



INF3110 – Programming Languages Runtime Organization part II

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Slides adapted from previous years' slides
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Higher-Order Functions

- Language features
 - Functions passed as arguments
 - Functions that return functions from nested blocks
 - Need to maintain environment of function
- Simpler case
 - Function passed as argument
 - Need pointer to activation record “higher up” in stack
- More complicated second case
 - Function returned as result of function call
 - Need to keep activation record of returning function

Functions as parameters: why?

Given a Person class and a list of Persons

Can you write a function that finds all persons that

- Are female?
- Are older than 50 years?
- Like drinking beer?

Functions as parameters: why?

Given a Person class and a list of Persons

Can you write a function that finds all persons that

- Are female?
- Are older than 50 years?
- Like drinking beer?

A first attempt:

```
List<Person> findPersonsThatAreFemale(List<Person> persons) {  
    List<Person> filteredList = new List<Person>();  
    for(Person p in persons) {  
        if(p.gender == "female") {  
            filteredList.Add(p);  
        }  
    }  
    return filteredList;  
}
```

Functions as parameters: why?

Given a Person class and a list of Persons

Can you write a function that finds all persons that

- Are female?
- Are older than 50 years?
- Like drinking beer?
- ...

A first attempt:

```
List<Person> findPersonsThatAreOlderThan50(List<Person> persons) {  
    List<Person> filteredList = new List<Person>();  
    for(Person p in persons) {  
        if(p.age > 50) {  
            filteredList.Add(p);  
        }  
    }  
    return filteredList;  
}
```

Why functions as parameters? – DRY!

```
List<Person> filterPersons(List<Person> persons, (Person → Boolean) filter) {  
    List<Person> filteredList = new List<Person>();  
    for(Person p in persons) {  
        if(filter(p)) {  
            filteredList.Add(p);  
        }  
    }  
    return filteredList;  
}
```

Imaginary func syntax



```
filterPersons( persons, function(Person p) { return p.age > 50 } );
```

```
filterPersons( persons, function(Person p) { return p.gender == "female" } );
```

Is this in use in languages today?

Traditional OO approach: make a class out of it!

E.g. in Java (pre v8):

```
Collections.sort(list, new Comparator<MyClass>() {  
    public int compare(MyClass a, MyClass b)  
    {  
        // compare objects here  
    }  
});
```

What is going on here?

- `Comparator<T>` is an interface, and `T` is a type parameter
- This interface has one method signature, `int compare(T a, T b)`;
- Starting from the first `{`, we have an anonymous class implementing this interface

Is this in use in languages today?

Functional approach:

- **Java (v8 and up):**
`list.sort((a, b) -> a.isGreaterThan(b));`
- **C#:**
`myList.Where(a => a.Number > 42).OrderBy(a => a.Number)
 .ThenBy(a => a.FooBar);`
- **Python:**
`Celsius = [39.2, 36.5, 37.3, 37.8]
Fahrenheit = map(lambda x: (float(9)/5)*x + 32, Celsius)`
- **JavaScript (Node):**
`app.get('/somepath/:date', function (req, res)
 res.setHeader('Content-Type', 'application/json');
 fetchStuff({ date: req.params.date}, function (error, result) {
 if (error) console.log(error);
 res.end(result);
 });
});`

Pass function as argument

There are two declarations of x

Which one is in scope for each usage of x?

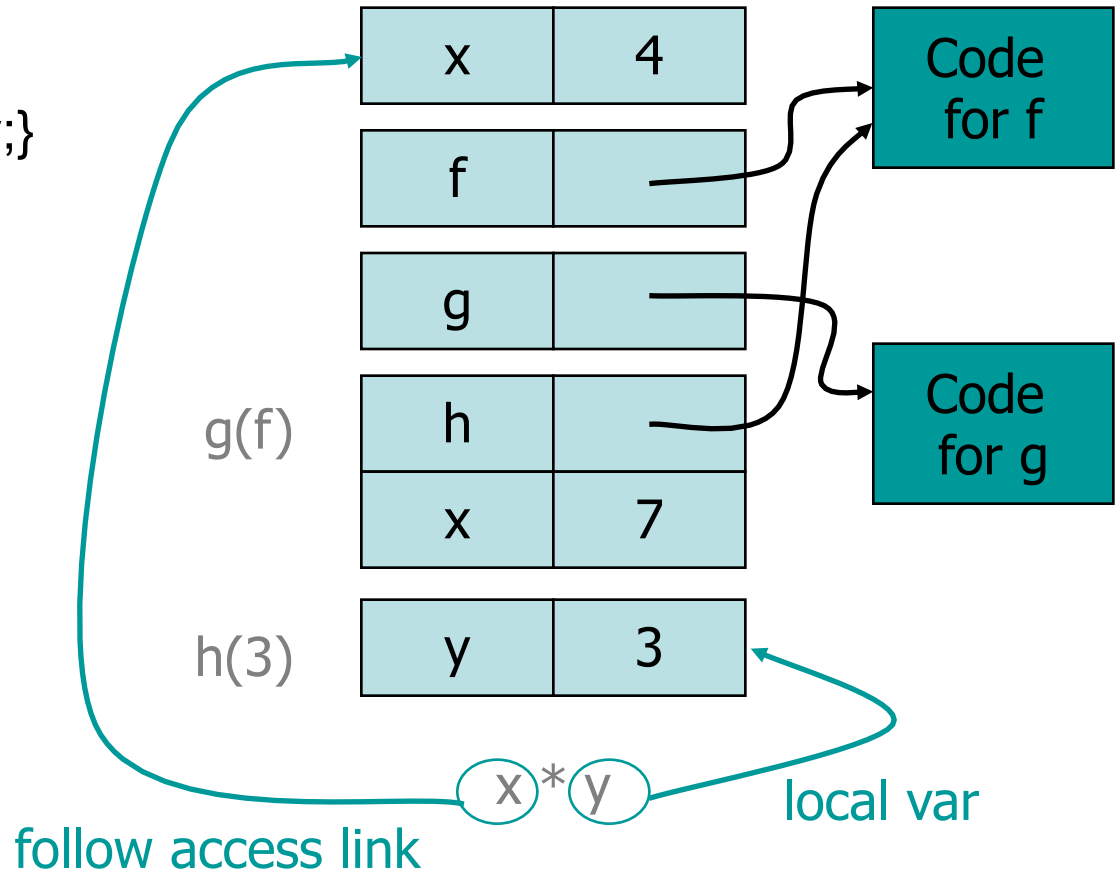
```
{ int x = 4;  
  { int f(int y) {return x*y;}  
    { int g(int→int h) {  
      int x=7;  
      return h(3) + x;  
    }  
    g(f);  
  }  
}
```

Formal function parameter

Actual function parameter

Static Scope for Function Argument

```
{ int x = 4;  
  { int f(int y) {return x*y;}  
    { int g(int→int h) {  
      int x=7;  
      return h(3) + x;  
    }  
  }  
}  
g(f);  
}  
}
```



How is access link for `h(3)` set? → Next slides

Closures

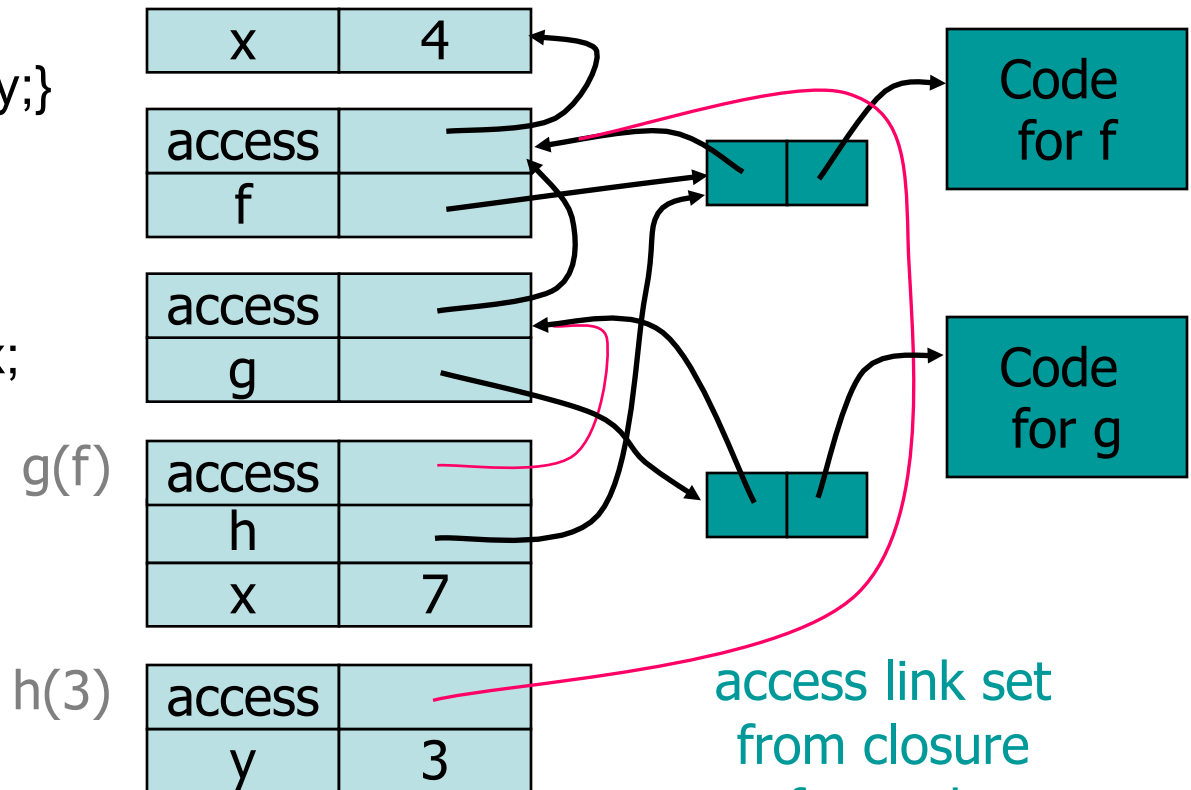
- Function value is pair *closure* = $\langle env, code \rangle$
- When a function represented by a closure is called
 - Allocate activation record for call (as always)
 - Set the access link in the activation record using the environment pointer from the closure

Function Argument and Closures

Run-time stack with access links

```

{ int x = 4;
  { int f(int y){return x*y;}
    { int g(int→int h) {
      int x=7;
      return h(3)+x;
    }
  }
}
g(f);
  
```



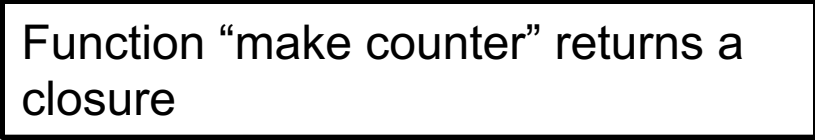
access link set from closure for each function call

Return Function as Result

- Language feature
 - Functions that return “new” functions
 - Need to maintain environment of function
- Function “created” dynamically
 - function value is closure = $\langle \text{env}, \text{code} \rangle$
 - code *not* compiled dynamically (in most languages)

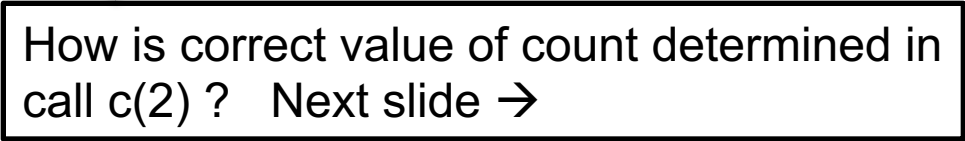
Example: Return function with private state

Function “make counter” returns a closure



```
{ int→int mk_counter (int init) {  
    int count = init;  
    int counter(int inc)  
        { return count += inc;}  
    return counter  
}  
int→int c = mk_counter(1);  
print c(2) + c(2);  
}
```

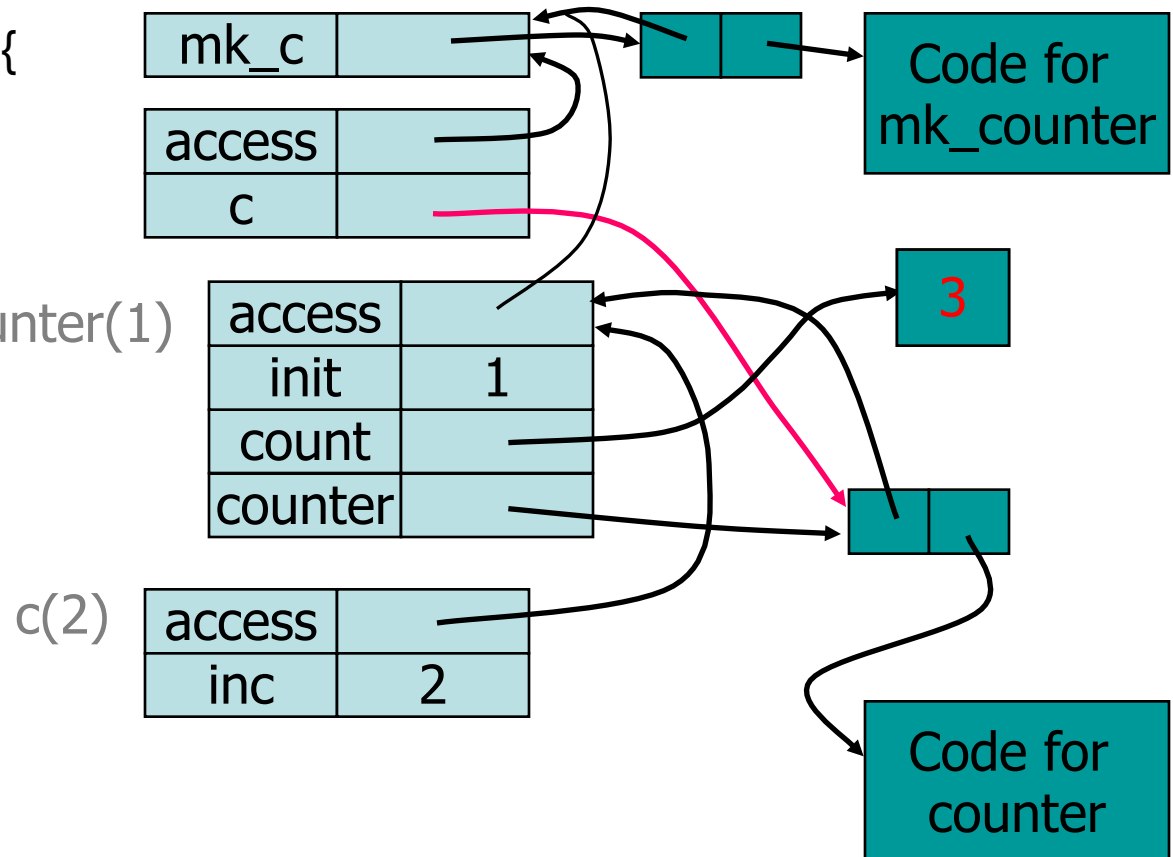
How is correct value of count determined in call c(2)? Next slide →



Function Results and Closures

```

{int→int mk_counter (int init) {
  int count = init;
  int counter(int inc)
    { return count+=inc;}
  return counter  mk_counter(1)
}
int→int c = mk_counter(1);
print c(2) + c(2);
}
    
```



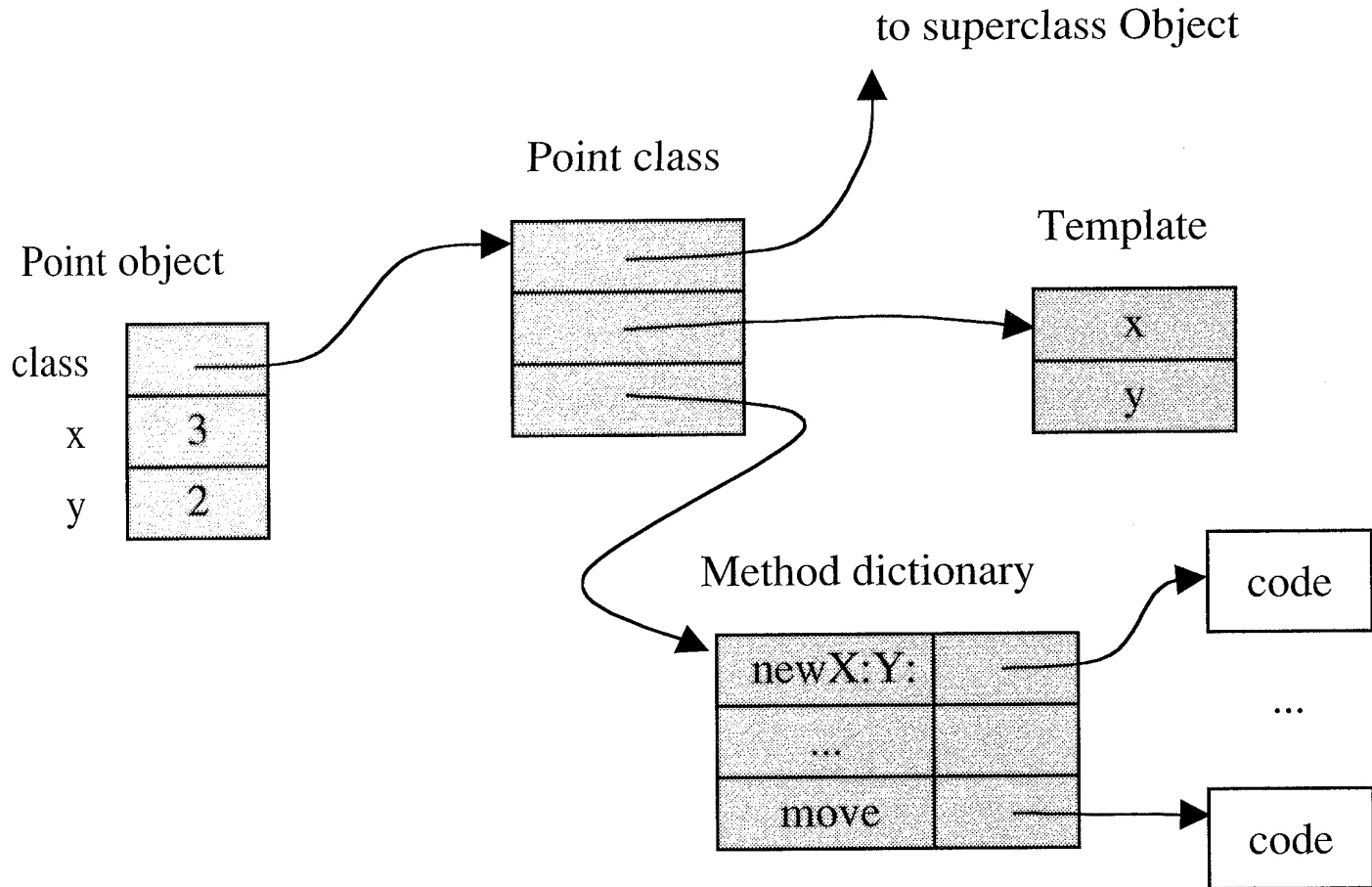
Call changes cell value from 1 to 3

Activation record associated with returned function cannot be deallocated upon function return

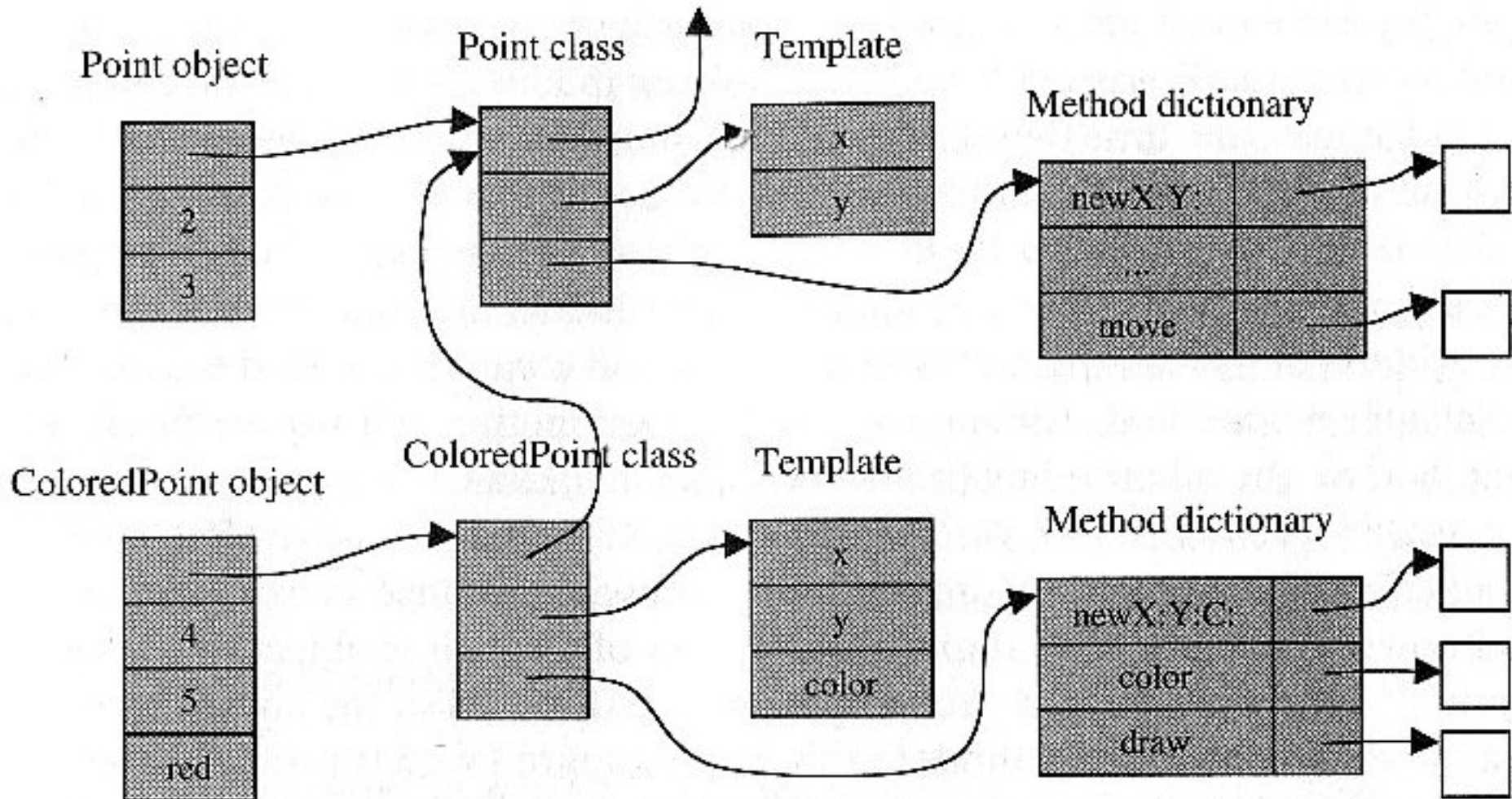
Classes and objects at runtime

- Data for each object stored with the object
 - E.g. x and y coordinates for a point
- Common functionality stored in shared location
 - Methods, static variables

Smalltalk – Point object and class



Smalltalk – runtime support for inheritance



Aside: not all scopes are equal

JAVASCRIPT



```
this.value = 42; //Global variable
```

```
var obj = {  
  value: 0, //Local field in object  
  increment: function() {  
    this.value++;  
    alert(this.value);  
  
    var innerFunction = function() {  
      alert(this.value);  
    }  
  
    innerFunction(); // Function invocation  
  }  
}  
obj.increment(); // Method call
```

What will be shown on screen?

Try it out yourselves: <http://jsfiddle.net/7jxw1r9v/1/>

How do we fix this?

```
this.value = 42; //Global variable
```

```
var obj = {  
  value: 0,  
  increment: function() {  
    this.value++;  
    alert(this.value);  
    var that = this;  
  
    var innerFunction = function() {  
      alert(that.value);  
    }  
  
    innerFunction(); //Function invocation  
  }  
}  
obj.increment(); //Method invocation
```

Why does this help?

Because this function is a closure that captures the «that» variable

JavaScript is not OO, but has objects!

```
var obj = {  
  data:'Hello World'  
}
```

```
var displayData = function() {  
  alert(this.data);  
}
```

```
displayData(); //undefined  
displayData.apply(obj); //Hello World
```

How do you think the access links for JavaScript are implemented?

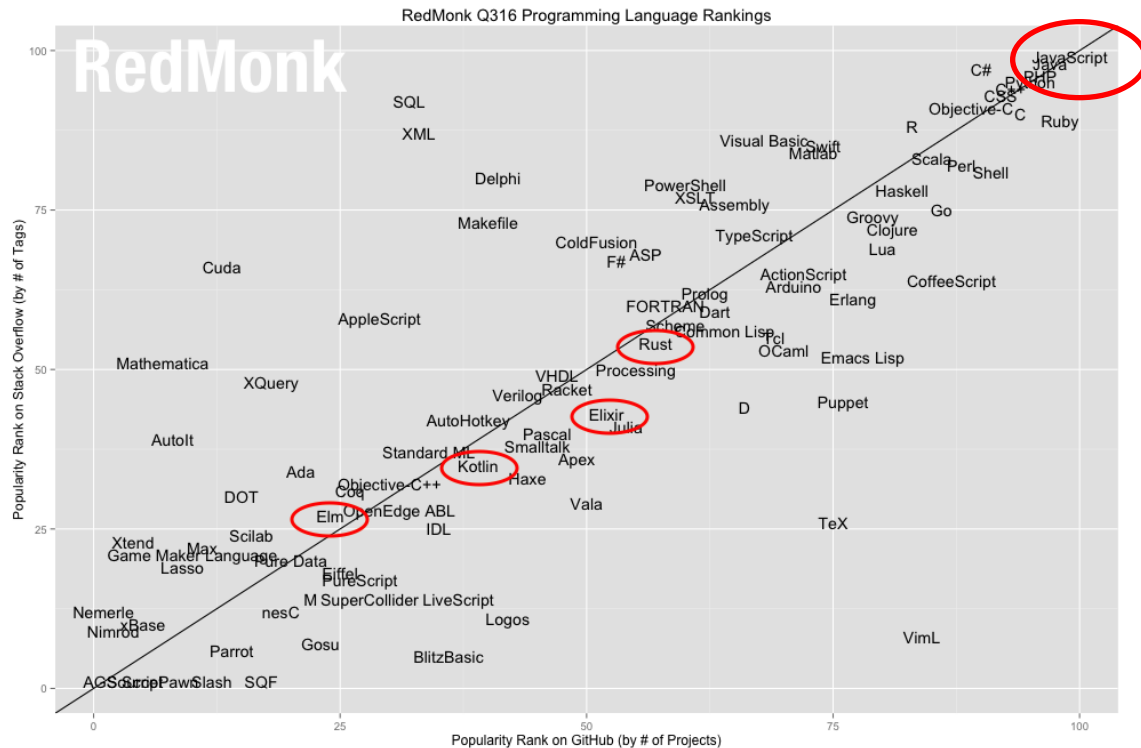
Everyone loves Javascript!

www.tiobe.com/index.php/content/paperinfo/tpci/index.html

Programming Language

The hall of fame listing all "Programming Language of the Year" language that has the highest rise in ratings in a year.

Year	Winner
2014	🏆 JavaScript
2013	🏆 Transact-SQL
2012	🏆 Objective-C
2011	🏆 Objective-C
2010	🏆 Python
2009	🏆 Go



<http://techbeacon.com/5-emerging-programming-languages-bright-future>

<https://www.destroyallsoftware.com/talks/wat>

But WHY all these WATs?

- JavaScript has automatic type coercion
 - It will try to convert types into something that matches the operand!
- $[] + [] = ""$
 - The + operator cannot operate on arrays, so the array is *coerced* to its string representation, which is a toString() of all the elements joined by commas.
- $\{\} + [] = 0$
 - The first is recognized as an empty code block.
 - The plus is thus unary, and [] is coerced to an empty string, which is in turned coerced to 0.
- $\{\} + \{\} = NaN$
 - The first is again an empty code block
 - The second {} is an empty object, which is coerced to [object Object], which is the toString() repr of objects
 - Which is again not a number, or NaN

More fun: scoping and blocks

Java:

```
void main() {  
    Integer x = 1;  
    System.out.println(x);  
    if (true) {  
        Integer x = 2;  
        System.out.println(x);  
    }  
    System.out.println(x);  
}
```

Output: «1», «2», «1»

JavaScript:

```
function main() {  
    var x = 1;  
    console.log(x);  
    if (true) {  
        var x = 2;  
        console.log(x);  
    }  
    console.log(x);  
}
```

Output: «1», «2», «2»!

- JavaScript has blocks, but not block scope!
- Declarations are always «hoisted» to the top of the function

More fun: scoping and blocks

Java:

```
void main() {  
    Integer x = 1;  
    System.out.println(x);  
    if (true) {  
        Integer x = 2;  
        System.out.println(x);  
    }  
    System.out.println(x);  
}
```

Output: «1», «2», «1»

JavaScript, explicit hoisting:

```
function main() {  
    var x;  
    var x;  
    x = 1;  
    console.log(x);  
    if (true) {  
        x = 2;  
        console.log(x);  
    }  
    console.log(x);  
}
```

Output: «1», «2», «2»!

- JavaScript has blocks, but not block scope!
- Declarations are always «hoisted» to the top of the function

Upcoming!

- More ML
- Part two of OO lecture
- Weekend!



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