Introduction to Databases

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Overview of this module

1. Today: Introduction to databases and the relational model
2. Next week: Basic SQL (answer queries)

Curriculum

- The curriculum of this module are the slides from the lectures
- the weekly exercises with solutions given on the semester page
- The mandatory assignment (will be published 19. october)
- The book *SQL Queries For Mere Mortals* should be used as supplement to the slides for more in-depth explanations, and more examples and exercises

Motivation

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All of these problems are solved by the filesystem!

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Python + Files

```
import csv
import os

filea = "a.csv"
fileb = "b.csv"
temp = "temp.csv"
source1 = csv.reader(open(filea,"r"),delimiter=",")
source2 = csv.reader(open(fileb,"r"),delimiter=",")

source2_dict = {}
for row in source2:
    source2_dict[row[0]] = row[1]

with open(temp, "w") as fout:
    csvwriter = csv.writer(fout, delimiter=delim)
    for row in source1:
        if row[1] in source2_dict:
            row[3] = source2_dict[row[1]]
        csvwriter.writerow(row)
os.rename(temp, filea)
```

SQL + Database

```
UPDATE a
   SET c4=b.c2
  FROM b
 WHERE a.c2 = b.c1;
```

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  - Database uses advanced techniques to find the most efficient way to execute queries
  - Also uses advanced data structures to store data for efficient retrieval

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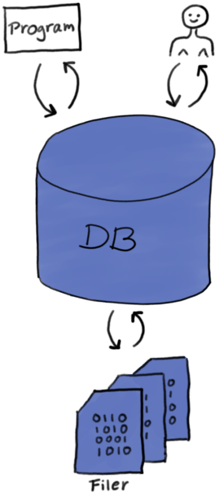

Database functions as an abstraction layer over the filsystem

- Makes it easier to search and manipulate data
- Easier to specify structure of the data
- More efficient and scalable

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- Relational databases: Stores tables (or relations) consisting of rows and columns
- We will focus on relational databases, the most used type of database

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  - a name,
  - a collection of columns
  - and a collection of rows (which is the data)
- A column has
  - a name,
  - and a type

Tables/Relations

Example table:

Patient

PatientID (int)	Name (text)	Birthdate (date)	BloodPressure (text)
0	Anna Consuma	1978-10-09	123/75
1	Peter Young	2009-03-01	150/81
2	Carla Smith	1986-06-14	101/53
3	Sam Penny	1961-01-09	127/82
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  - a patient must have `PatientID`, `Name`, `Birthdate`, `BloodPressure`

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- The most used type of database

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- A database can have multiple such database schemas, and each schema has a name
- Schemas are used to group related tables together (e.g. one schema for tables related to patients, one schema for tables related to hospitals, etc.)

Example Database

Academia

BI

Employees

EmployeeID	Name	Startdate
0	Peter Svensen	01-03-1999
1	Kari Smith	11-04-1977
2	Petrine Lye	07-01-2002
⋮	⋮	⋮

Students

StudentID	Name	Level
0	Ove Persson	Bachelor
1	Ingrid Olava	Master
2	Marge Smith	Bachelor
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Courses

CourseID	Name	Level
0	Analytics	Master
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UIO

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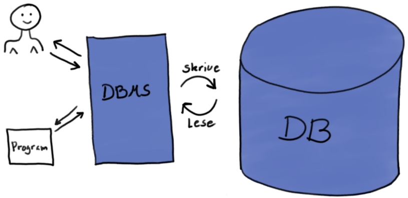
- easy access from programming languages (like Python)
- systems for security and control of who has access to the data
- systems that secure the integrity of the data
- support for much larger volumes of data and much more complex structure

Database systems

- A database is really just a collection of data (not a system/program)

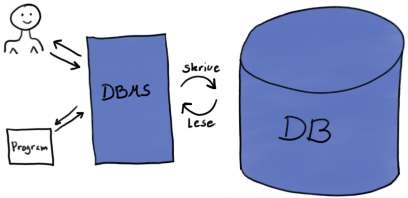
Database systems



- A database is really just a collection of data (not a system/program)
- A database management system (DBMS) is

  a system that let users define, create, maintain and control access to data.

Database systems



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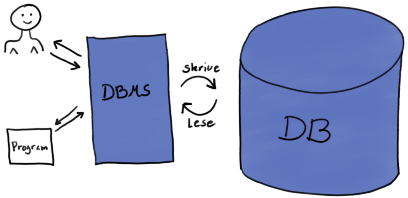
  a database management system over relational databases.

Database systems



- A database is really just a collection of data (not a system/program)
- A database management system (DBMS) is

  a system that let users define, create, maintain and control access to data.

- A relational database manegement system (RDBMS) is

  a database management system over relational databases.

- Often use the word “database” for both data, program, and the combaintion of these

Schema violations

The database system will not allow you to insert values violating the database schema.

Thus, the following is not allowed (errors marked in red):

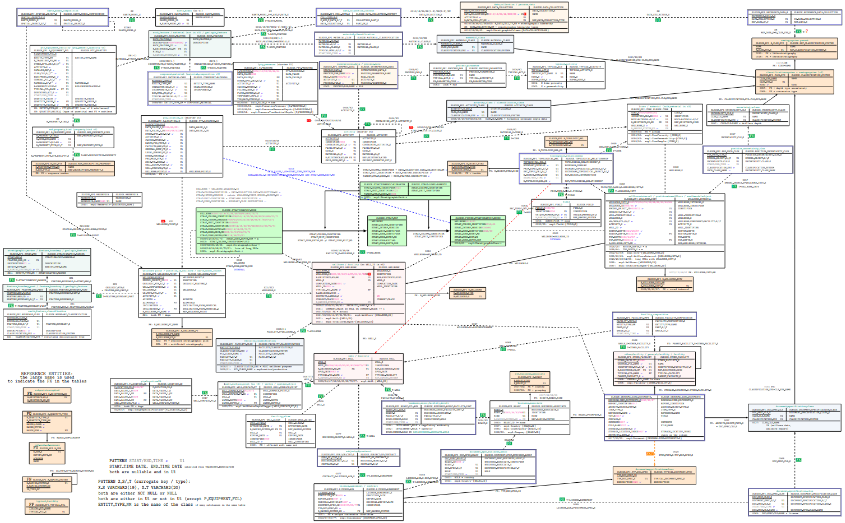
Patient

PatientID (int)	Name (text)	Birthdate (date)	BloodPressure (text)
0	Anna Consuma	1978-10-09	123/75
1	Peter Young	2009-03-01	150/81
2	2	1986-06-14	101/53
3	Sam Penny	long ago	127/82
four	John Mill	1989-11-16	147/92
5 6	Yvonne Potter	1971-04-12	122/74

Database design

- So relational databases store data as tables with a rigid structure
- But how should we represent information as data in tables?
- The structure of the data, i.e. which tables and columns we make, affects how easy it is to use and maintain the data
- Need to have a good *database design*

Complex database schemas



Database design: Blood pressure

Assume we want to keep track of pasient's blood pressure over time. We could then make a table looking like this:

Patient

PatientID	Name	Birthdate	Telephone	BloodPressure	TestTime
0	Anna Consuma	1978-10-09	12345678	123/75	2022-09-23
1	Peter Young	2009-03-01	21679921	150/81	2022-09-20
2	Carla Smith	1986-06-14	98765432	101/53	2022-08-07
3	Sam Penny	1961-01-09	91827364	127/82	2022-09-28
4	John Mill	1989-11-16	56473829	147/92	2022-09-13
5	Yvonne Potter	1971-04-12	91298833	122/74	2022-09-04

Blood pressure: More tests

Patient

PatientID	Name	Birthdate	Telephone	BloodPressure	TestTime
0	Anna Consuma	1978-10-09	12345678	123/75	2022-09-23
1	Peter Young	2009-03-01	21679921	150/81	2022-09-20
2	Carla Smith	1986-06-14	98765432	101/53	2022-08-07
3	Sam Penny	1961-01-09	91827364	127/82	2022-09-28
4	John Mill	1989-11-16	56473829	147/92	2022-09-13
5	Yvonne Potter	1971-04-12	91298833	122/74	2022-09-04
0	Anna Consuma	1978-10-09	12345678	125/73	2022-10-01
1	Peter Young	2009-03-01	21679921	143/80	2022-10-03
4	John Mill	1989-11-16	56473829	146/92	2022-10-03
5	Yvonne Potter	1971-04-12	91298833	124/75	2022-10-04
0	Anna Consuma	1978-10-09	12345678	126/74	2022-10-05
3	Sam Penny	1961-01-09	91827364	126/80	2022-10-08
1	Peter Young	2009-03-01	21679921	141/79	2022-10-11

Problems with bad design

- Difficult to maintain data
  - If a patient changes name or phone number, need to change multiple rows
- Difficult to add data
  - Cannot insert new patient without also inserting blood pressure and testtime
- Duplicate data takes up more disk space and is slower to work with

Blood pressure: (Failed) attempt at better structure

Patient

PatientID	Name	Birthdate	Telephone	BloodPressure
0	Anna Consuma	1978-10-09	12345678	(123/75, 2022-09-23),(125/73, 2022-10-01),...
1	Peter Young	2009-03-01	21679921	(150/81, 2022-09-20),(143/80, 2022-10-03),...
2	Carla Smith	1986-06-14	98765432	(101/53, 2022-08-07)
3	Sam Penny	1961-01-09	91827364	(127/82, 2022-09-28),(126/80, 2022-10-08)
4	John Mill	1989-11-16	56473829	(147/92, 2022-09-13),(146/92, 2022-10-03)
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2	Carla Smith	1986-06-14	98765432	(101/53, 2022-08-07)
3	Sam Penny	1961-01-09	91827364	(127/82, 2022-09-28),(126/80, 2022-10-08)
4	John Mill	1989-11-16	56473829	(147/92, 2022-09-13),(146/92, 2022-10-03)
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- Blood pressure values now contained deep inside a single value

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2	Carla Smith	1986-06-14	98765432	(101/53, 2022-08-07)
3	Sam Penny	1961-01-09	91827364	(127/82, 2022-09-28),(126/80, 2022-10-08)
4	John Mill	1989-11-16	56473829	(147/92, 2022-09-13),(146/92, 2022-10-03)
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- Blood pressure values now contained deep inside a single value
- Need to ”parse”/”unwrap” this complex value to get blood pressure values

Blood pressure: (Failed) attempt at better structure

Patient

PatientID	Name	Birthdate	Telephone	BloodPressure
0	Anna Consuma	1978-10-09	12345678	(123/75, 2022-09-23),(125/73, 2022-10-01),...
1	Peter Young	2009-03-01	21679921	(150/81, 2022-09-20),(143/80, 2022-10-03),...
2	Carla Smith	1986-06-14	98765432	(101/53, 2022-08-07)
3	Sam Penny	1961-01-09	91827364	(127/82, 2022-09-28),(126/80, 2022-10-08)
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- Blood pressure values now contained deep inside a single value
- Need to ”parse”/”unwrap” this complex value to get blood pressure values
- Makes working with these values very complex (both for humans and computer)

Blood pressure: (Failed) attempt at better structure

Patient

PatientID	Name	Birthdate	Telephone	BloodPressure
0	Anna Consuma	1978-10-09	12345678	(123/75, 2022-09-23),(125/73, 2022-10-01),...
1	Peter Young	2009-03-01	21679921	(150/81, 2022-09-20),(143/80, 2022-10-03),...
2	Carla Smith	1986-06-14	98765432	(101/53, 2022-08-07)
3	Sam Penny	1961-01-09	91827364	(127/82, 2022-09-28),(126/80, 2022-10-08)
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5	Yvonne Potter	1971-04-12	91298833	(122/74, 2022-09-04),(124/75, 2022-10-04)

- Blood pressure values now contained deep inside a single value
- Need to ”parse”/”unwrap” this complex value to get blood pressure values
- Makes working with these values very complex (both for humans and computer)
- Generally: Columns should have simple values!

Blood pressure: Better structure

Patient

PatientID	Name	Birthdate	Telephone
0	Anna Consuma	1978-10-09	12345678
1	Peter Young	2009-03-01	21679921
2	Carla Smith	1986-06-14	98765432
3	Sam Penny	1961-01-09	91827364
4	John Mill	1989-11-16	56473829
5	Yvonne Potter	1971-04-12	91298833

BloodPressure

PatientID	BloodPressure	TestTime
0	123/75	2022-09-23
1	150/81	2022-09-20
2	101/53	2022-08-07
3	127/82	2022-09-28
4	147/92	2022-09-13
5	122/74	2022-09-04
0	125/73	2022-10-01
1	143/80	2022-10-03
4	146/92	2022-10-03
5	124/75	2022-10-04
0	126/74	2022-10-05
3	126/80	2022-10-08
1	141/79	2022-10-11

Students and courses

Want to store information about students, courses and grades:

Students and courses

Want to store information about students, courses and grades:

- For students: Username, name, surename, address...
- For courses: Coursecode, title, description, credits...
- Grades: Which student got which grade in which course

Students and courses

Want to store information about students, courses and grades:

- For students: Username, name, surename, address...
- For courses: Coursecode, title, description, credits...
- Grades: Which student got which grade in which course

Naive solution: Everything in one table!

Students and courses: Schema

StudentCourse

Username	Name	Surename	Address	Coursecode	Title	Desc.	Credits	Grade
evgenit	Evgenij	Thorstensen	Addr1	IN2090	Databaser	Beskr...	10	B
peternl	Petter	Nilsen	Addr2	IN2090	Databaser	Beskr...	10	A
evgenit	Evgenij	Thorstensen	Addr1	IN2080	Beregn...	Descr...	10	A
leifhka	Leif H.	Karlsen	Addr3	IN2090	Databaser	Beskr...	10	B
leifhka	Leif H.	Karlsen	Addr3	IN3110	Program...	Desc2...	5	C

Insert and delete

StudentCourse

Username	Name	Surename	Address	Coursecode	Title	Desc.	Credits	Grade
evgenit	Evgenij	Thorstensen	Addr1	IN2090	Databaser	Beskr...	10	B
peternl	Petter	Nilsen	Addr2	IN2090	Databaser	Beskr...	10	A
evgenit	Evgenij	Thorstensen	Addr1	IN2080	Beregn...	Descr...	10	A
leifhka	Leif H.	Karlsen	Addr3	IN2090	Databaser	Beskr...	10	B
leifhka	Leif H.	Karlsen	Addr3	IN3110	Program...	Desc2...	5	C

Data duplication makes it more difficult to insert and update data:

Insert and delete

StudentCourse

Username	Name	Surename	Address	Coursecode	Title	Desc.	Credits	Grade
evgenit	Evgenij	Thorstensen	Addr1	IN2090	Databaser	Beskr...	10	B
peternl	Petter	Nilsen	Addr2	IN2090	Databaser	Beskr...	10	A
evgenit	Evgenij	Thorstensen	Addr1	IN2080	Beregn...	Descr...	10	A
leifhka	Leif H.	Karlsen	Addr3	IN2090	Databaser	Beskr...	10	B
leifhka	Leif H.	Karlsen	Addr3	IN3110	Program...	Desc2...	5	C

Data duplication makes it more difficult to insert and update data:

- Need to insert all the info about student and course, even if only want to insert a new grade

Insert and delete

StudentCourse

Username	Name	Surename	Address	Coursecode	Title	Desc.	Credits	Grade
evgenit	Evgenij	Thorstensen	Addr1	IN2090	Databaser	Beskr...	10	B
peternl	Petter	Nilsen	Addr2	IN2090	Databaser	Beskr...	10	A
evgenit	Evgenij	Thorstensen	Addr1	IN2080	Beregn...	Descr...	10	A
leifhka	Leif H.	Karlsen	Addr3	IN2090	Databaser	Beskr...	10	B
leifhka	Leif H.	Karlsen	Addr3	IN3110	Program...	Desc2...	5	C

Data duplication makes it more difficult to insert and update data:

- Need to insert all the info about student and course, even if only want to insert a new grade
- Impossible to insert a new student, without also inserting a course

Insert and delete

StudentCourse

Username	Name	Surename	Address	Coursecode	Title	Desc.	Credits	Grade
evgenit	Evgenij	Thorstensen	Addr1	IN2090	Databaser	Beskr...	10	B
peternl	Petter	Nilsen	Addr2	IN2090	Databaser	Beskr...	10	A
evgenit	Evgenij	Thorstensen	Addr1	IN2080	Beregn...	Descr...	10	A
leifhka	Leif H.	Karlsen	Addr3	IN2090	Databaser	Beskr...	10	B
leifhka	Leif H.	Karlsen	Addr3	IN3110	Program...	Desc2...	5	C
				IN9999	Quantum	Beskr3	10	

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peternl	Petter	Nilsen	Addr2	IN2090	Databaser	Beskr...	10	A
evgenit	Evgenij	Thorstensen	Addr1	IN2080	Beregn...	Descr...	10	A
leifhka	Leif H.	Karlsen	Addr3	IN2090	Databaser	Beskr...	10	B
leifhka	Leif H.	Karlsen	Addr3	IN3110	Program...	Desc2...	5	C
				IN9999	Quantum	Beskr3	10	
abcdef	Aber C.	Deflan	Addr4					

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Insert and delete

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peternl	Petter	Nilsen	Addr2	IN2090	Databaser	Beskr...	10	A
evgenit	Evgenij	Thorstensen	Addr1	IN2080	Beregn...	Descr...	10	A
leifhka	Leif H.	Karlsen	Addr3	IN2090	Databaser	Beskr...	10	B
leifhka	Leif H.	Karlsen	Addr3	IN3110	Program...	Desc2...	5	C
				IN9999	Quantum	Beskr3	10	
abcdef	Aber C.	Deflan	Addr4					

Data duplication makes it more difficult to insert and update data:

- Need to insert all the info about student and course, even if only want to insert a new grade
- Impossible to insert a new student, without also inserting a course
- Updates must be performed consistently and update all duplicates

Anomalier: Sletting

StudentCourse

Username	Name	Surename	Address	Coursecode	Title	Desc.	Credits	Grade
evgenit	Evgenij	Thorstensen	Addr1	IN2090	Databaser	Beskr...	10	B
peternl	Petter	Nilsen	Addr2	IN2090	Databaser	Beskr...	10	A
evgenit	Evgenij	Thorstensen	Addr1	IN2080	Beregn...	Descr...	10	A
leifhka	Leif H.	Karlsen	Addr3	IN2090	Databaser	Beskr...	10	B
leifhka	Leif H.	Karlsen	Addr3	IN3110	Program...	Beskr2...	5	C

- Deleting a course may delete a student

Anomalier: Sletting

StudentCourse

Username	Name	Surename	Address	Coursecode	Title	Desc.	Credits	Grade
evgenit	Evgenij	Thorstensen	Addr1	IN2090	Databaser	Beskr...	10	B
peternl	Petter	Nilsen	Addr2	IN2090	Databaser	Beskr...	10	A
evgenit	Evgenij	Thorstensen	Addr1	IN2080	Beregn...	Descr...	10	A
leifhka	Leif H.	Karlsen	Addr3	IN2090	Databaser	Beskr...	10	B
leifhka	Leif H.	Karlsen	Addr3	IN3110	Program...	Beskr2...	5	C

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evgenit	Evgenij	Thorstensen	Addr1	IN2090	Databaser	Beskr...	10	B
peternl	Petter	Nilsen	Addr2	IN2090	Databaser	Beskr...	10	A
evgenit	Evgenij	Thorstensen	Addr1	IN2080	Beregn...	Descr...	10	A
leifhka	Leif H.	Karlsen	Addr3	IN2090	Databaser	Beskr...	10	B
leifhka	Leif H.	Karlsen	Addr3	IN3110	Program...	Beskr2...	5	C

- Deleting a course may delete a student
- Deleting a student may delete a course

Anomalier: Sletting

StudentCourse

Username	Name	Surename	Address	Coursecode	Title	Desc.	Credits	Grade
evgenit	Evgenij	Thorstensen	Addr1	IN2090	Databaser	Beskr...	10	B
peternl	Petter	Nilsen	Addr2	IN2090	Databaser	Beskr...	10	A
evgenit	Evgenij	Thorstensen	Addr1	IN2080	Beregn...	Descr...	10	A
leifhka	Leif H.	Karlsen	Addr3	IN2090	Databaser	Beskr...	10	B
leifhka	Leif H.	Karlsen	Addr3	IN3110	Program...	Beskr2...	5	C

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Anomalier: Sletting

StudentCourse

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peternl	Petter	Nilsen	Addr2	IN2090	Databaser	Beskr...	10	A
evgenit	Evgenij	Thorstensen	Addr1	IN2080	Beregn...	Descr...	10	A
leifhka	Leif H.	Karlsen	Addr3	IN2090	Databaser	Beskr...	10	B
leifhka	Leif H.	Karlsen	Addr3	IN3110	Program...	Beskr2...	5	C

- Deleting a course may delete a student
- Deleting a student may delete a course
- Difficult to fix this with this structure

Fix the structure

Username	Name	Surename	Address	Coursecode	Title	Desc.	Credits	Grade
evgenit	Evgenij	Thorstensen	Addr1	IN2090	Databaser	Beskr...	10	B
peternl	Petter	Nilsen	Addr2	IN2090	Databaser	Beskr...	10	A
evgenit	Evgenij	Thorstensen	Addr1	IN2080	Beregn...	Descr...	10	A
leifhka	Leif H.	Karlsen	Addr3	IN2090	Databaser	Beskr...	10	B
leifhka	Leif H.	Karlsen	Addr3	IN3110	Program...	Beskr2...	5	C

Fix the structure

Username	Name	Surename	Address	Coursecode	Title	Desc.	Credits	Grade
evgenit	Evgenij	Thorstensen	Addr1	IN2090	Databaser	Beskr...	10	B
peternl	Petter	Nilsen	Addr2	IN2090	Databaser	Beskr...	10	A
evgenit	Evgenij	Thorstensen	Addr1	IN2080	Beregn...	Descr...	10	A
leifhka	Leif H.	Karlsen	Addr3	IN2090	Databaser	Beskr...	10	B
leifhka	Leif H.	Karlsen	Addr3	IN3110	Program...	Beskr2...	5	C

Fix the structure

Username	Name	Surename	Address	Coursecode	Title	Desc.	Credits	Grade
evgenit	Evgenij	Thorstensen	Addr1	IN2090	Databaser	Beskr...	10	B
peternl	Petter	Nilsen	Addr2	IN2090	Databaser	Beskr...	10	A
evgenit	Evgenij	Thorstensen	Addr1	IN2080	Beregn...	Descr...	10	A
leifhka	Leif H.	Karlsen	Addr3	IN2090	Databaser	Beskr...	10	B
leifhka	Leif H.	Karlsen	Addr3	IN3110	Program...	Beskr2...	5	C

Student

Username	Name	Surename	Address
evgenit	Evgenij	Thorstensen	Addr1
peternl	Petter	Nilsen	Addr2
leifhka	Leif H.	Karlsen	Addr3

Fix the structure

Username	Name	Surename	Address	Coursecode	Title	Desc.	Credits	Grade
evgenit	Evgenij	Thorstensen	Addr1	IN2090	Databaser	Beskr...	10	B
peternl	Petter	Nilsen	Addr2	IN2090	Databaser	Beskr...	10	A
evgenit	Evgenij	Thorstensen	Addr1	IN2080	Beregn...	Descr...	10	A
leifhka	Leif H.	Karlsen	Addr3	IN2090	Databaser	Beskr...	10	B
leifhka	Leif H.	Karlsen	Addr3	IN3110	Program...	Beskr2...	5	C

Student

Username	Name	Surename	Address
evgenit	Evgenij	Thorstensen	Addr1
peternl	Petter	Nilsen	Addr2
leifhka	Leif H.	Karlsen	Addr3

Fix the structure

Username	Name	Surename	Address	Coursecode	Title	Desc.	Credits	Grade
evgenit	Evgenij	Thorstensen	Addr1	IN2090	Databaser	Beskr...	10	B
peternl	Petter	Nilsen	Addr2	IN2090	Databaser	Beskr...	10	A
evgenit	Evgenij	Thorstensen	Addr1	IN2080	Beregn...	Descr...	10	A
leifhka	Leif H.	Karlsen	Addr3	IN2090	Databaser	Beskr...	10	B
leifhka	Leif H.	Karlsen	Addr3	IN3110	Program...	Beskr2...	5	C

Student

Username	Name	Surename	Address
evgenit	Evgenij	Thorstensen	Addr1
peternl	Petter	Nilsen	Addr2
leifhka	Leif H.	Karlsen	Addr3

Fix the structure

Username	Name	Surename	Address	Coursecode	Title	Desc.	Credits	Grade
evgenit	Evgenij	Thorstensen	Addr1	IN2090	Databaser	Beskr...	10	B
peternl	Petter	Nilsen	Addr2	IN2090	Databaser	Beskr...	10	A
evgenit	Evgenij	Thorstensen	Addr1	IN2080	Beregn...	Descr...	10	A
leifhka	Leif H.	Karlsen	Addr3	IN2090	Databaser	Beskr...	10	B
leifhka	Leif H.	Karlsen	Addr3	IN3110	Program...	Beskr2...	5	C

Student

Username	Name	Surename	Address
evgenit	Evgenij	Thorstensen	Addr1
peternl	Petter	Nilsen	Addr2
leifhka	Leif H.	Karlsen	Addr3

Course

Coursecode	Title	Desc.	Credits
IN2090	Databaser	Beskr...	10
IN2080	Beregn...	Descr...	10
IN3110	Program...	Beskr2...	5

Fix the structure

Username	Name	Surename	Address	Coursecode	Title	Desc.	Credits	Grade
evgenit	Evgenij	Thorstensen	Addr1	IN2090	Databaser	Beskr...	10	B
peternl	Petter	Nilsen	Addr2	IN2090	Databaser	Beskr...	10	A
evgenit	Evgenij	Thorstensen	Addr1	IN2080	Beregn...	Descr...	10	A
leifhka	Leif H.	Karlsen	Addr3	IN2090	Databaser	Beskr...	10	B
leifhka	Leif H.	Karlsen	Addr3	IN3110	Program...	Beskr2...	5	C

Student

Username	Name	Surename	Address
evgenit	Evgenij	Thorstensen	Addr1
peternl	Petter	Nilsen	Addr2
leifhka	Leif H.	Karlsen	Addr3

Course

Coursecode	Title	Desc.	Credits
IN2090	Databaser	Beskr...	10
IN2080	Beregn...	Descr...	10
IN3110	Program...	Beskr2...	5

Fix the structure

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evgenit	Evgenij	Thorstensen	Addr1	IN2090	Databaser	Beskr...	10	B
peternl	Petter	Nilsen	Addr2	IN2090	Databaser	Beskr...	10	A
evgenit	Evgenij	Thorstensen	Addr1	IN2080	Beregn...	Descr...	10	A
leifhka	Leif H.	Karlsen	Addr3	IN2090	Databaser	Beskr...	10	B
leifhka	Leif H.	Karlsen	Addr3	IN3110	Program...	Beskr2...	5	C

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- Note: Same columns and same values!

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- Good structure: Separate table for students, courses and grades

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Rules of thumb for database design:

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- One table per relationship (grade)
- One table per multi-valued property (blood pressure)

Not only one good structure!

- There can be many good ways of structuring the same information into tables.
- The following two tables on *life expectancy in Norway* contain the same information without data duplication, but is structured differently:

Year	Men	Women
2017	80.9	84.3
2018	81.0	84.5
2019	81.2	84.7
2020	81.5	84.9

Gender	Year	LE
men	2017	80.9
men	2018	81.0
men	2019	81.2
men	2020	81.5
women	2017	84.3
women	2018	84.5
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women	2020	84.9

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- These columns are called *keys* or *primary keys*

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- E.g. in `BloodPressure` we used `PatientID` to reference a patient and in `Grade` we used `Username` to reference students and `Coursecode` to reference courses
- Such references are known as *foreign keys*

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- When organizations/companies merge, they need to merge their data
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- Data needs to be of the same format/same structure
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- Merging data into a common format/structure known as *data integration*

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- Standardization helps solving these problems

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