

# Robustness – the art of preparing for the unexpected

**Version 071123** 



#### The exceptional

- Data may have strange syntax or values
  - we apply common data-parsing techniques
- An unexpected signal arrives
  - we explicitly describe every conceivable transition
- No signal arrives
  - we guard our protocols with timers
- Security issues
  - authentication + logging + statistics
- Availability issues
  - self tests

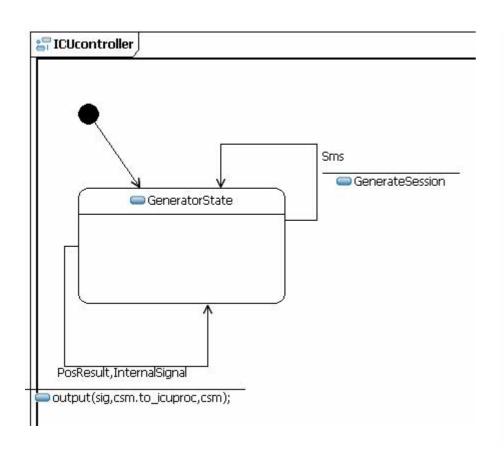


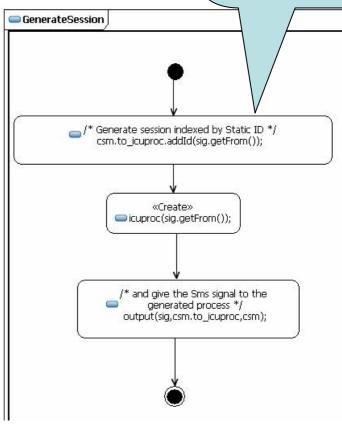
#### Handling an error or exceptional situation

- The invalid situation is due to an inadequate user input
  - then we know what caused it and the user should be notified
    - ICU: The user is notified by an SMS
- The invalid situation is due to an internal error
  - the reason is unclear, but the situation has become erroneous
  - The correct recovery may be hard to specify, but we believe that terminating the whole program is probably the last resort
    - ICU: different responses:
      - Try and send SMS to the user (if the appropriate user is known)
      - Dump the call stack on console (syserr) (very low level)
      - Terminate the session (and notify the session owner by an SMS)

#### ICUcontroller's GenerateSession

Data invariant is that Static ID should not already be used







#### **Checking the data invariant**

- Our task is to check whether STAT-ID is already the ID of another ICUprocess
- Here are the checking strategies:
  - checking directly the data of the routing port
    - simple, but on low (Java) level
  - sending a probe signal and wait for its possible consequences
    - more protocol needed, and possibly changing the forward() operation
    - if the normal response is that a timer must expire, this will be slow
  - recording which static ids are active (in the Archive)
    - lots of book-keeping, slower, overkill
- We go for the simple java-oriented solution this time

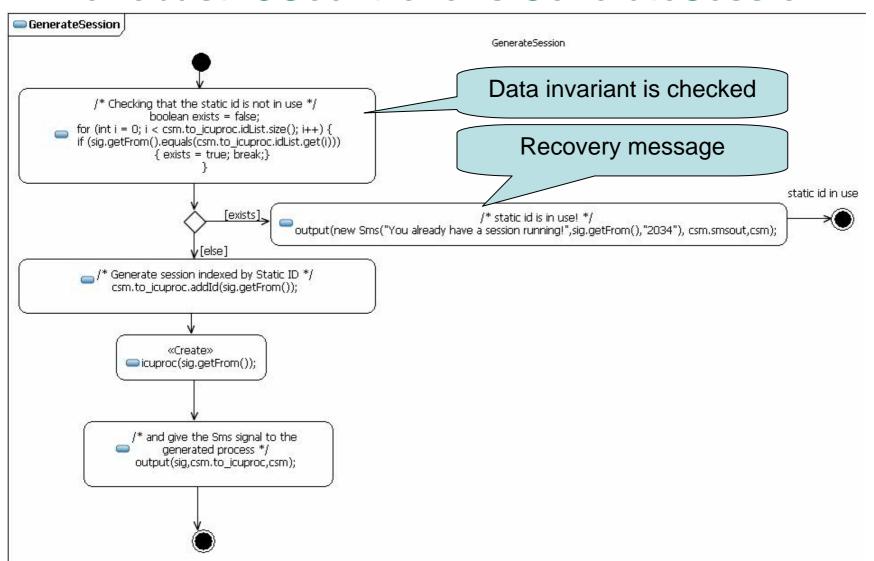


#### Error recovery for the static id re-use

- The cause of the static id re-use is most probably because the user has sent two requests in quick sequence
- We should response by returning an error message to the user
  - This will imply fixing the composite structure
- Move to a final state
  - in our service-oriented architecture, the service session is the natural unit of recovery, i.e. canceling the current service session is often the best approach

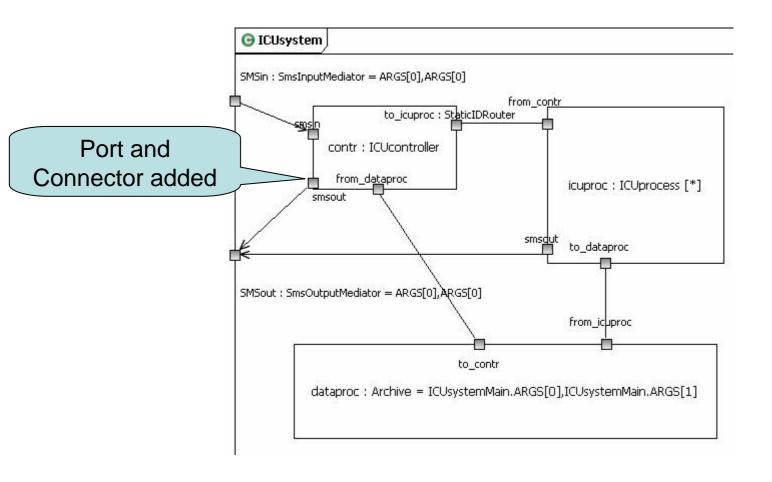


#### The robust ICUcontroller's GenerateSession





#### **Modified Composite Structure**





#### ICUcontroller: the exceptional

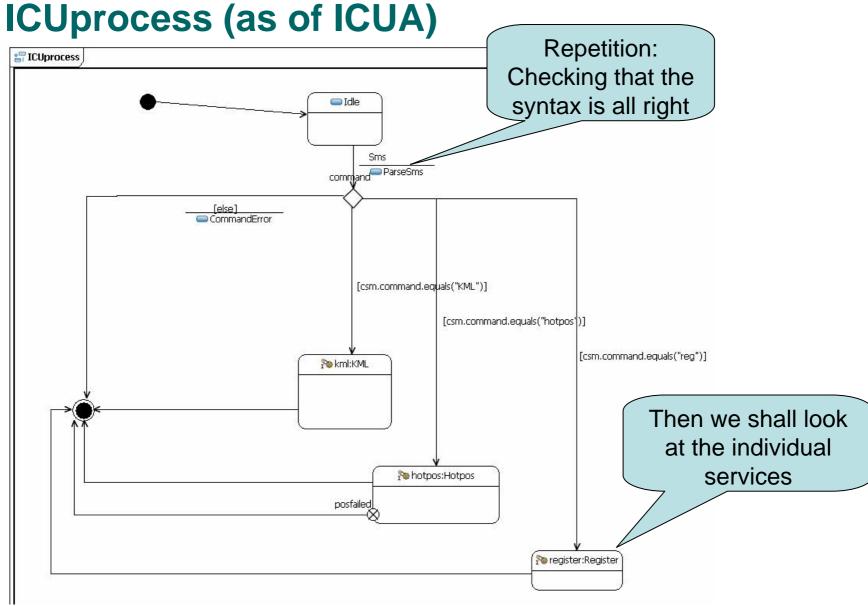
- Data may have strange syntax or values
  - Now checking for static id already in use
- An unexpected signal arrives
  - ICUcontroller handles Sms, PosResult and InternalSignal in all states
  - We are going to look at unexpected signals for ICUprocess
- No signal arrives
  - ICUcontroller does not have such waiting situations (?)
  - we shall guard our protocols/services with timers (ICUprocess)
- Security issues
  - authentication + logging + statistics
    - Authentication is not needed to enter ICUcontroller
    - we are going to check for registration in ICUprocess
- Availability issues
  - self tests
    - We could use ICUcontroller to test availability of PATS (but don't)
  - we will consider this with the Archive



#### **Explicit transitions please!**

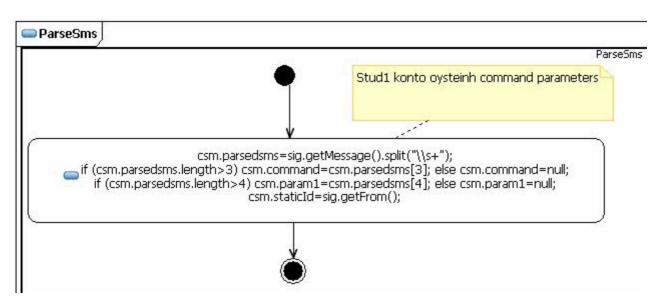
- Finite State Machines have a great advantage by being finite!
  - there is a finite set of transitions to execute
  - we can make sure to cover them all
- UML State Machines also define default transitions
  - where the signal is just discarded/consumed
  - We believe that default transitions are a warning of design flaw
- Not all signals can be properly handled at any time
  - We may defer a signal to a state where the signal can be dealt with







#### **ParseSms**



- If the Sms does not start with "Stud1 konto username" it will not come to the program at all
  - Still we may choose to check for it due to running on FakePats
- If there are more than 1 parameter, there is also an error
  - at least for the set of services of ICU that we have up to now
- We should give user syntax error messages right away
  - and not hide it by letting command be null

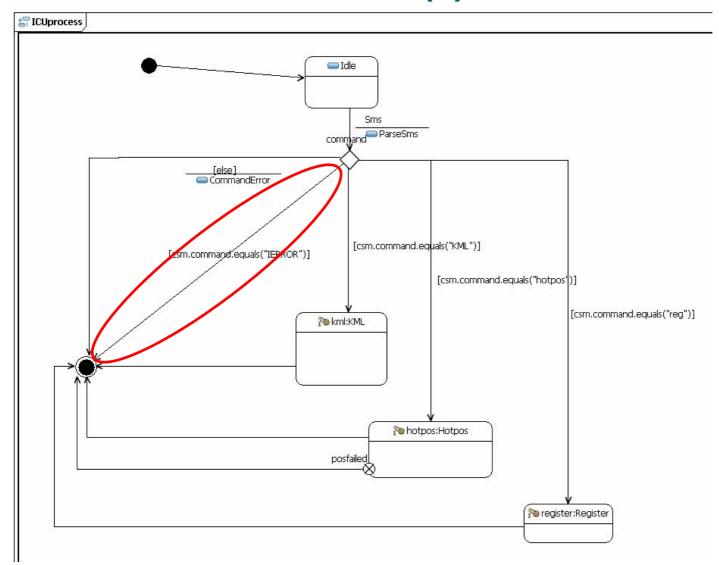


#### ParseSms robustified (1)

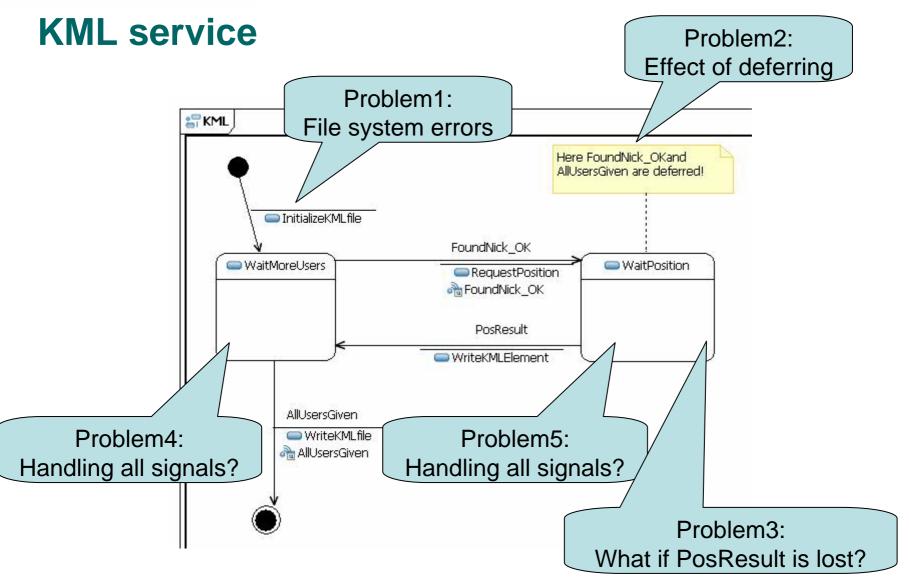
```
csm.parsedsms=sig.getMessage().split("\\s+");
/* check for existence of necessary prefix - very rudimentary */
if (csm.parsedsms.length<=3)
{ output(new Sms("ICU: Syntax error - no command", sig.getFrom(), "2034"), csm.smsout, csm);
 csm.command = "IERROR";
else
{ csm.command=csm.parsedsms[3];
 /* check for only one parameter */
 if (csm.parsedsms.length>5)
 { output(new Sms("ICU: Too many parameters!",sig.getFrom(),"2034"), csm.smsout,csm);
  csm.command = "IERROR";
 else
 { if (csm.parsedsms.length>4) csm.param1=csm.parsedsms[4]; else csm.param1=null;
csm.staticId=sig.getFrom();
```



# ParseSms robustified (2)





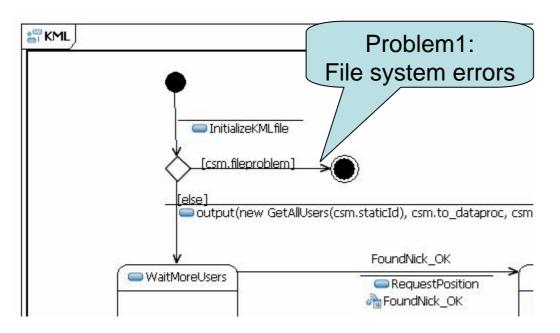




#### KML problems (1)

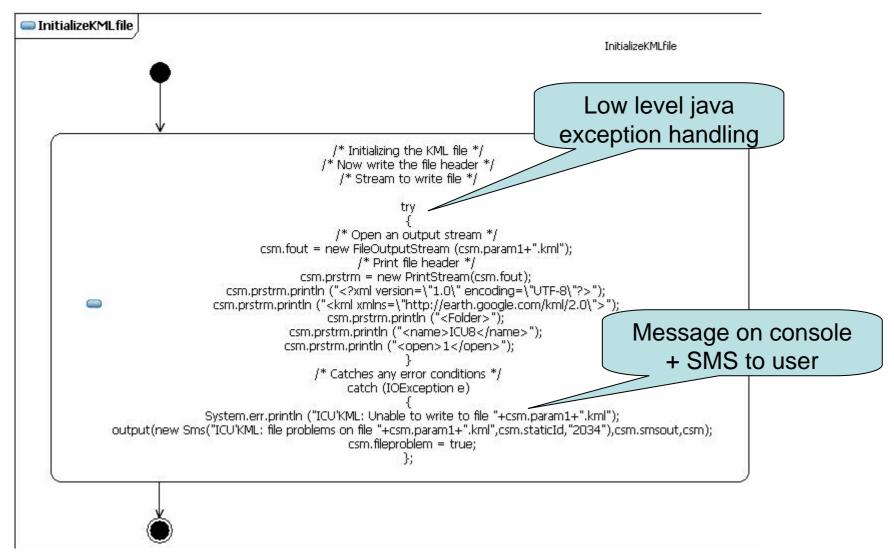
- 1: File writing problems
  - currently error dumped on console, and proceed as if no problem has arisen
    - not adequate: if the initialization fails to write on the file, the session should terminate, and double messages given (to console and user)

#### 1: File writing problems - New KML machine





# 1: File writing problems (2)





# KML problems (2)

- 2: The deferring of FoundNick\_OK is motivated by wanting to handle one positioning at the time
  - but the effect is the need to handle many defers
    - since the database produces users faster than PATS positions them
    - actually the #defers are in the order of #users<sup>2</sup>
  - and decreased efficiency due to this defer-handling and since positioning requests may be done in parallel (possibly)
    - but in fact sending too many positioning requests very quickly seems to stress PATS such that sometimes requests are lost
  - The optimal solution may be to introduce a little more protocol to sequentialize such that the Archive is explicitly asked to give the next user
    - rather than giving all users in a stream of messages
    - ... but we keep to the *defer* solution to show in detail how *defer* is



# JFTrace of the deferring KML

Filtered Trace from /127.0.0.1:54321 at 2007-04-29 23:48:08.759				
New ICUsystem_Archive@4e224cb9				
New ICUsystem_ICUcontroller@40				
ICUsystem_Archive@4e224cb9	null	StartMessage@4f430cb9		Idle
ICUsystem_ICUcontroller@406f0cb9	null	StartMessage@41278cb9		GeneratorState
ICUsystem_ICUcontroller@406f0cb9	GeneratorState	Sms@3e18cba (Stud1 konto oysteinh KML,2034,91390900)	New ICUsystem_ICUprocess@1dadccba Output Sms@3e18cba (Stud1 konto oysteinh KML,2034,91390900)	GeneratorState
ICUsystem_ICUprocess@1dadccba	null	StartMessage@1d884cba		Idle
ICUsystem_ICUprocess@1dadccba	Idle	Sms@3e18cba (Stud1 konto oysteinh KML,2034,91390900)	Output GetAllUsers@298accha (91390900)	WaitMoreUsers^kml
ICUsystem_Archive@4e224cb9	ldle	GetAllUsers@298accba (91390900)	Output FoundNick_OK@2fc1ccba (No91390900, 91390900, 91390900) Output FoundNick_OK@2e7c8cba (No66688899, 66688899, 91390900) Output FoundNick_OK@3bad8cba (No09090909, 09090909, 91390900) Output AllUsersGiven@3a474cba (91390900)	ldle
ICUsystem_ICUcontroller@406f0cb9	GeneratorState	FoundNick_OK@2fc1ccba (No91390900, 91390900, 91390900)	Output FoundNick_OK@2fc1ccba (No91390900, 91390900, 91390900)	GeneratorState
CUsystem_ICUcontroller@406f0cb9	GeneratorState	FoundNick_OK@2e7c8cba (No66688899, 66688899, 91390900)	Output FoundNick_OK@2e7c8cba (No66688899, 66688899, 91390900)	GeneratorState
		FoundNick_OK@3bad8cba (No09090909, 09090909, 91390900)	Output FoundNick_OK@3bad8cba (No09090909, 09090909, 91390900)	GeneratorState
		AllUsersGiven@3a474cba (91390900)	Output AllUsersGiven@3a474cba (91390900)	GeneratorState
		FoundNick_OK@2fc1ccba (No91390900, 91390900, 91390900)	Output PosRequest@be80cb9	/vaitr ssition^kml
		FoundNick_OK@2e7c8cba (No66688899, 66688899, 91390900)	F1 00 890 94000	Saved
		FoundNick_OK@3bad8cba (No09090909, 09090909, 91390900)		Saved
		AllUsersGiven@3a474cba (91390900)		Saved
		Poskesult@31584cb9	Output PosResult@31584cb9	GeneratorState
		PosResult@31584cb9	P 10 (00.0) Particle of the control	WaitMoreUsers^kml
		FoundNick_OK@2e7c8cba (No66688899, 66688899, 91390900)	Output PosRequest@53838cb9	Waitr ssition/kml
		FoundNick_OK@3bad8cba (No09090909, 09090909, 91390900)	7 (5) (5) (5)(1)	Saved
		AllUsersGiven@3a474cba (91390900)		Saved
		PosResult@54868cb9	Output PosResult@54868cb9	SeperatorState
		PosResult@54868cb9	10 CS17	WaitMoreUsers^kml
			Output PosRequest@6c4dccb9	vVaitPosition^kml
		AllUsersGiven@3a474cba (91390900)	VI OS SVIII	Saved
		PosResult@78c88cb9	Output PosResult@78c88cb9	GeneraterState
		PosResult@78c88cb9		vvaitivoreUsers^kml
ICUsystem_ICUprocess@1dadccba	WaitMoreUsers^kml	AllUsersGiven@3a474cba (91390900)	Output Sms@6d174cb9 (null.kml:E0104541,N595627,91390900,2034)	FinalState
	State Machine New ICUsystem_Archive@4e224cb9 New ICUsystem_ICUcontroller@40 ICUsystem_IcUcontroller@40f0cb9 ICUsystem_ICUcontroller@40f0cb9 ICUsystem_ICUcontroller@40f0cb9 ICUsystem_ICUprocess@1dadccba ICUsystem_ICUprocess@1dadccba ICUsystem_ICUcontroller@40f0cb9 ICUsystem_ICUcontroller@40f0cb9 ICUsystem_ICUcontroller@40f0cb9 ICUsystem_ICUcontroller@40f0cb9 ICUsystem_ICUcontroller@40f0cb9 ICUsystem_ICUprocess@1dadccba	State Machine	State Machine	State Machine

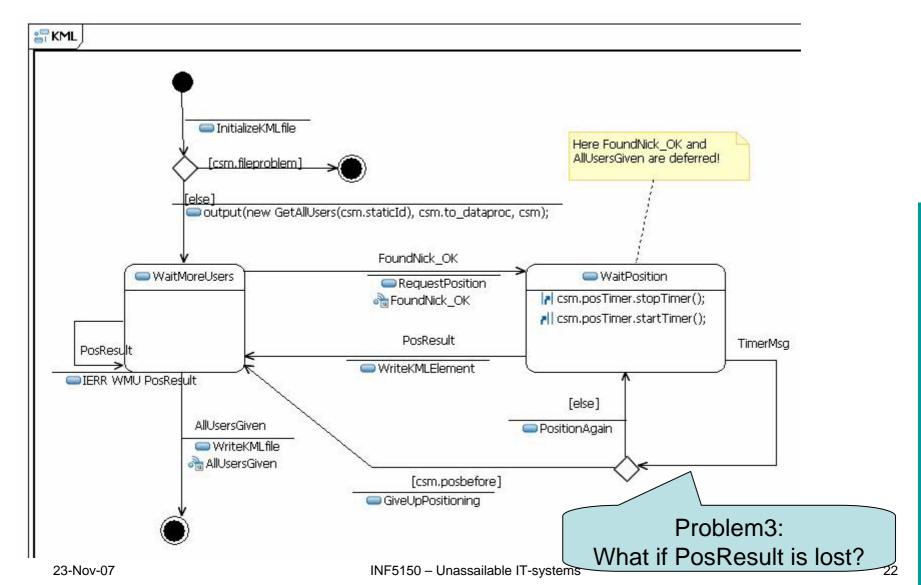


# KML problems (3) – needing timers

- 3: Even when using defer, we have no guarantee that the PosRequest results in a corresponding PosResult
  - We shall have to guard the PosResult by a timer
  - What then to do if the guarding timer expires?
    - Giving an SMS to the user for every non-positioned phone may be too many SMSes
      - and we could cut off after a small number of such messages (say 3)
        - and then give a more general error message and terminate KML session
    - We could try again to position the failed one (one retry)
  - What if the timer has expired, recovery has been done, and then the PosResult appears very late?
    - In our case this will have a cascading effect of PosResult appearing when it should not
      - this actually becomes rather tricky! (will be covered later)



# 3: Including the timer in KML

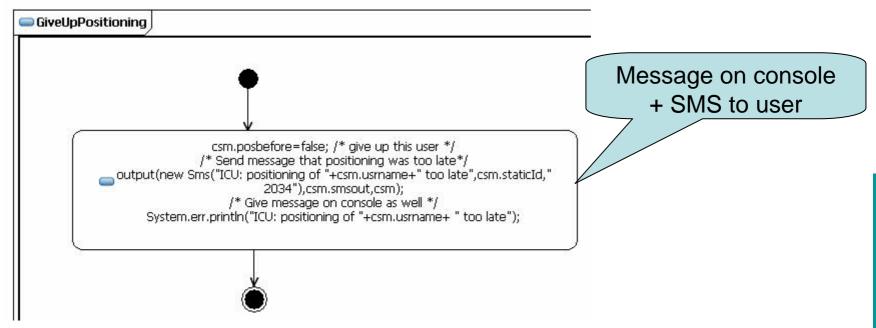




#### 3: Guarding PosResult with a timer

- Declare a timer
  - Define a public attribute of the ICUprocess state machine
    - never define attributes of substate machines
  - Add stereotype «TimerMsg» to that attribute (it has no type)
- Set the timer
  - To guard an expected response, start the timer in entry-clause
  - Set default time in property of timer attribute (here: 10000 ms)
- Stop the timer
  - when normal response happens (here: PosResult transition)
  - or easier: stop the timer in the exit-clause
- Timer expires
  - TimerMsg is the type of the trigger (always)
    - sig == the defined Timer attribute
      - to distinguish between several timers

# 3: Giving up positioning after one re-try



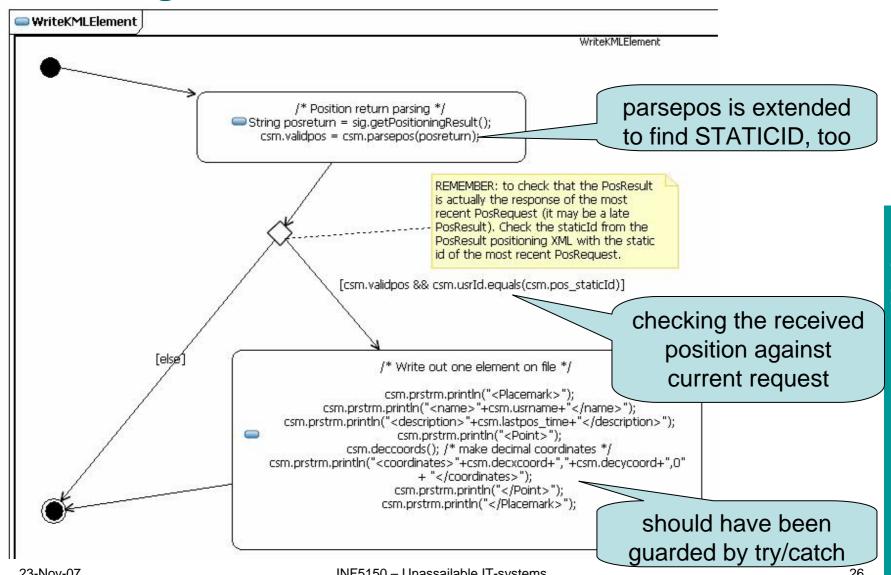


#### Timer expires

- To give PATS one more chance with this user
  - we need to define a variable to control re-positioning
- Having given up we must still cope with the PosResult coming later
  - This is more tricky than meets the eye since
    - when positioning is given up there is normally several FoundNick\_OK signals in the queue
    - and a late PosResult will follow those, but
    - that PosResult will come before any PosResult that is the result of new PosRequests
  - Thus, we must make sure that the PosResult is actually matched with the right nickname
    - We need to check the static id of the PosResult with that of the most recent PosRequest



# Checking the static id of the PosResult

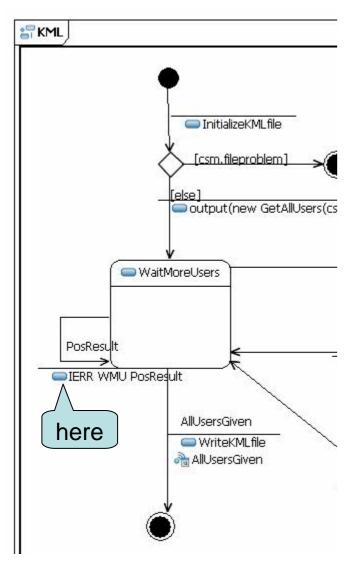


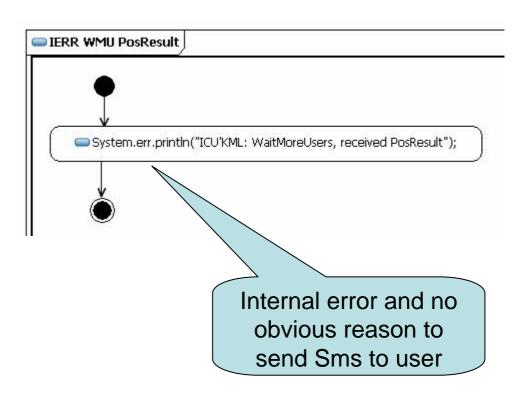


# KML problems (4 & 5)

- 4: Default transitions of WaitMoreUsers
  - PosResult and Sms are not handled
    - Sms cannot come to KML =>
      - internal error, handled on enclosing level
    - PosResult should not come =>
      - internal sequencing error, give message on console and ignore signal
- 5: Default transitions of WaitPosition
  - There are non-KML signals that should be covered (as Sms)
    - we will cover that on enclosing level

#### 4: PosResult received at WaitMoreUsers



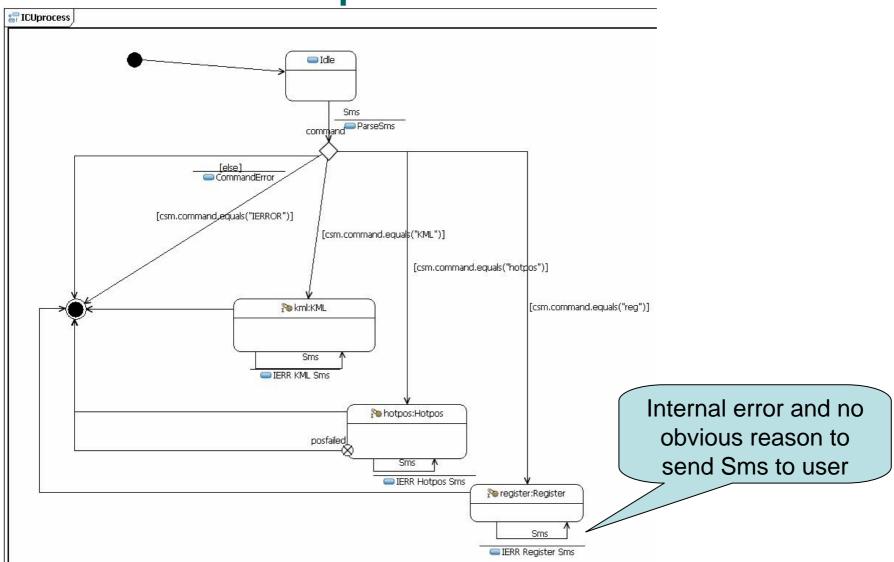


#### 4&5: Problems best solved on ICUprocess level

- The unexpected Sms signal
  - Neither KML, Hotpos or Register cater for receiving Sms
  - ...but they do not need to since Sms is always handled by ICUcontroller by creating a new ICUprocess
    - true, but will it always be that way?
- Covering the unexpected also makes the software more robust for the future
- The normal situation being that Sms will not occur in ICUprocess it may be handled on ICUprocess level



#### The modified ICUprocess





#### The exceptional

- Data may have strange syntax or values
  - We have looked at data checks for ICUcontroller
- An unexpected signal arrives
  - we explicitly describe every conceivable transition
  - We have looked at this for ICUprocess'KML
- No signal arrives
  - we guard our protocols/services with timers (ICUprocess'KML)
- Security issues
  - authentication + logging + statistics
  - Check for registration in ICUprocess'Hotpos
- Availability issues
  - self tests (we shall improve the Archive)



#### Services revisited

#### Hotpos

- Only registered users should be able to position others
- Positioning must be accepted by the positioned user
  - for the sake of showing more advance protocol for authentication

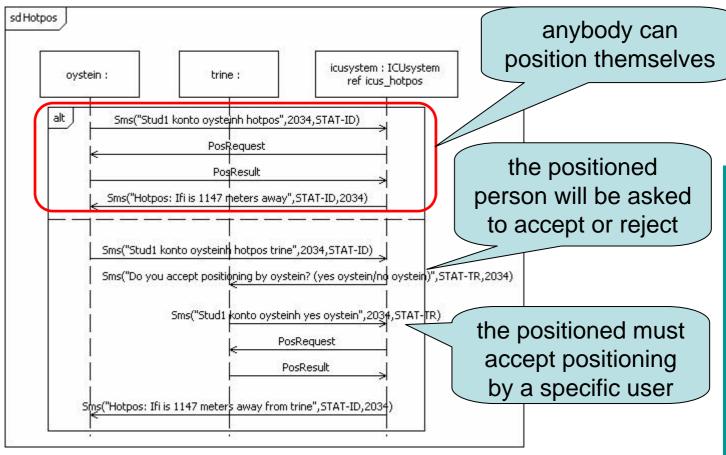
#### KML

- will not get the same full treatment
  - because asking every registered user is too tedious
- This shows that a "buddy group" concept probably needs to be introduced to continue to offer KML service

#### Register

– will of course not require that users are registered!

#### **Hotpos revisited**





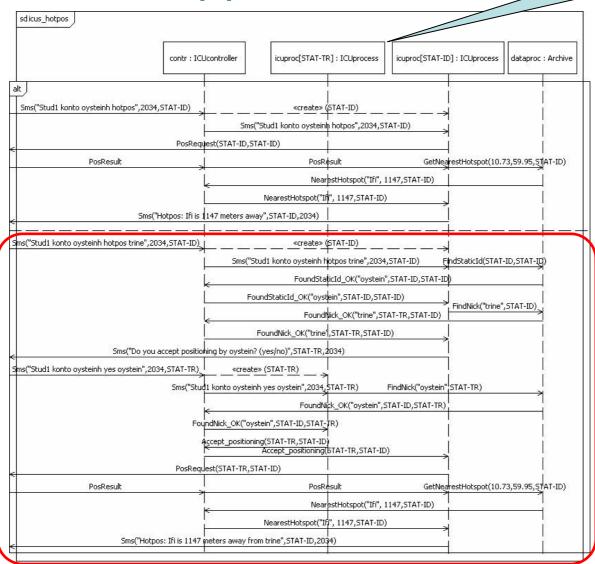
#### Problems when Trine should accept Oystein

- We need to know that Trine really accepts Oystein and not somebody else
  - we need to connect Trine's response to Oystein's session
- Trine's response is an Sms and that will in our design spawn another session!
  - which may not be a bad idea!
- Let us make a new service a yesno service
  - The yesno service will take an Sms with the following syntax:
    - "yes nickname" or "no nickname"
  - The yesno service will send a signal to the session identified by the nickname
    - Accept\_positioning or Reject\_positioning depending on yes/no



# Hotpos in detail (1)

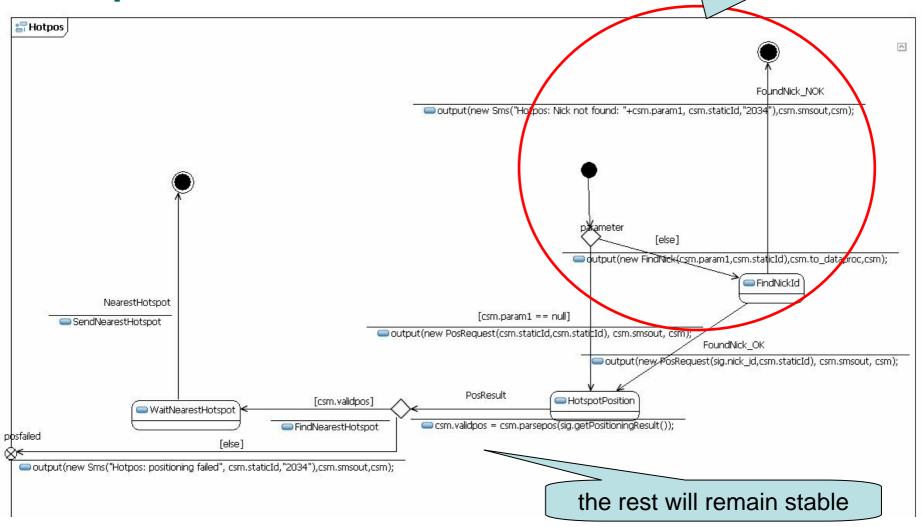
#### two sessions!





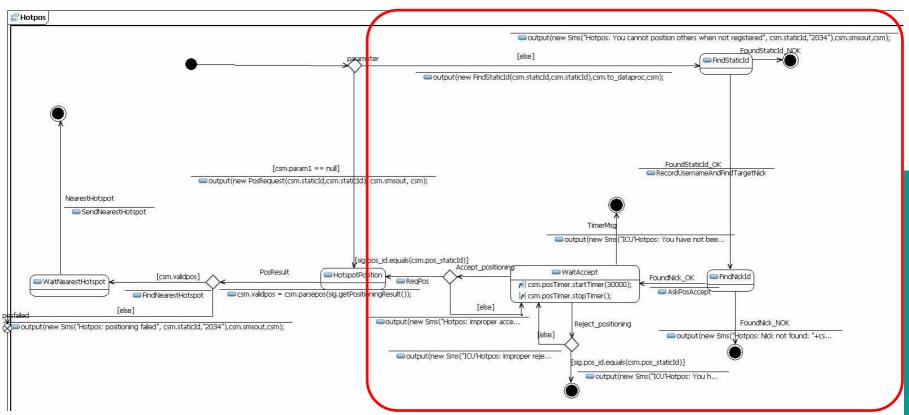
## Hotpos state machine in ICUB

#### This will change



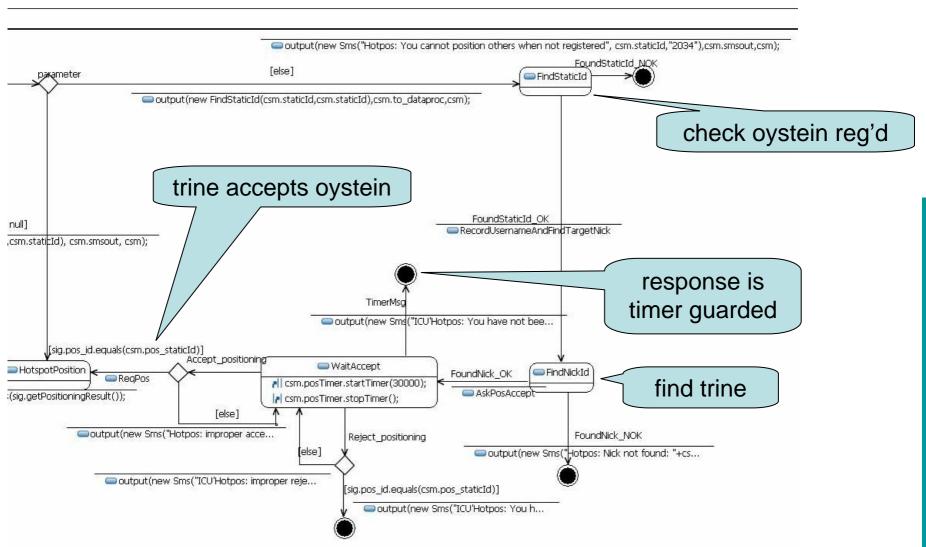


## **Hotpos in ICUC**



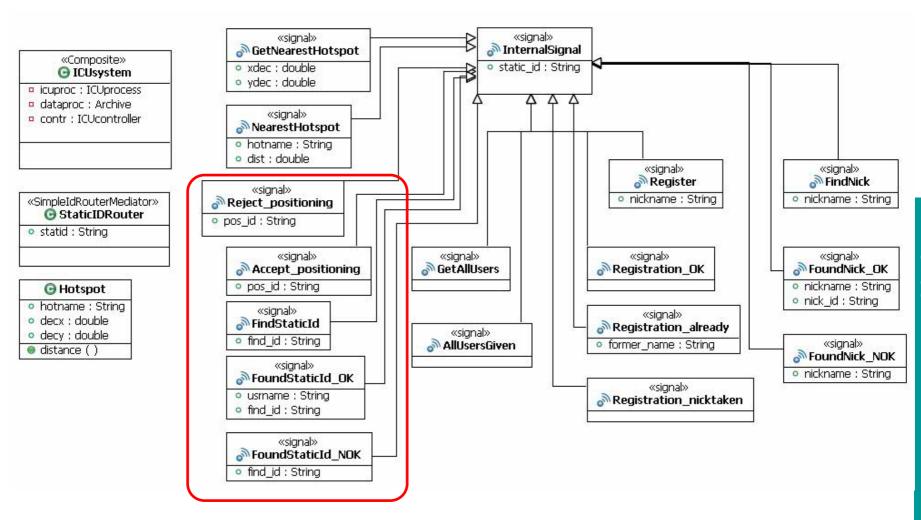


### Hotpos – the new features

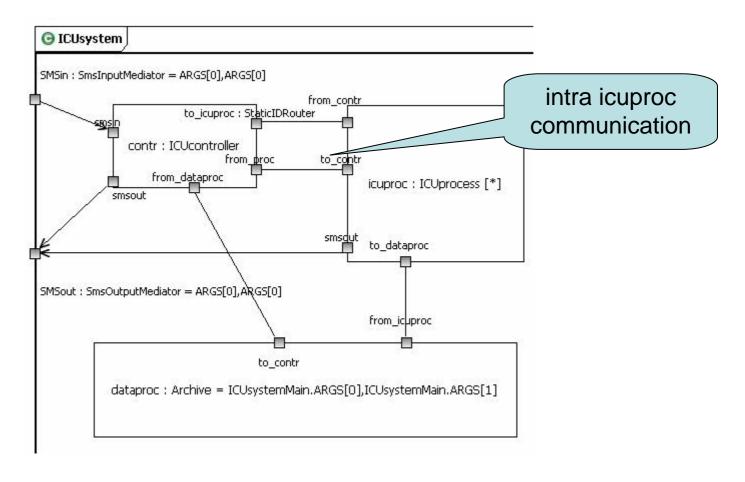




## **New internal signals**

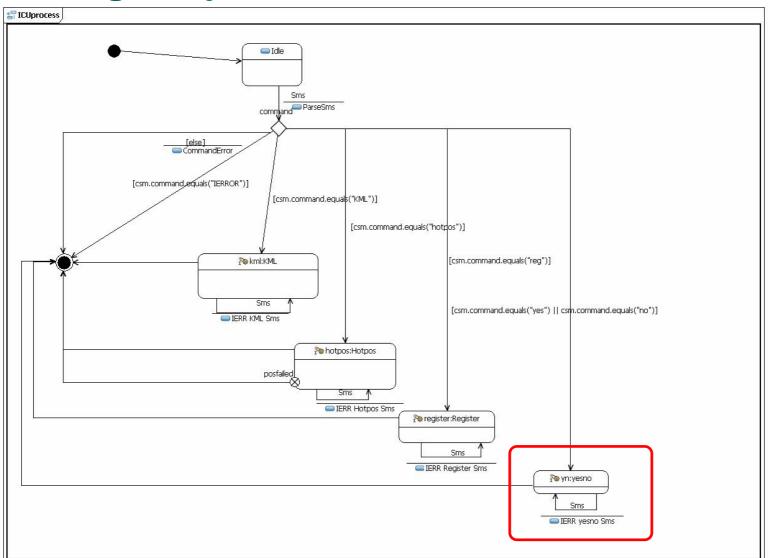


### New communication path must be added





### Adding the yesno service

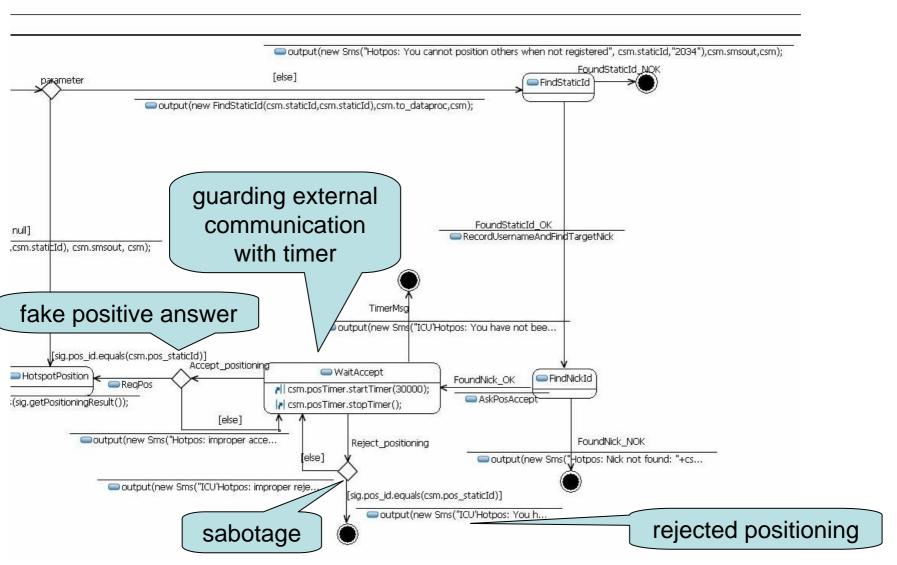




The yesno service decided not to check for registration due to possible 🚝 yesno reuse of the utility service output(new FindNick(csm.param1,csm.staticId),csm.to\_dataproc,csm); FoundNick NOK FindNick output(new Sms("Hotpos: Nick not found: "+csm.param1, csm.staticId,"2034"),csm.smsout,csm); FoundNick\_OK [else] output(new Reject\_positioning(csm.staticId,sig.nick\_id), csm.to\_contr,csm); [csm.command.equals("yes")] output(new Accept\_positioning(csm.staticId,sig.nick\_id), csm.to\_contr,csm);



### Hotpos – more issues





#### Points to make

- trivial additions to the Archive
  - finding the registered username from static id
- n+1
  - new signals introduced means new signals to cope with everywhere
- stability for parts of the state machine
  - emphasizing that a state is enough to determine the history
- services that use other services
  - Hotpos uses yes-no service
  - therefore we need new connection (and new ports) between icuproc and contr



### More points to make

- Guarding the external communication with a timer
  - WaitAccept where the positioning must be confirmed
- what about yes-no service?
  - out of protocol we must check on receiving side that the yes-no has the appropriate static id
    - since otherwise anybody could just send a "yes oystein" in place of the reply from Trine
  - also a reject must be checked against the static id
    - since otherwise anybody could just send a "no oystein" in place of the reply from Trine!
  - we will not require that yes-no needs registration
    - it is a utility, and may be used more freely at a later stage



### The exceptional applied to new Hotpos

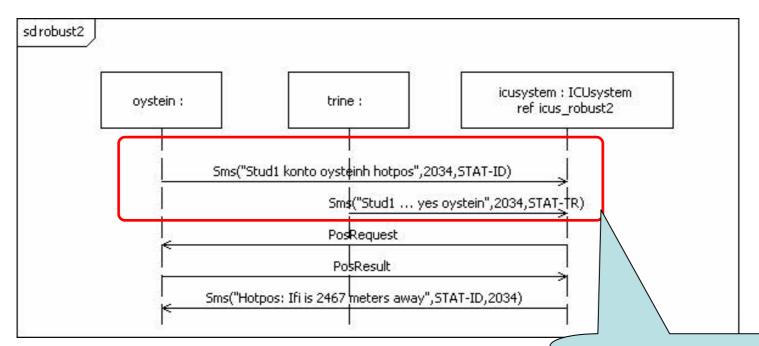
- Data may have strange syntax or values
  - checking static ids of the accept/reject messages
- An unexpected signal arrives
  - We should probably gone through the new signals everywhere
    - especially the accept and reject signals
- No signal arrives
  - we guard our external communication with timers (WaitAccept)
- Security issues
  - authentication (+ logging + statistics)
  - Check that user is registered
  - Check expected static id
    - prevents faked positive acceptance or negative service sabotage
  - Denial of service
    - keep faking will give resetting of the timer



#### n+1

- When we add functionality, we add signals
  - and those added signals should be covered in all states
  - in ICUC this is not the case!
- We have added external legal services yes and no
  - These services may produce internal signals Accept\_positioning or Reject\_positioning to other ICUprocesses
  - Those services may not be ready for those inputs!
    - if yes/no has been sent for no purpose or the nickname is misspelled
      - and the misspelled person really has a service going (rather improbable)

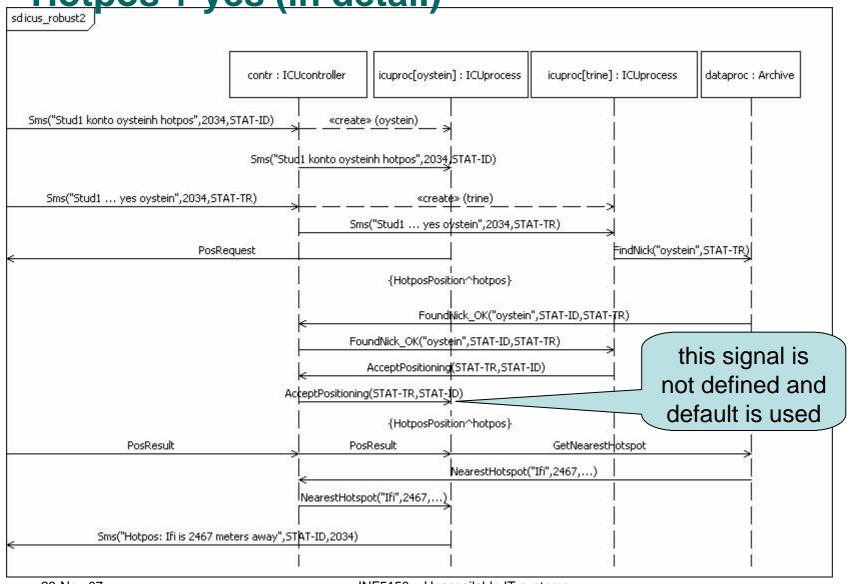
## Hotpos + yes (resulting in a default transition)



competing intiatives and yes is really wrong

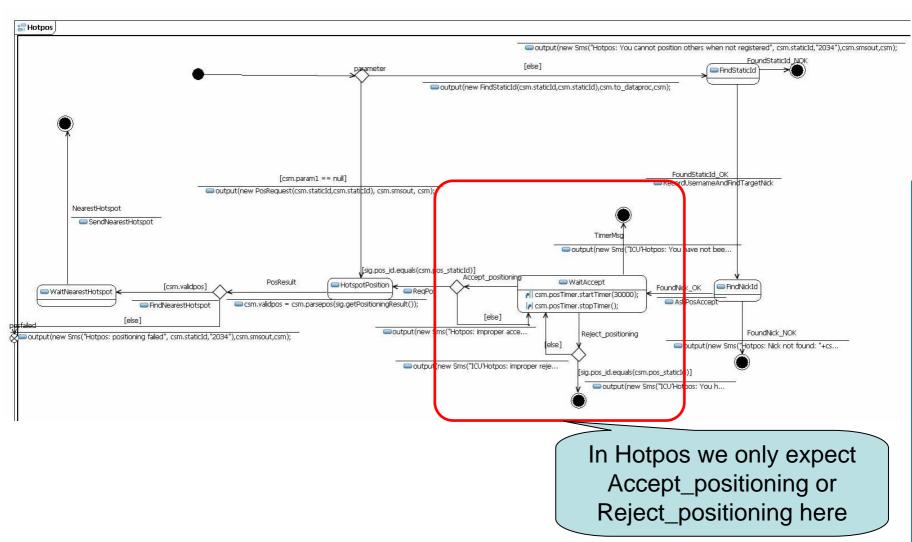


Hotpos + yes (in detail)

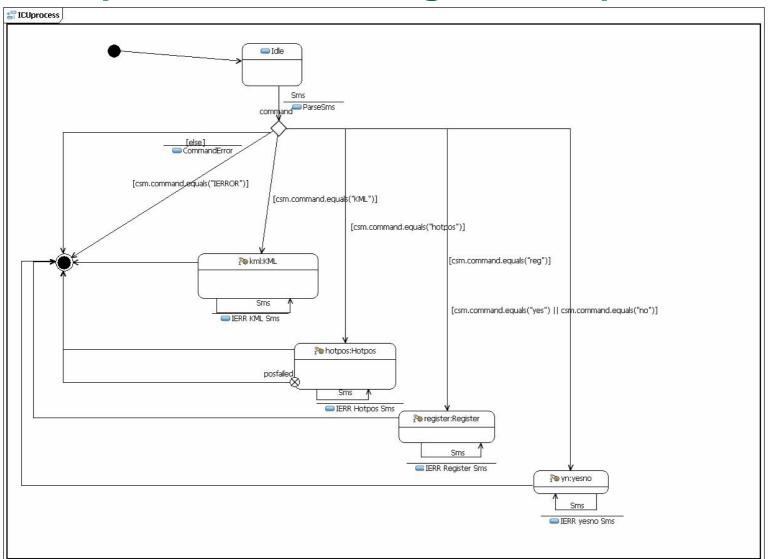




## **Misplaced Internal Signals**



## No capture of InternalSignal on top level

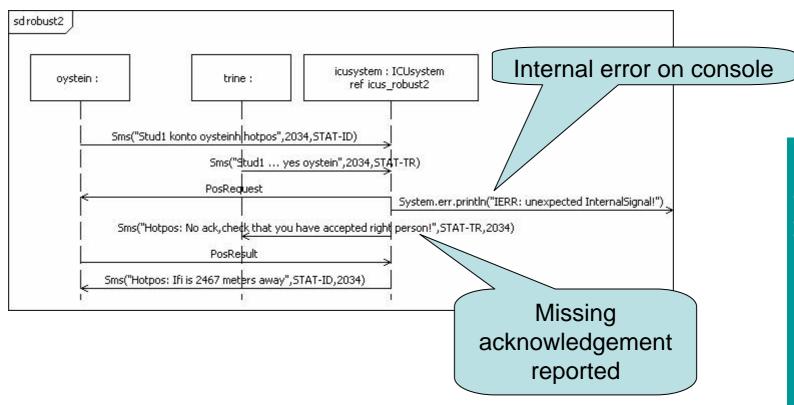




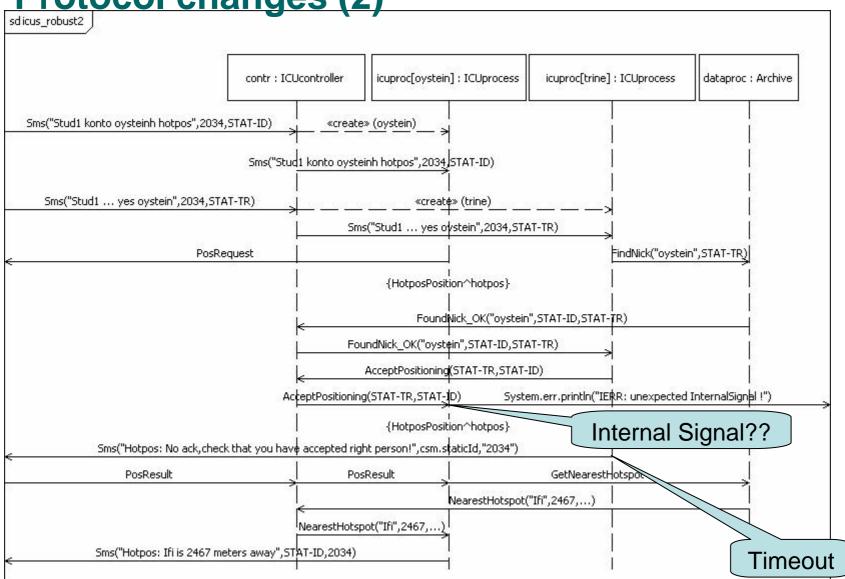
## Using the yes/no service

- Normal (sd ICUC'Hotpos):
  - oystein asks "hotpos trine"
  - trine accepts (or rejects) by saying "yes oystein" or "no oystein"
- Exceptional 1(sd ICUC'robust2):
  - oystein positions himself by "hotpos"
  - trine for no reason concurrently says "yes oystein"
- Exceptional 2:
  - oyvind asks "hotpos trine"
  - trine misreads oyvind's nick and says "yes oystein"
    - trine gets no message that her supposed acceptance fails!
    - oyvind will time out waiting for trine's approval
  - Possibly we should need double acknowledgment protocol
    - trine should be confirmed that her acceptance succeeded?!
    - or she should get an error message back when not acknowledged

## Protocol changes in detail (sequence diagrams)

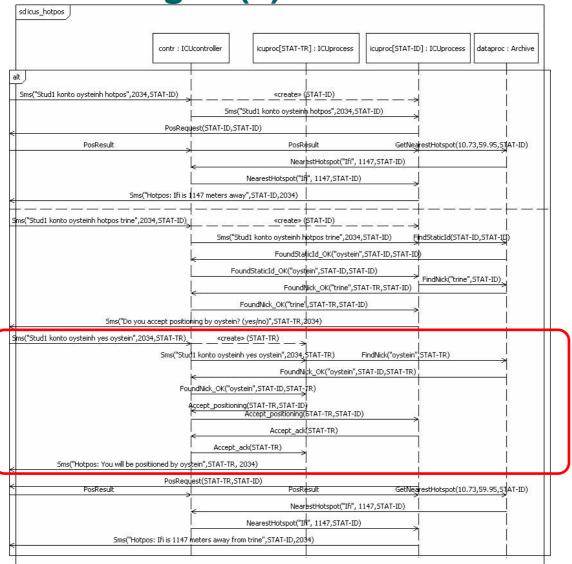


**Protocol changes (2)** 



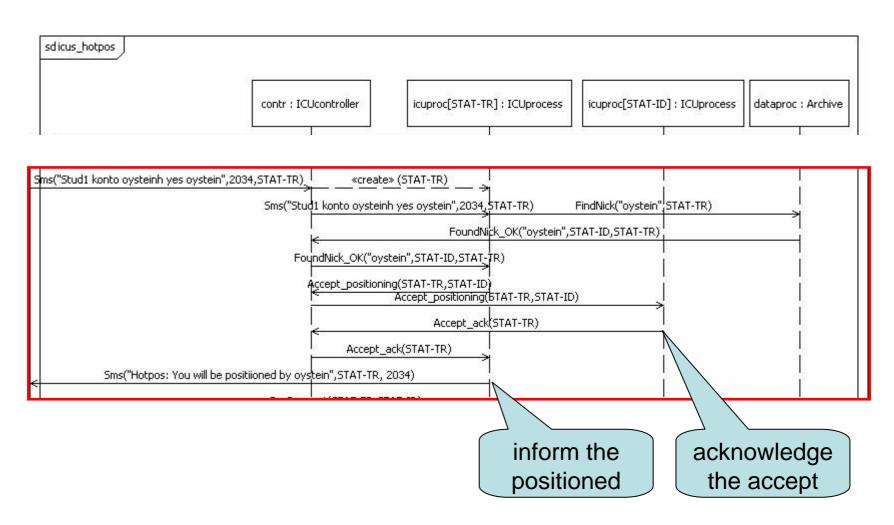


Protocol changes (3) – the normal Hotpos



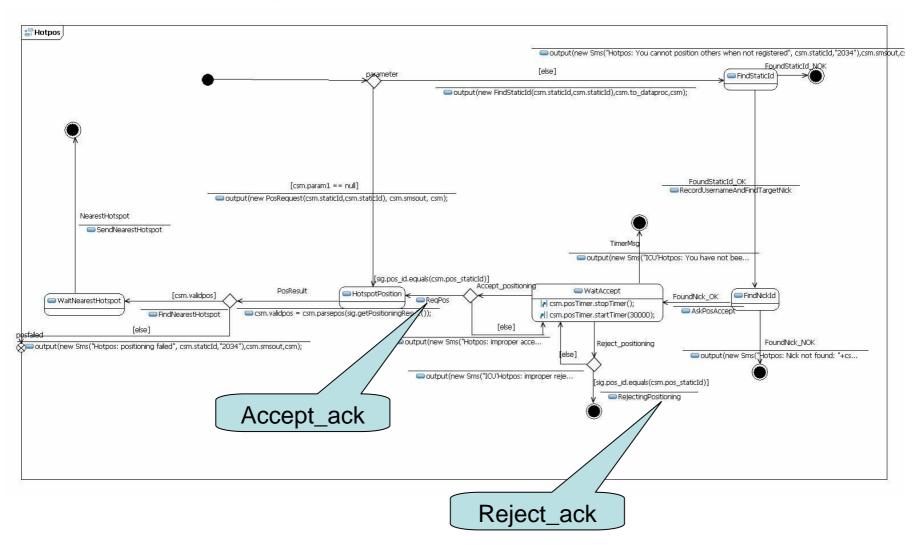


## Protocol changes (4) – the big view



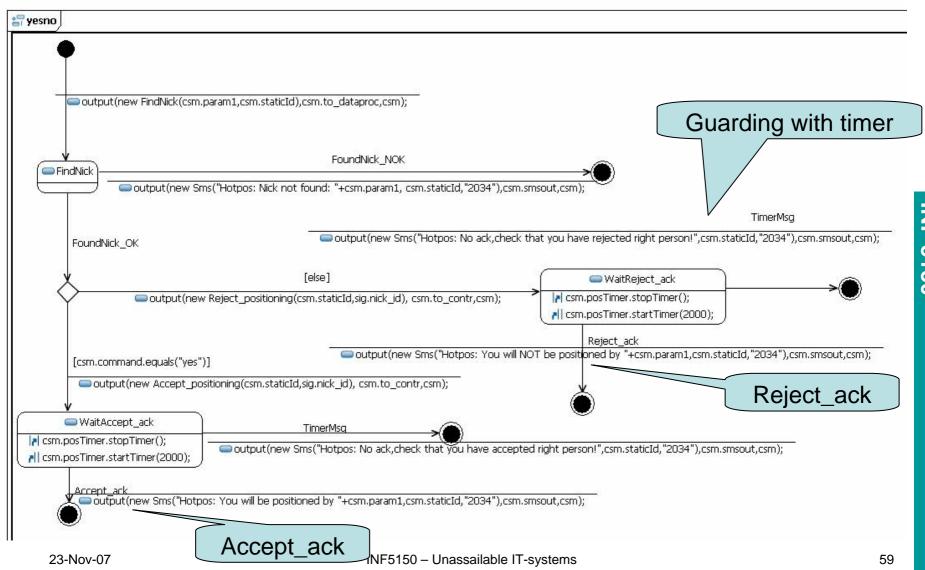


## Implementing the new protocol (Hotpos)



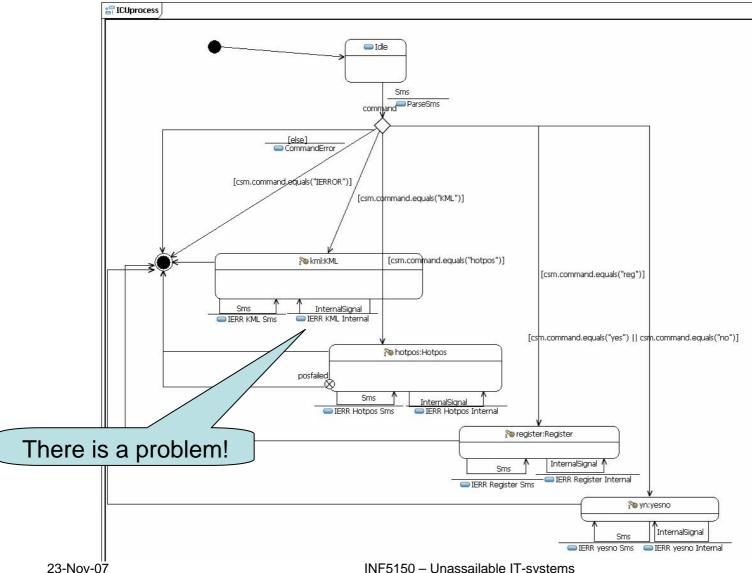


## Implementing the new protocol (yesno)





## Implementing the new protocol (ICUprocess)





## A story about history

- Unexpected signals can be caught on outer levels
  - but we want the net effect to be ignoring them after giving an error message
    - This can be done in UML with History states
- History states assures that when returning through a history point into a submachine state, execution will return where it left off in there
  - there are shallow and deep histories (one level, or all levels)
- UML has history states, but JavaFrame has not!
  - 1. implement history in JavaFrame? (not done yet)
  - 2. let transition return into the state anyway? (will restart the state)
  - 3. let transition end in a final state (terminating service which means that there is a way to perform denial of service)
  - 4. flatten the outer level error transitions into the inner levels



#### A lesson learned

- What you have not checked, may not work
  - We did not manage to check the Sms errors and therefore did not manage to discover the history problem
- What is defined in a standard, may not be implemented
  - History states are found in UML 2, but are not implemented in JavaFrame
- The optimal solution is not always obvious
  - 1. Implementing History states in JavaFrame
    - good for the future, but time-consuming now
  - 2. Restart the state
    - will also restart the service and that is not in general attractive
  - 3. Terminate
    - simple solution that actually hurts an innocent user
  - 4. Flatten the transitions down
    - not very elegant, but requires only finite time to do
    - not very future-oriented

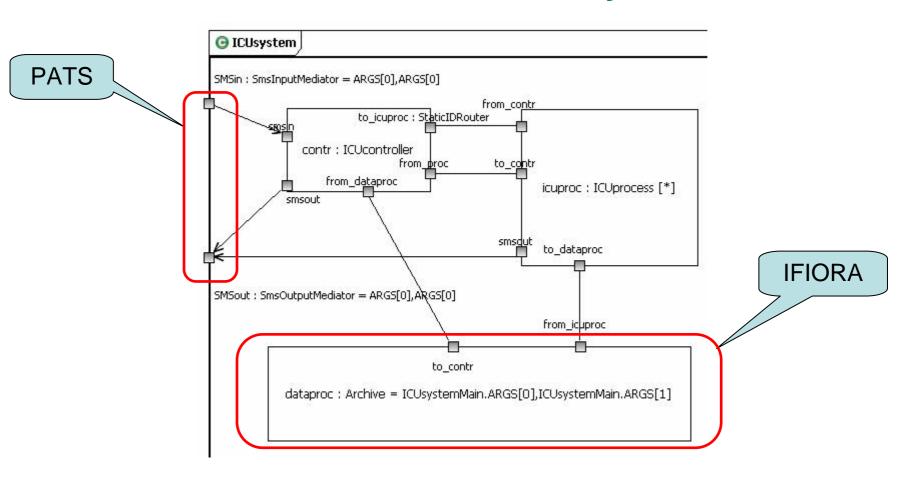


## **Availability**

- Availability
  - That authorized users can get the services they want when they want them
- It may be too late to check the availability when the service is being asked for
  - It may be necessary to check regularly regardless of demand
- External resources upon which the service depends
  - should be checked regularly
- Internal resources
  - may be trusted as they may only be divisions of the program
  - may be checked if they involve external resource (like network)



# **External resources of the ICU system**





#### **PATS**

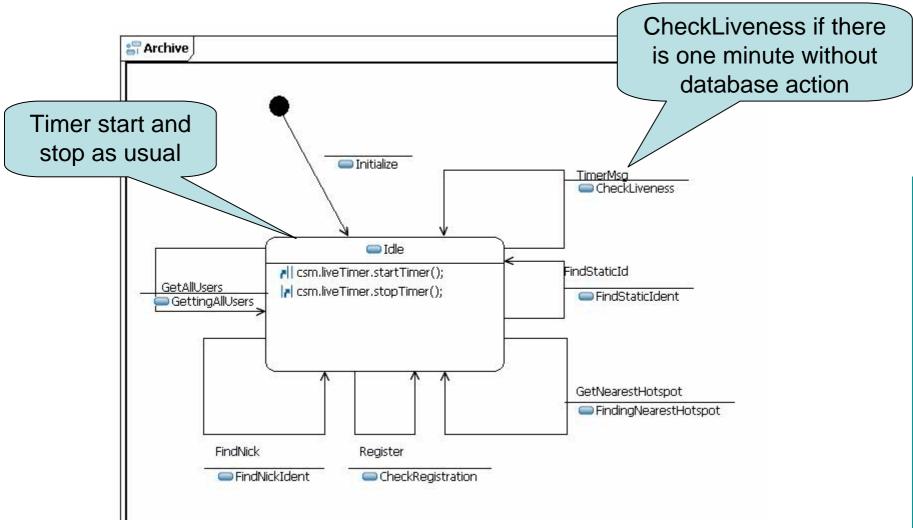
- The connection to PATS is controlled by the IFI lower level software
  - This is not always enough to make sure that PATS really works the way it is expected for our purposes
- In a normal situation there will be frequent requests to PATS and malfunction would be reported through the robustness means that we have already applied
  - If PATS connection is dead, nothing would reach our program
- Extra liveness checks would actually cost money (for commercial utilization of PATS)
- For ICU we decide not to introduce extra liveness checks against PATS

#### IFIORA – the IFI Oracle database server

- IFIORA will also be invoked frequently and failure reported through the exception handling
  - which should be improved from stack dump!
- For the sake of demonstration we also include a liveness check for IFIORA
  - We assume that the exception handling implicit in jdbc will always capture availability exceptions
  - An extra liveness check will be implemented through a regular timer-driven transition that performs a simple SQL-command
  - An availability exception will be reported back to the calling service through a special internal error signal (DataError)
    - on which the service may react by issuing a message to the user
- Many small cascading effects around in the model



#### Archive – with added liveness timer



### **CheckLiveness**

```
/* Liveness check by performing the simplest kind of SQL command */
try {
       Statement stmt = csm.con.createStatement();
       String theQuery = "SELECT COUNT(*) FROM gsmuser";
       ResultSet r = stmt.executeQuery(theQuery);
 catch (Exception e) {
  System.err.println("ICU'Archive: Liveness check fails! Reconnecting!");
  try {
       DriverManager.registerDriver(new oracle.jdbc.driver.OracleDriver())
       Properties props = new Properties();
       props.put("user", csm.oracleAccount);
       props.put("password", csm.oraclePasswd);
       String url = "jdbc.:oracle:thin:@delphinium.ifi.uio.no:1521:IFIORA";
       csm.con = DriverManager.getConnection(url, props);
   catch (SQLException ee) {
       System.err.println("ICU'Archive: Error when reconnecting!");
```

cheap SQL statement

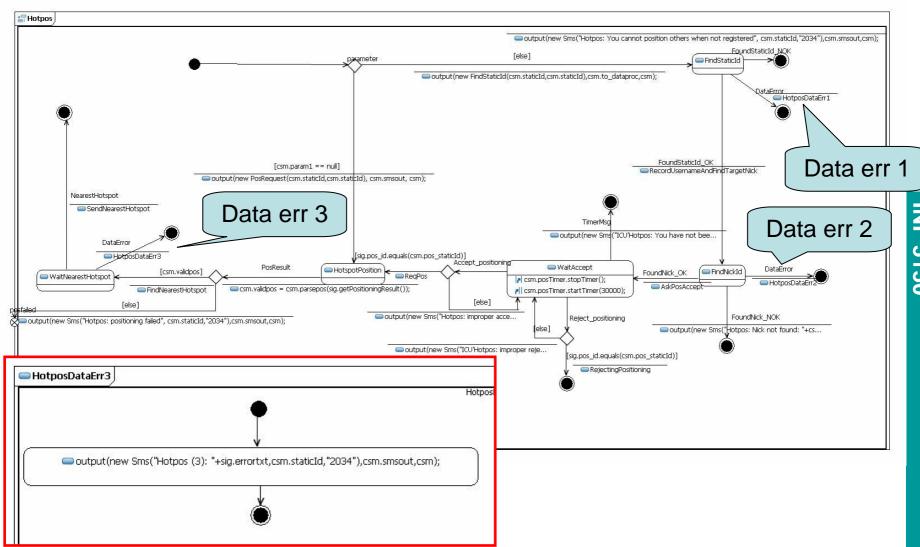
must reconnect to work again



#### **FindStaticIdent**

```
/* look in gsmusers for static id */
try {
  Statement stmt = csm.con.createStatement();
  String theQuery = "SELECT * FROM gsmuser WHERE staticid = "+ sig.find_id + "";
  ResultSet r = stmt.executeQuery(theQuery);
  if (r.next())
  { /* Static id found*/
     output(new FoundStaticId_OK(r.getString("nickname"),sig.find_id,sig.static_id),
     csm.to_contr,csm);
                                                Double error messages: to the
  else
                                                console and the calling service
  { /* Static id not found */
     output(new FoundStaticId_NOK(sig.find_id,
                                                        rc_id),csm.to_contr,csm);
} catch (Exception e) {
 System.err.println("ICU'Archive: Error when Selecting staticid from gsmuser");
 output(
  new DataError("ICU'Archive: Error when Selecting staticid from gsmuser",
  sig.static_id), csm.to_contr, csm);
```

# Catching the DataError message in Hotpos



#### The robustification summarized

- Data may have strange syntax or values
  - We have looked at data checks for ICUcontroller
- An unexpected signal arrives
  - we explicitly describe every conceivable transition
  - We will look at this again for "n+1" situation
- No signal arrives
  - we guard our protocols/services with timers
- Security issues
  - authentication + logging + statistics
  - Check for registration in ICUprocess'Hotpos
- Availability issues
  - liveness tests (Archive)



### What more robustification could we have done?

- KML and yesno are still without authentication
  - in practice we would need a "buddy" concept
- PATS is not checked
  - we could have covered sending Sms/PosRequest
  - probably best on lower level, but would cause some problems
- We have not tested every peculiar (but imagined) situation
  - because it is difficult/tedious to do
  - will require a very precise testing environment
- Probably should have had one more iteration of cleaning up the diagrams
  - aesthetics is important for understanding