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Deadline: 18 October 2023, 23:59

### Exercise O2.1

(Validity and Proof Calculi)

Consider the following formulae.

$$F_1: ((p \rightarrow q) \wedge (q \rightarrow r)) \rightarrow (p \rightarrow r) \quad (\text{“Transitivity”})$$

$$F_2: \forall x \exists y (p(x) \wedge (p(y) \rightarrow q(x))) \rightarrow \forall z q(z) \quad (\text{“A modus ponens”})$$

Prove the validity of formulae  $F_1$  and  $F_2$  in the *resolution calculus*. First, translate the negated formula into clausal form.

### Exercise O2.2

(Adding a Logical Operator)

The logical operator  $\uparrow$  is defined as follows:  $A \uparrow B \equiv \neg(A \wedge B)$ .

Resolution works on formulae in clause form. For a resolution-based theorem proving programme to work with full 1st-order or propositional formulae, these are transformed to clause form before starting resolution. What would have to be changed in such a theorem prover to make it accept formulae with a  $\uparrow$  operator?