IN3060/4060 – Semantic Technologies – Spring 2021

Lecture 11: OWL: Loose Ends

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Mandatory exercises

- Oblig 6 published after lecture.
- First attempt by April 16th.
- Second attempt two weeks after feedback.

Outline

Make it simple!

- “Data level” with resources
- “Ontology level” with properties and “classes”
- Can have `rdf:type` relation between data objects and classes
- Allow a fixed vocabulary for relations between classes and properties
- Interpret:
  - Class as set of data objects
  - Property as relation between data objects

OWL 2 TBox and ABox

- The TBox
  - is for *terminological knowledge*
  - is independent of any actual instance data
  - is a set of axioms:
    - Class inclusion $\sqsubseteq$, equivalence $\equiv$
    - roles symmetric, asymmetric, reflexive, irreflexive, transitive,. . .
    - roles functional, inverse functional
    - inverse roles: $\mathit{hasParent} \equiv \mathit{hasChild}^{-1}$
    - role inclusion $\mathit{hasBrother} \sqsubseteq \mathit{hasSibling}$
    - role chains $\mathit{hasParent} \circ \mathit{hasBrother} \sqsubseteq \mathit{hasUncle}$
  - Only certain combinations allowed

OWL 2 TBox and ABox

- The ABox
  - is for *assertional knowledge*
  - contains facts about concrete instances $a, b, c, \ldots$
  - A set of (negative) concept assertions $C(a)$, $\neg D(b)$ ...
  - and (negative) role assertions $R(b, c)$, $\neg S(a, b)$
  - also `owl:sameAs`: $a = b$ and `owl:differentFrom`: $a \neq b$.

A Strange Catalogue

- We have seen many nice things that can be said in OWL
- Why the strange restrictions, e.g. on role axioms?
- Why not use 1st-order logic, could say much more?

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- Because of the reasoning
  - Class satisfiability ($C \not\equiv \bot$)
  - Classification ($C \sqsubseteq D$)
  - Instance Check ($C(a)$)
  - ...
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- Add a little more to OWL, and this is lost

Outline

1. Reminder: OWL
2. Disjointness and Covering Axioms
3. Keys
4. Punning
5. More about Datatypes
6. What can't be expressed in OWL 2
7. OWL 2 profiles

Single and Married

- Try to model the relationship between the concepts *Person*, *Married* and *Single*:
- First try:

$$\begin{aligned} \mathit{Single} &\sqsubseteq \mathit{Person} \\ \mathit{Married} &\sqsubseteq \mathit{Person} \end{aligned}$$

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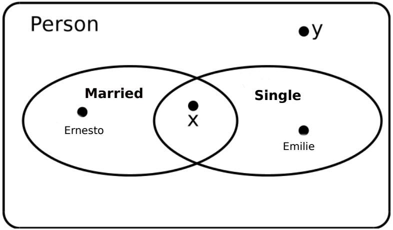
- General shape of a model:



- x is both *Single* and *Married*, y is neither but a *Person*.

Disjointness Axioms

- Nothing should be both a *Single* and a *Married*
- Add a *disjointness* axiom for *Single* and *Married*
- Equivalent possibilities:

$$\begin{aligned} \mathit{Single} \sqcap \mathit{Married} &\equiv \bot \\ \mathit{Single} &\sqsubseteq \neg \mathit{Married} \\ \mathit{Married} &\sqsubseteq \neg \mathit{Single} \end{aligned}$$

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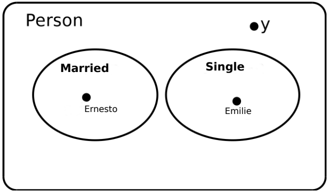

- Specific support in OWL (`owl:disjointWith`) and Protégé

Covering Axioms

- Any *Person* should be either *Single* or *Married*.
- Add a *covering axiom* $Person \sqsubseteq Married \sqcup Single$

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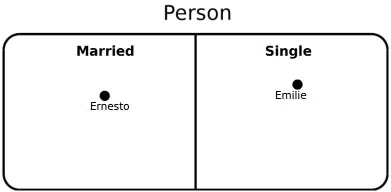

- Specific support in Protégé (Edit Menu: "Add Covering Axiom")

Meat and Veggies

- Careful: not all subclasses are disjoint and covering
- Subclasses can be covering but not disjoint.
- E.g.

$$\begin{array}{rcl} \mathit{MeatEatingMammal} & \sqsubseteq & \mathit{Mammal} \\ \mathit{VeggieEatingMammal} & \sqsubseteq & \mathit{Mammal} \end{array}$$

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- $\mathit{Mammal} \sqsubseteq \mathit{MeatEatingMammal} \sqcup \mathit{VeggieEatingMammal}$
- But there are mammals eating both
- No disjointness axiom for *MeatEatingMammal* and *VeggieEatingMammal*

Cats and Dogs

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- Nothing is both a cat and a dog: $Cat \sqsubseteq \neg Dog$
- But there are mammals which are neither
- No covering axiom with subclasses *Cat* and *Dog* for *Mammal*

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- Subclasses can be neither disjoint nor covering.
- E.g.

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- A course is uniquely determined by code, semester, year.
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- hasKey: if two named instances of the class coincide on values for each of key properties, then these two individuals are the same.

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OWL 2 Keys

- OWL 2 includes special “hasKey” axioms
- Example: `Course hasKey {hasCode, hasSemester, hasYear}`
- Works for object properties and datatype properties.
- OWL Keys apply only to explicitly **named instances**
  - Makes reasoning tractable.
  - It may not be supported by all OWL 2 reasoners

Reasoning with OWL Keys

- Given:
  - `:Norwegian hasKey {:personnr}`
  - `:david a :Norwegian`
  - `:david :personnr "12345698765"`
  - `:davidC a :Norwegian`
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- Given:
  - `:Singleton hasKey {:id}`
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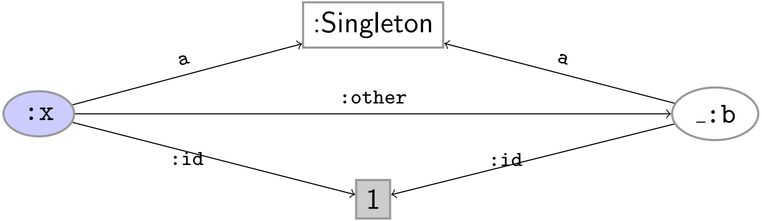

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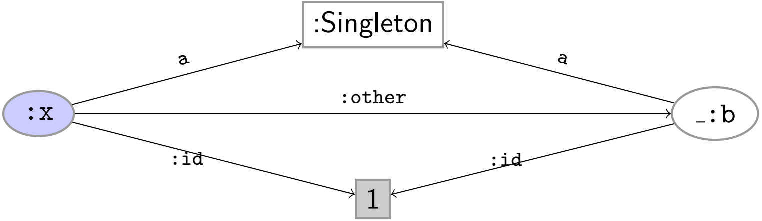

- Since `_:b` is a blank node, and therefore not an explicitly named instance,
- the reasoner does not infer `:x owl:sameAs _:b`.

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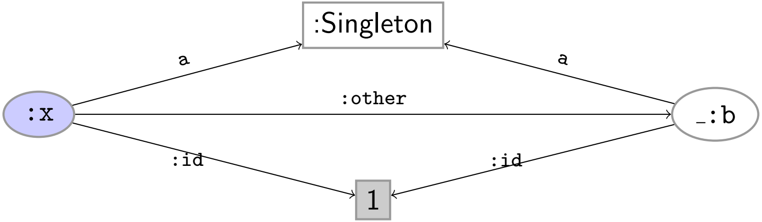

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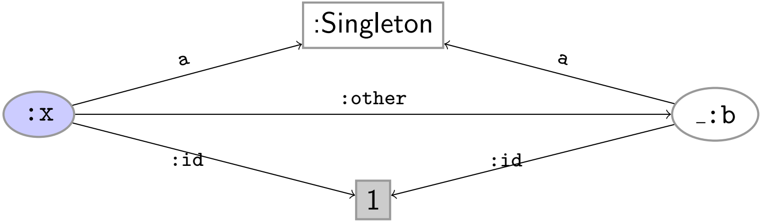

- This is *not* inconsistent.
- Distinct keys only required for explicitly named individuals.

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(1):Service     rdf:type    owl:Class .
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- Remember: In OWL strict separation of classes, properties and individuals. However, not entirely correct...
- OWL 2 introduces *punning*, allowing one URI to be used for, e.g., both a class and an individual, both an individual and property
- Restriction: not both a class and a datatype property, or for different property types.
- Example:

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(1):Joe       rdf:type     :Eagle .
(2):Eagle     rdf:type     :Species .
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`:Eagle` is both a class and an individual.

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  `:Eagle` is both a class and an individual.
- However, semantically, "punned" URI are treated as different terms. (under the hood)
  - Meaning, the class `:Eagle` is different from the individual `:Eagle`.
  - Axioms about the class is not transferred to the individual, or vice versa.

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- Why?
  - This says that Scotch has at least 3 *different ages*
  - For instance -1, 0, 15

A possible solution

- Idea: don't use age.
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    - range *int*
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e.g. `:young :casked "2000"^^int, "2001"^^int, "2002"^^int`

- $\mathit{Scotch} \sqsubseteq \mathit{Whisky} \sqcap \ \geq_3 \mathit{casked}.\mathit{int}$

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- Can't express e.g. that the years are consecutive
  - Knowing a whisky is casked in 2000 and 2009 doesn't imply it is casked for 10 years.
- Reasoning about $\geq_n$ often works by generating n sample instances
  - $Town \equiv \geq_{10000} inhabitant.Person$
  - $Metropolis \equiv \geq_{1000000} inhabitant.Person$
  - Will kill almost any reasoner

Reminder: Datatype properties

- OWL distinguishes between
  - object properties: go from resources to resources
  - datatype properties: go from resources to literals
- OWL (2) prescribes a list of available built-in datatypes for literals
  - Numbers: real, rational, integer, positive integer, double, long,. . .
  - Strings
  - Booleans
  - Binary data
  - IRIs
  - Time Instants
  - XML Literals
- Varying tool support (e.g., depending on editor and reasoner)

Data Ranges

- Like concept descriptions, only for data types
- Boolean combinations allowed (Manchester syntax)
  - `xsd:integer or xsd:string`
  - `xsd:integer and not xsd:byte`

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  - `xsd:integer or xsd:string`
  - `xsd:integer and not xsd:byte`
- Each basic datatype can be restricted by a number of *facets*
  - `xsd:integer[`$\geq$ `9]` – integers ≥ 9.
  - `xsd:integer[`$\geq$ `9,` $\leq$ `11]` – integers between 9 and 11.
  - `xsd:string[length 5]` – strings of length 5.
  - `xsd:string[maxLength 5]` – strings of length ≤ 5.
  - `xsd:string[minLength 5]` – strings of length ≥ 5.
  - `xsd:string[pattern "[01]*"]` – strings consisting of 0 and 1.

Range Examples

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  `Whisky and age some integer[>= 12]`

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- Note: often makes best sense with functional properties
  Why?

Pattern Examples

- An integer or a string of digits
  - `xsd:integer or xsd:string[pattern "[0-9]+"]`
- ISBN numbers: 13 digits in 5 “–”-separated groups, first 978 or 979, last a single digit.
  - Book $\sqsubseteq$ `ISBN some string[length 17 ,`

    `pattern "97[89]-[0-9]+-[0-9]+-[0-9]+-[0-9]"]`

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  - `Book` $\sqsubseteq$ `ISBN some string[length 17 , pattern "97[89]-[0-9]+-[0-9]+-[0-9]+-[0-9]"]`
- Reasoning about patterns:
  - `R` a functional datatype property
  - $A \equiv$ `R some string[pattern "(ab)*"]`
  - $B \equiv$ `R some string[pattern "a(ba)*b"]`
  - Reasoner can find out that $B \sqsubseteq A$.

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Expressivity

- Certain *relationships* between concepts and properties can't be expressed in OWL
- E.g.
  - Given that property *hasSibling* and class *Male* are defined...
  - ...cannot say that $hasBrother(x, y)$ iff $hasSibling(x, y)$ and $Male(y)$.
- Usually, adding such missing relationships would lead to undecidability
- *Not* easy to show that something is not expressible
  - We look at some examples, not proofs

Brothers

- Given terms

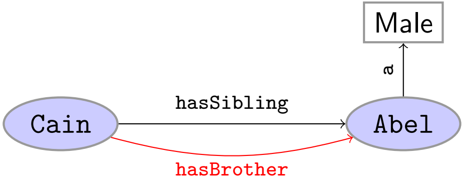
  hasSibling *Male*

Brothers

- Given terms

hasSibling *Male*

- . . . a brother is *defined* to be a sibling who is male



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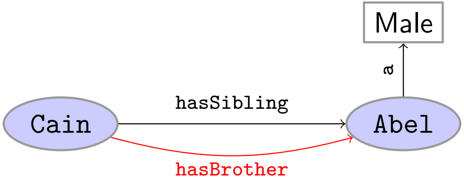

- Best try:

$$\begin{aligned} &hasBrother \sqsubseteq hasSibling \\ &\top \sqsubseteq \forall hasBrother.Male \qquad \text{or: } \mathrm{rg}(hasBrother, Male) \\ &\exists hasSibling.Male \sqsubseteq \exists hasBrother.\top \end{aligned}$$

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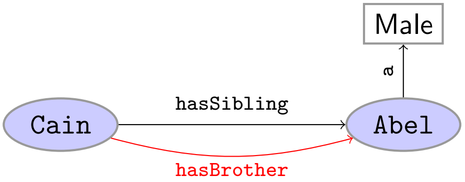

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- Not enough to infer that one's male sibling is one's brother.

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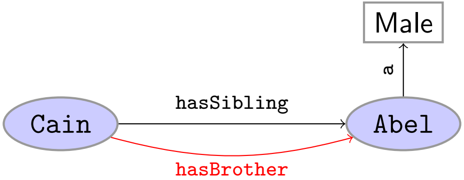

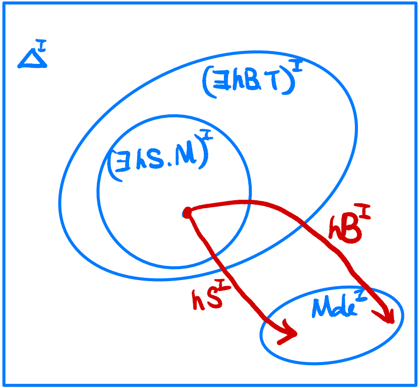

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- Given terms

hasParent *hasBrother*

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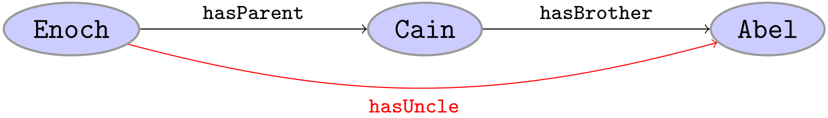

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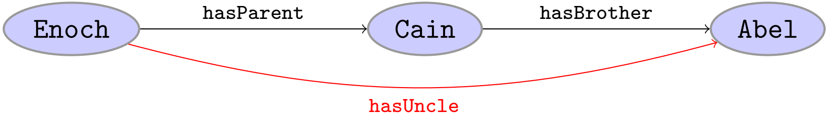

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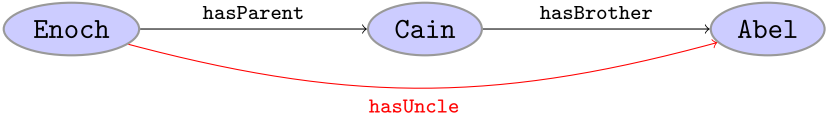

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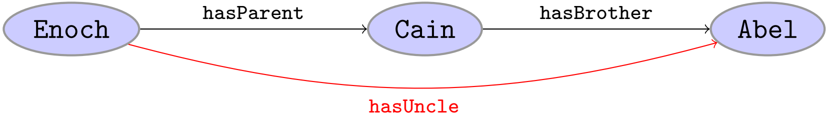

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- properties cannot be declared sub-properties of property chains in OWL 2.
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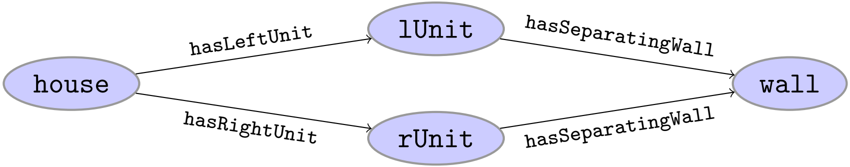

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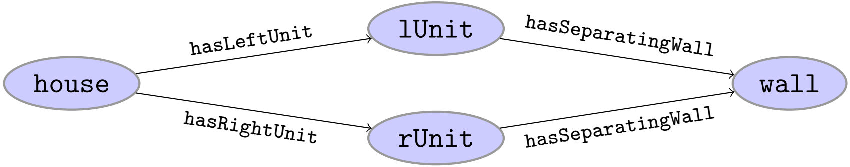

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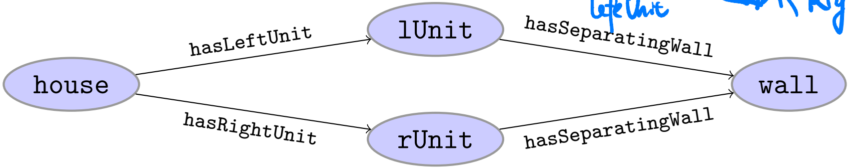

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$$\mathit{TwinParent} \equiv \mathit{Person} \sqcap \geq_2 \mathit{hasChild}.\exists \mathit{hasBirthday}[\ldots]$$

- Still no way of connecting the birthdays

Reasoning about Numbers

- Reasoning about natural numbers is undecidable in general.
- DL Reasoning is decidable
- Therefore, general reasoning about numbers can’t be “encoded” in DL
- Cannot encode addition, multiplication, etc.
- Note: a lot can be done with other logics, but not with DLs
    - Outside the intended scope of Description Logics

Combining OWL 2 and Rules

Some limitation may be addressed

- SWRL: Semantic Web Rule Language
- Uses XML syntax based on RuleML
- OWL 2 + unrestricted SWRL leads to undecidability
- Restricted SWRL + OWL is decidable and very powerful
- A bit more in the next SPARQL lesson

Outline

1. Reminder: OWL
2. Disjointness and Covering Axioms
3. Keys
4. Punning
5. More about Datatypes
6. What can’t be expressed in OWL 2
7. **OWL 2 profiles**

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- OWL Full: Anything goes: classes, relations, individuals, ... like in RDFS, are not kept apart. Highly expressive, not decidable. But we want OWL’s reasoning capabilities, so stay away if you can—and you almost always can.

OWL 2 Validator: `http://mowl-power.cs.man.ac.uk:8080/validator/`

OWL EL

Based on DL $\mathcal{EL}^{++}$.

$\mathcal{EL}^{++}$ concept descriptions, simplified

$$
\begin{array}{rlcl}
C, D \rightarrow & A & \mid & \text{(atomic concept)} \\
& \top & \mid & \text{(universal concept)} \\
& \bot & \mid & \text{(bottom concept)} \\
& \{a\} & \mid & \text{(}\textit{singular}\text{ enumeration)} \\
& C \sqcap D & \mid & \text{(intersection)} \\
& \exists R.C & \mid & \text{(existential restriction)}
\end{array}
$$

Axioms

- $C \sqsubseteq D$ and $C \equiv D$ for concept descriptions D and C.
- $R \sqsubseteq S$, $R \equiv S$, $R \circ S \sqsubseteq R$, $dom(R) \sqsubseteq C$ and $ran(R) \sqsubseteq D$ for concept descriptions D, C and roles R, S.
- $C(a)$ and $R(a, b)$ for concept C, role R and individuals a, b.

OWL EL contd.

Not supported, simplified:

- negation, (NB, disjointness of classes: $C \sqcap D \sqsubseteq \bot$ possible),
- disjunction,
- universal quantification,
- cardinalities,
- inverse roles,
- plus some role characteristics.
- reduced list of datatypes (e.g., not supported "boolean" nor "double")

Complete list: `http://www.w3.org/TR/owl2-profiles/#Feature_Overview`.

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- Checking ontology consistency, class expression subsumption, and instance checking is in **P**.
- “Good for large ontologies.”
- Used in many biomedical ontologies (e.g. SNOMED CT).

OWL QL

Based on DL-Lite$_R$.

DL-Lite$_R$ concept descriptions

$$\begin{array}{lcl} B & := & A \mid \exists R.\top \\ C & := & B \mid \neg B \end{array}$$

DL-Lite$_R$ role descriptions

$$\begin{array}{lcl} Q & := & R \mid R^{-} \\ S & := & Q \mid \neg Q \end{array}$$

DL-Lite$_R$ Axioms

- Concept inclusions $B \sqsubseteq C$ for concept descriptions B and C.
- Role inclusions $Q \sqsubseteq S$ for roles Q, S.
- $A(a)$ and $R(a, b)$ for atomic concept A, role R and individuals a, b.

OWL QL contd.

Not supported, simplified:

- disjunction,
- universal quantification,
- cardinalities,
- functional roles, keys,
- $=$ (`SameIndividual`)
- enumerations (closed classes),
- subproperties of chains, transitivity
- reduced list of datatypes (e.g., not supported "boolean" nor "double")

Complete list: `http://www.w3.org/TR/owl2-profiles/#Feature_Overview_2`.

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Complete list: `http://www.w3.org/TR/owl2-profiles/#Feature_Overview_2`.

- Captures language for which queries can be translated to SQL.
- "Good for large datasets."

OWL RL

- Puts constraints in the way in which constructs are used (i.e., syntactic subset of OWL 2).
- So that OWL 2 RL axioms can be directly translated into datalog rules.
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- It is closely related to Description Logic Programs (DLP).
- Syntactic restriction on Class Expressions in OWL 2 RL can be found: `http://www.w3.org/TR/owl2-profiles/#Feature_Overview_3`.
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- Reasoning in RL is possible in polynomial time, but in other cases results may be incomplete.

Next

- 9 April: SPARQL 1.1
- 16 April: RDF Validation
- 23 April: Application in Norway (Aibel, DNV)
- 30 April, 7 May: OTTR Templates (T.B.A.)
- 14 May: Open RDF Data