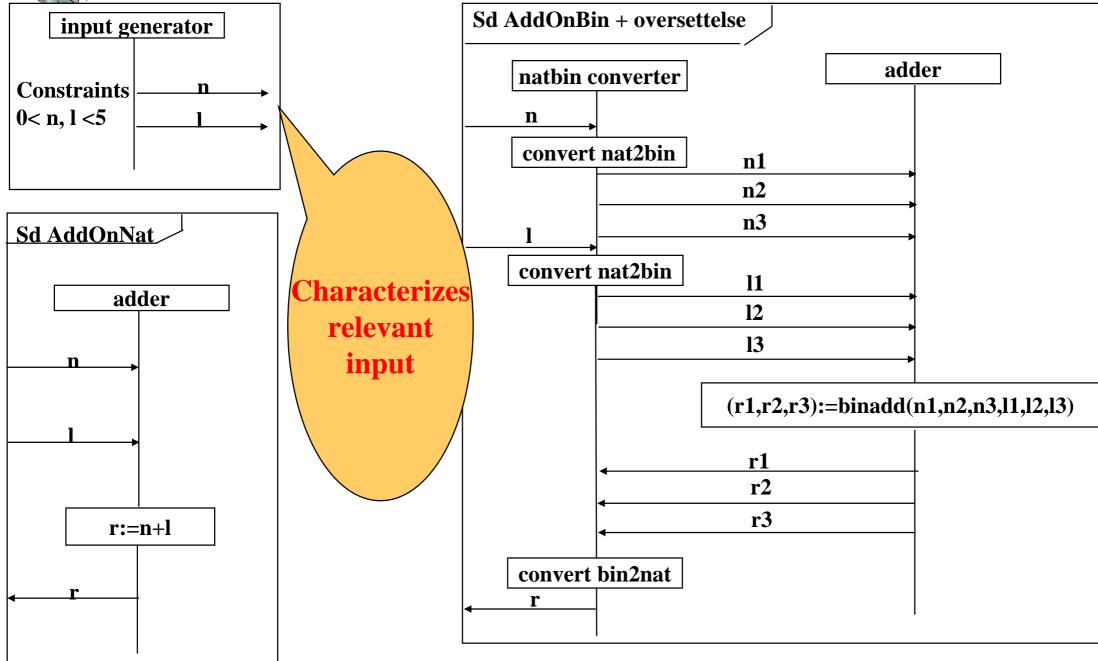
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Conditional interface refinement (another form of detailing)



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Literature on refinement

The classics:

- C.A.R. Hoare. An axiomatic basis for computer programming.
 Communications of the ACM, 12:576-583, 1969
- C.A.R. Hoare. Proof of correctness of data representations. Acta Informatica, 1:271-282, 1972
- O-.J. Dahl, C.A.R. Hoare. Hierarchial program structures. In structured programming, pages 175-220. Academic Press, 1972
- E.W.Dijkstra. A discipline of programming. Prentice-Hall, 1976

STAIRS is to a large extent based on ideas taken from:

 Manfred Broy, Ketil Stølen. Specification and Development of Interactive Systems - FOCUS on Streams, Interfaces and Refinement. ISBN 0-387-95073-7, Springer, 2001



Security analysis

First of three lectures October 28

We will

- classify dependability concepts
- introduce motivate and explain a basic terminology for risk management and risk analysis
- relate risk management to system development
- describe and motivate the different processes of risk management and how they interact