Repetition & Highlights: SQL from the outside and the inside



Details: https://www.postgresql.org/docs/current/sql.html

SQL is the inter-galactic data-speak with several sub-speaks:

- Data Definition Language (DDL):
 CREATE-queries, creating schemas
- Data Query Language (DQL):
 SELECT-queries
- Data Manipulation Language (DML): INSERT/UPDATE, write-queries
- Data Control Language (DCL):
 GRANT/REVOKE, access & user management



- Works on **tables** (or **relations**):
 - Relation name
 - **Relation schema**: set of attribute names with associated datatypes

Tutorials

- Set(!) of **Relation records:** tuples of elements conforming the schema
 - tutorial topic site ID w3schools SQL 2003STD Database 1 2 w3schools WebDev HTML 5 w3schools 3 CSS 3 WebDev SQL 2003STD w3resource Database 4 5 w3resource MySQL Database

• Example:



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o Example (Select-project-join-query)
SELECT p.Name, m.Id FROM Person p JOIN Movie m ON
p.Id = m.ActorId
WHERE m.DirectorId = 1234;

- Many other operators (LEFT/RIGHT JOIN, UNION, GROUP BY, subqueries, etc.)
- Always to keep in mind: bag semantics by default, NULL may behave strangely, rich datatypes



Inside SQL: Indices

- An **index** on an attribute A is a data structure that facilitates finding the elements with a certain values in A (A is called the **search key**)
- The index is **organised** (e.g., sorted) on the search key
- For each value of the search key, the index has a list of pointers to the corresponding records
- More than one index in the same table (or file) means
 - Faster search
 - More complexity (changes will lead to updated indices as well)
 - Increased storage requirement, larger files



a b c ...

Ζ

Primary indices Dense vs. Sparse

A **dense index** has one lookup for each value of the search key



A **sparse index** has one lookup for each data block





Types of indices

- Dense/sparse
- Primary/Cluster/Secondary
 - Determines how main data file is sorted
- Multilevel (indices on indices)
- Multidimentional (simultaneously on many attributes)
- Organisation: List/B+/hash/...



Inside SQL: The (Typical) Journey of a Query









- Each node in a parse tree is
- an *atom (primitive)* i.e., a lexical element like a keyword, name, constant, parentheses or operators ... (leaf node)
- a syntactic category part of the query ... (inner node)



Generating the logical query plan



Relational Algebra

• **Domain:** relations (tables)

• Set operators:

- 1. (Set) Union R U S
- 2. (Set) Difference R S
- 3. Projection $\pi_{L}(R)$
- 4. Selection $\sigma_{\rm C}({\rm R})$
- 5. Cartesian product (a.k.a., Cross Product and Cross Join) $R \times S$
- 6. Renaming $\rho_{S(A1,A2,...,An)}(R)$
- 7. + Expressible operators (Intersection, Other joins, Division, etc.)

• Bag operators:

- 1. Versions of the set operators
- 2. Duplicate elimination
- 3. Aggregation
- Beyond bags:
 - 1. Sorting, etc.



From SQL to RA Example



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Query Optimisation: Equivalent Plans

- There are many RA lows preserving query results:
 - $_{\odot}\,$ associativity and commutativity of many operators
 - \circ e.g., R⋈S = S⋈R, R⋈ (S ⋈ T) = (R ⋈ S) ⋈ T, etc.
 - $_{\odot}\,$ special rules for selection

$$\circ$$
 e.g., $\sigma_{a}(R \bowtie S) = \sigma_{a}(R) \bowtie \sigma_{a}(S)$, $\sigma_{a AND b}(R) = \sigma_{a}(\sigma_{b}(R))$
 \circ etc.

From one LQP we can get many equivalent ones



Query Plan Cost Estimation



Estimating LQPs Costs

• To assess LQPs, we need some way to calculate cost

- DBMS estimates the costs
- We use **size of temporary relations** (#tuples x TupleSize)
 - Estimate the result of each operator
 - Add costs to the tree
 - The cost of the plan is the sum of all the costs in the tree (except the leaves and the root)



Comparing Logical Query Plans - Example





Physical Plans



Physical operators

- A query can be expressed in terms of a relational algebra expression
 - A physical query plan (PQP) is implemented by algorithms for relational algebra operators
 - They are based on basic algorithms for reading (scanning) a relation, sorting a relation, etc.
 - Two-Phase Multiway Merge Sort (TPMMS), etc.
- To choose a good PQP, we need to estimate the cost of algorithms
 - We can use the number of disk IOs as the cost



Physical Plans



Repetition & Highlights: Graph DBs



(Labelled) graph databases

- For data that is natural to describe and traverse as graphs
 - Each node has an inner structure describing its properties
 - The edges indicate relationships between the nodes
 - The edges can carry information in the same way as the nodes
- <u>Can be schema-free</u>
 - New nodes and edges (with new inner structures) can be introduced dynamically
 - Existing nodes and edges can be expanded with new properties
- Search: Specify how the graph should be navigated
 - The graph is traversed directly via pointers to neighboring nodes (the traverse requires no indexes and no join operations)



Linked Data, Semantic Knowledge Graphs

- Method for publishing data on the Web
- Self-describing data and relations
- Interlinking
- Accessed using semantic queries
- A set of open standards developed by W3C
 - Data format: RDF
 - Knowledge representation: RDFS/OWL
 - Query language: SPARQL









Labeled Property Graphs (e.g., Neo4j)



- Node, Relationship, Property, Label, Relationship type
- Cypher query language

