

Exercise Session 10

Towards Home Exam 2

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IN[59]570: Distributed Objects

April 29, 2019 (Revised on May 6, 2019)

The source code for these slides is maintained here:

<https://github.com/emerald/in5570v19/tree/master/exercise-sessions/10>

Agenda

1. Parametric Polymorphism (in Emerald)
2. The PCRTType from Home Exam 2
3. Emulating Unavailability (using Docker)

Polymorphism in Emerald

In a language with **polymorphism**, we can write code that works for many datatypes, not just one particular type of data

- ▶ Emerald supports **inclusion polymorphism**, due to its *conformance relation*: In place of a particular type, Emerald happily accepts a value of a different, but conforming type
 - ▶ You know about conformance from previous weeks
- ▶ Emerald also supports **parametric polymorphism**, where types depend on the actual parameters (e.g., the parameters to a method call)
 - ▶ See also Section 7.4 on pages 18-19 of [Raj et al., 1991]

Passing the Type Parameter Implicitly

Showing the contents of `implicit.m`

```
const main <- object main
  op showType[a : t]                % Where does t come from?
    forall t
      stdout.putstring[t$name || "\n"]
  end showType
  initially
    self.showType[5]                % From here!
  end initially
end main
```

```
$ emx implicit.x
integertype
```

- ▶ The `forall` clause introduces an *unconstrained type variable*
- ▶ We can then (somewhat backwards) use `t` in the signature preceding the `forall` clause (`op showType[a : t]`)
- ▶ We must use a `forall` clause, as otherwise `t` is undefined ☹
- ▶ `t` gets the type `ConcreteType`; `t` can be inspected at runtime

There Are No Runtime Costs

All types are determined at compile-time!

- ▶ Watch out for “type must be manifest” errors from the Emerald compiler; if you get these, it means that the type of some expression cannot be determined at compile-time

Passing the Type Parameter Explicitly

Showing the contents of `explicit.m`

```
const main <- object main
  op showType[t : Type, a : t]      % t is an explicit parameter
    stdout.putstring[(typeof a)$name || "\n"]
  end showType
  initially
    self.showType[Integer, 5]      % but still comes from here
  end initially
end main
```

```
$ emx explicit.x
integertype
```

- ▶ Recall that, in Emerald, types are also objects
- ▶ As another example, recall how you must explicitly pass a type to the `Array.of` constructor, to get an `Array` of that type
- ▶ Unfortunately, we can't do much with `t` directly (see line 3)
- ▶ Values of type `Type` are assumed to only be available at runtime

A Brief Interlude: `typeof` VS. `syntactictypeof`

Quite unlike many popular languages, Emerald provides two ways to ask for the type of an expression — `typeof` and `syntactictypeof`:

- ▶ `typeof` gives the actual type at *runtime*
- ▶ `syntactictypeof` gives the type determined at *compile time*

The Emerald system guarantees that the runtime type of an expression will conform to its compile-time type.

typeof VS. syntactictypeof Illustrated

What happens if we ask for `typeof t` instead of `typeof a` above?

```
$ diff explicit.m typeof.m
3c3
<     stdout.putstring[(typeof a)$name || "\n"]
---
>     stdout.putstring[(typeof t)$name || "\n"]
$ emx typeof.x
pat
```

What about `syntactictypeof t`?

```
$ diff explicit.m syntactictypeof.m
3c3
<     stdout.putstring[(typeof a)$name || "\n"]
---
>     stdout.putstring[(syntactictypeof t)$name || "\n"]
$ emx syntactictypeof.x
typetype
```


Constraining Type Variables Such That ...

Showing lines 1–8 of replicate.m

```
const RType <- typeobject RType
  operation replicate[X : t, N : Integer]
    forall t
      suchthat
        t *> typeobject ot
          op clone -> [result : t]
        end ot
      end RType
```

- ▶ Use a **suchthat** clause
- ▶ *> means *conforms to*, and the expression on the right-hand side can be any type-valued expression

Building Values with Types Such That ...

Showing lines 10–21 of replicate.m

```
const Replicator : RType <- object Replicator
  export operation replicate[X : t, N : Integer]
    forall t
      suchthat
        t *> typeobject ot
          op clone -> [result : t]
        end ot
      for i : Integer <- 0 while i < N by i <- i + 1
        const c <- X.clone[]
      end for
    end replicate
end Replicator
```

- ▶ t has to be available at compile-time
- ▶ Unfortunately, the only way to do so is with a **forall** clause
- ▶ Parametric polymorphism can be quite verbose in Emerald ☹

Relicating Integers and Strings

Showing lines 23–40 of replicate.m

```
const RInt <- class RInt[value : Integer]  
  export operation clone -> [result : RIntType]  
    stdout.putstring["Cloning " || value.asstring || "..\n"]  
    result <- RInt.create[value]  
  end clone  
end RInt  
const RString <- class RString[value : String]  
  export operation clone -> [result : RStringType]  
    stdout.putstring["Cloning " || value || "..\n"]  
    result <- RString.create[value]  
  end clone  
end RString  
const main <- object main  
  initially  
    Replicator.replicate[RInt.create[5], 3]  
    Replicator.replicate[RString.create["Hello"], 5]  
  end initially  
end main
```

Constructing Dependent Types

Showing the contents of `replicas.m`

```
const RaType <- typeobject RaType
  operation replicas[X : t] -> [Array.of[rt]]
  forall t
  where
    rt <- typeobject rt
      operation read -> [o : t]
      operation write[o : t]
    end rt
end RaType
```

- ▶ Use a **where** clause
- ▶ The type `rt` *depends on* the given type `t`
- ▶ Constructing a value of this particular type however, is even more tricky than for `RType`; that is, without resorting to type assertions (**view ... as ...**)

The PCRTType in Home Exam 2

Showing the contents of typedefs.m

```
const PCRTType <- typeobject PCRTType
  operation replicate[X : t, N : Integer]
  forall t
  suchthat
    t *> typeobject ot
    op clone -> [result : t]
    end ot
  operation replicas[X : t] -> [Array.of[rt]]
  forall t
  where
    rt <- typeobject rt
    operation read -> [o : t]
    operation write[o : t]
    end rt
end PCRTType
```

Emulating Unavailability (using Docker)

- ▶ Docker containers connect to the web via a network “bridge”
- ▶ You can connect and disconnect containers from such a bridge
- ▶ If a container is not connected to a network bridge, for all intents and purposes, it is offline
- ▶ This way, we can simulate temporary node unavailability

Creating A (New) Network Bridge

Although a Docker container is by default connected to a default network bridge, you can exert grander control by creating your own network bridge

- ▶ To create a network bridge:

```
$ docker network create \  
  --subnet=172.18.0.0/24 \  
  --ip-range=172.18.0.0/24 \  
  --driver=bridge \  
  unavail
```

- ▶ The subnet and IP range arguments effectively make the following IP address available for containers to use:
 - ▶ 172.18.0.2
 - ▶ 172.18.0.3
 - ▶ ...
 - ▶ 172.18.0.254
- ▶ This bridge is named `unavail` (see last argument)

Connecting Running Containers to the Bridge

- ▶ Start a Docker container
 - ▶ Let it have the container ID 85a87446465
- ▶ To connect 85a87446465 to unavail at address 172.18.0.2:

```
$ docker network connect --ip=172.18.0.2 unavail 85a87446465
```

- ▶ To disconnect 85a87446465 from unavail:

```
$ docker network disconnect unavail 85a87446465
```

In a similar vein, you can connect up a range of containers, and methodically take them offline one-by-one.

- ▶ See attached `monitor.m` for a sample program that monitors the list of available nodes

More Network Operations

As you experiment with Docker and bridge networks, you might find the following useful:

- ▶ To inspect the state `unavail` (e.g., see list of connected containers):

```
$ docker network inspect unavail
```

- ▶ To remove `unavail`

```
$ docker network rm unavail
```

Further Reading



Raj, Tempero, Levy, Black, Hutchinson, and Jul (1991),
Technical Report: The Emerald Programming Language

<https://www.uio.no/studier/emner/matnat/ifi/INF5510/v15/pensum/Report.pdf>