INF3510 Information Security University of Oslo Spring 2010

Lecture 3 Key Management and PKI



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Outline

Key management
Key establishment
Public key infrastructure
Digital certificates
PKI trust models
PKI components

Key Management

- The security of cryptographically protected information depends on:
 - The strength/size of the keys
 - The robustness of cryptographic algorithms/protocols
 - The protection and management afforded to the keys
- Key management provides the foundation for the secure generation, storage, distribution, and destruction of keys.
- Proper key management is essential to the robust use of cryptography for security.
- Poor key management may easily compromise strong algorithms.

Key Usage

- A single key should be used for only one purpose
 - e.g., encryption, authentication, key wrapping, random number generation, or digital signatures
- Using the same key for two different processes may weaken the security of one or both of the processes.
- Limiting the use of a key limits the damage that could be done if the key is compromised.
- Some uses of keys interfere with each other
 - e.g. an asymmetric key pair should only be used for either encryption or digital signatures, not both.

Types of Cryptographic Keys

- How many types of keys are there?
- They're classified according to:
 - Whether they're public, private or symmetric
 - Their use
 - For asymmetric keys, also whether they're static or ephemeral
 - NIST Special Publication 800-57, Recommendation for Key Management – Part 1: General, August 2005, defines **19** types of cryptographic keys.

Key types

- 1. Private signature key
- 2. Public signature verification key
- 3. Symmetric authentication key
- 4. Private authentication key
- 5. Public authentication key
- 6. Symmetric data encryption key
- 7. Symmetric key wrapping key
- 8. Symmetric and asymmetric RNG seeds/keys (RNG: Random Number Generation)
- 9. Symmetric master key

Key types (continued)

- 10. Private key transport key
- 11. Public key transport key
- 12. Symmetric key agreement key
- 13. Private static key agreement key
- 14. Public static key agreement key
- 15. Private ephemeral key agreement key
- 16. Public ephemeral key agreement key
- 17. Symmetric authorization key
- 18. Private authorization key
- 19. Public authorization key

Random Number Generator Seeds

- RNG keys are used to initialise the generation of random symmetric/asymmetric keys
- Knowing the seed may determine the key uniquely
 - Requires confidentiality and integrity protection
 - Period of protection, e.g.:
 - a. Used once and destroyed
 - b. From generation until no longer needed for subsequent reseeding
 - c. Shall be destroyed at the end of the protection period

Additional Key Parameters

- Domain parameters
- Initialization vectors
- Shared secrets
- RNG seeds
- Other public information;
 e.g. nonce

- Intermediate results
- Key control information;
 e.g. identifier, purpose,
 counter
- Random numbers
- Passwords
- Audit information

Crypto Period

- The crypto period is the time span during which a specific key is authorized for use
- The crypto period is important because it:
 - Limits the amount of information protected by a given key that is available for cryptanalysis.
 - Limits the amount of exposure if a single key is compromised.
 - Limits the use of a particular algorithm to its estimated effective lifetime.

Crypto Period (continued)

The crypto period is important because it:

- Limits the time available for attempts to penetrate physical, procedural, and logical access mechanisms that protect a key from unauthorized disclosure.
- Limits the period within which information may be compromised by inadvertent disclosure of keying material to unauthorized entities.
- Limits the time available for computationally intensive cryptanalytic attacks. (in applications where long-term key protection is not required).

Factors Affecting Crypto Periods

- In general, as the sensitivity of the information or the criticality of the processes increase, the length of the cryptoperiods **should** decrease in order to limit the damage that might result from each compromise.
- Short cryptoperiods may be counter productive, particularly where denial of service is the paramount concern, and there is a significant potential for error in the re-keying, key update or key derivation process.

Factors Affecting Crypto Periods (cont.)

Communications versus Storage

- Cryptoperiods are generally made shorter for keys used for protection of communication, e.g. limited to a single session
- Cryptoperiods are generally made longer for stored data because the overhead of re-encryption associated with changing keys may be burdensome.

Cost of Key Revocation and Replacement

- In some cases, the costs associated with changing keys are painfully high. (e.g. decryption and subsequent re-encryption of very large databases, and revocation and replacement of a very large number of keys)
- In such cases, the expense of the security measures necessary to support longer cryptoperiods may be justified.

Key Usage Periods

- A key is used for both <u>protecting</u> and <u>processing</u> the protected information.
- The protection period is called the originator usage period.
- The processing period is called the recipient usage period.
- A symmetric key shall not be used to provide protection after the end of the originator usage period.
- The recipient usage period normally extends beyond the originator usage period..
- The cryptoperiod of a symmetric key is the period from the beginning of the originator usage period to the end of the recipient usage period.

Ref: NIST SP 800-57

Кеу Туре	Cryptoperiod		
	Originator Usage Period (OUP)	Recipient Usage Period (RUP)	
1. Private Signature Key	1-3 years		
2. Public Signature Key	Several years (depends on key size)		
3. Symmetric Authentication Key	<= 2 years	<= OUP + 3 years	
4. Private Authentication Key	1-2 years		
5. Public Authentication Key	1-2 years		
6. Symmetric Data Encryption Keys	<= 2 years	<= OUP + 3 years	
7. Symmetric Key Wrapping Key	<= 2 years	<= OUP + 3 years	

Ref: NIST SP 800-57

Кеу Туре	Cryptoperiod	
	Originator Usage Period (OUP)	Recipient Usage Period (RUP)
8. Symmetric and asymmetric RNG Keys	Upon reseeding	
9. Symmetric Master Key	About 1 year	
10. Private Key Transport Key	<= 2 years	
11. Public Key Transport Key	1-2 years	
12. Symmetric Key Agreement Key	1-2 years	
13. Private Static Key Agreement Key	1-2 years	

Ref: NIST SP 800-57

Кеу Туре	Cryptoperiod	
	Originator Usage Period (OUP)	Recipient Usage Period
14. Public Static Key Agreement Key	1-2 years	
15. Private Ephemeral Key Agreement Key	One key agreement transaction	
16. Public Ephemeral Key Agreement Key	One key agreement transaction	
17. Symmetric Authorization Key	<= 2 years	
18. Private Authorization Key	<= 2 years	
19. Public Authorization Key	<= 2 years	

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Keys and Keying Material Compromise

- Key compromise occurs when the protective mechanisms for the key fail, and the key can no longer be trusted
- When a key is compromised, all use of the key to protect information shall cease and the compromised key shall be revoked.
 - However, the continued use of the key under controlled circumstances to remove or verify the protections may be warranted, depending on the risks of continued use and an organization's Key Management Policy.
- The continued use of a compromised key shall be limited to processing protected information.
 - In this case, the entity that uses the information shall be made fully aware of the dangers involved.

Key Compromise Recovery Plan

- A compromise recovery plan **should** contain:
 - The identification of the personnel to notify.
 - The identification of the personnel to perform the recovery actions.
 - The re-key method.
 - Any other recovery procedures, such as:
 - Physical inspection of equipment.
 - Identification of all information that may be compromised.
 - Identification of all signatures that may be invalid due to the compromise of a signing key.
 - Distribution of new keying material, if required.

Undetected Key Compromise

- The worst form of key compromise is when it is not detected.
- Nevertheless certain protective measures can be taken.
 Key management systems (KMS) should be designed:
 - to mitigate the negative effects of a key compromise.
 - so that the compromise of a single key has limited consequences, e.g., a single key could be used to protect only a single user or a limited number of users, rather than a large number of users.
 - Often, systems have alternative methods to authenticate communicating entities that do not rely solely on the possession of keys.
 - Avoid building a system with catastrophic weaknesses.



Key State Transitions

 Transitions between states are triggered by events, such as the expiration of a cryptoperiod or the detection of a compromise of a key.



1) Pre-Activation

 The key has been generated but is not yet authorized for use

4) Active

The key may be used to cryptographically protect information or cryptographically process previously protected information, or both. When a key is active, it may be designated to protect only, process only, or both.



6) Deactivated

 A key whose cryptoperiod has expired but is still needed to perform cryptographic processing is deactivated until it is destroyed.

2), 7) Destroyed

The key is destroyed. Even though the key no longer exists in this state, certain key attributes (e.g. key name, type and cryptoperiod) may be retained.



3), 5), 8) Compromised

Generally, keys are compromised when they are released to or determined by an unauthorized entity. If the integrity or secrecy of the key is suspect, it is revoked. The key is not used to apply protection to information. In some cases, the key may be used for processing.



9), 10) Destroyed Compromised

 The key is destroyed after a compromise, or the key is destroyed and a compromise is later discovered. Key attributes may be retained.

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Key Management Phases Ref: NIST SP 800-57



Management State Transitions

Key management determines the state and usage of keys. Each key state requires a different management procedure.

Key Management Phases (cont.) Ref: NIST SP 800-57



1) Pre-operational Phase

The keying material is not yet available for normal cryptographic operations.
Keys may not yet be generated, or may be in the pre-activation state. System or enterprise attributes are established during this phase as well.

4) Operational Phase

 The keying material is available and in normal use. Keys are in the active state. Keys may be designated as protect only, process only, or protect and process.

Key Management Phases (cont.) Ref: NIST SP 800-57



3), 5), 6) Post-operational Phase

The keying material is no longer in normal use, but access to the keying material is possible and the keying material may be used for process only in certain circumstances. Keys are in the deactivated or compromised states. Keys are archived when not processing data.

2), 7) Destroyed Phase

Keys are no longer available. All records of their existence may have been deleted. Although the keys themselves are destroyed, the key attributes may be retained.

Key Management States and Phases Ref: NIST SP 800-57



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Key Generation

- Most sensitive of all cryptographic functions.
- Need to prevent unauthorized disclosure, insertion, and deletion of keys.
- Automated devices that generate keys and initialisation vectors (IVs) should be physically protected to prevent:
 - disclosure, modification, and replacement of keys,
 - modification or replacement of IVs.
- Keys should be randomly chosen from full range of key space.

Key Generation Examples

- Stream cipher
 - capture true random stream (OTP), or
 - generate pseudorandom stream with keystream generator using initial short random key (initial state or seed).
- AES symmetric block cipher
 - ensure key is as probable as any other
- RSA asymmetric cipher
 - make sure *n* (modulus) is too large to be factored

Key Protection

Keys should be

- accessible for authorised users,
- protected from unauthorised users
- Old keys must be kept, if messages encrypted under these keys are stored
 - Where will they be kept?
 - How will they be kept securely?
 - Who will know how to access them when required?

Key Protection Examples

symmetric ciphers

- Never stored or transmitted 'in the clear'
- May use hierarchy: session keys encrypted with master
- Master key protection:
 - Locks and guards
 - Tamper proof devices
 - Passwords/passphrases
 - Biometrics

asymmetric ciphers

- Need to protect private keys only
- Public keys made public but must be authentic

Key destruction

- No key material should reside in volatile memory or on permanent storage media after destruction
- Key destruction methods, e.g.
 - Simple delete operation on computer
 - may leave undeleted key e.g. in recycle bin or on disk sectors
 - Special delete operation on computer
 - that leaves no residual data, e.g. by overwriting
 - Magnetic media degaussing
 - Destruction of physical device e.g high temperature
 - Master key destruction which logically destructs subordinate keys

Outline

Key management <u>Key Establishment</u> Public key infrastructure Digital certificates PKI trust models PKI components

Key Establishment: The Problem

- Two parties wish to communicate securely using cryptography.
- Symmetric ciphers are much more efficient than asymmetric ciphers.
 - In practice, for large amounts of data, symmetric ciphers are preferred over asymmetric ciphers.
- Therefore need to establish a new shared key: a session key.
Key Establishment: Options

Three options to distribute session key

- 1. Use existing shared keys
 - Only possible in small group (see next slide)
- Use a trusted third party (server) who shares a symmetric (long-term) key with each user
 - Need protocol such as Kerberos (lecture on authent.)
- 3. Use asymmetric cipher to protect session key
 - Requires PKI (see later)

Key Establishment: Simple 'Out-Of-Band' Solution

- Communicating parties must share a secret key.
- With *n* participants in a network, then each participant needs *n*-1 long-term shared secret keys.
- Total number of keys to be exchanged 'out-of-band' is (n-1) + (n-2) + ... + 2 + 1
 - = n(n-1)/2
 - e.g. for *n*=100, 4950 keys must be established.



Outline

Key management Key Establishment <u>Public-key infrastructure</u> Digital certificates PKI trust models PKI components

Why the interest in PKI?

Cryptography solves security problems in open networks, ... but creates key management complexity.



Public-key cryptography simplifies the key management, ... but creates trust management problems.

Public Keys and the Spoofing Problem

What happens when an attacker inserts false keys in the public-key register?



How does this affect the security of:

- a digital signature on messages received from A? Why?
- a confidential message sent to A? Why?

Digital signature notations

- Cryptographic terminology and notation:
 - Private Key K_{priv}: confidential key only known by owner
 - Public Key K_{pub}: publicly known key
 - Plaintext message M: the original message or data
 - Hash function H: used to create hash block
 - Digital signature Sig: cryptographic authentication code
 - Signature generation S: Function for creating the digital signature Sig on message M or on hash H(M)
 - Recovery function R: Function for recovering the message M or hash H(M) from the digital signature Sig

Digital signature



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Public-key infrastructure

Integrity and trustworthiness of public keys:

- How can a user be sure who a public key belongs to?
- How can a user be sure a public key has not been altered - intentionally or unintentionally?
- How can public keys be made available in a trusted way?
 - Use public-key certificates (aka. digital certificates) issued by a trusted third party
 - a Certification Authority (CA).
 - A public-key certificate is a public key digitally signed by a CA.
 - A hierarchy of public-key certificates becomes a PKI.

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Public-key infrastructure

- Public-key cryptography needs a PKI to work
- A PKI is a set of
 - Policies (to define the rules for managing certificates)
 - Technologies (to implement the policies and generate, store and manage certificates)
 - **Procedures** (related to key management)

that enables practical application of public key cryptography often in large, distributed settings

Public-Key Certificates

- A public-key certificate is simply a public key with a digital signature
- Certification Authorities (CA) sign public keys.
- The CA's public key is needed in order to validate certificates
- Relying party is an entity that needs to validate a certificate (i.e. to verify that the public key is authentic)

X.509 Digital Certificate

- Version
- Serial Number
- Algorithm Identifier
- CA Name
- CA Unique Identifier
- User Name
- User Unique Identifier
- Public Key
- Validity Period

CA Digital

Signature

ZYVY

Extensions

How to generate a digital certificate?

- 1. Assemble the information in single record Rec
- 2. Hash the record
- 3. Sign the hashed record
- 4. Append the digital signature to the record



Public-key certificates

- Public-key certificates (aka. digital certificates) logically bind a public key to an identifier ("distinguished name")
- A digital certificate contains:
 - the user's public key
 - the user's identifier
 - + some other information e.g. validity period
- A Certificate Authority (CA) creates and digitally signs the certificate

Digital certificates Digital certificates in use

• X.509 standard

- most widely used standard (still evolving: now v3)
- based on X.500 Distinguished Naming (DN) scheme
- Important fields in X.509 digital certificates are:
 - Version number
 - Serial Number (set by the CA)
 - Signature Algorithm identifier (Algorithm used for dig sigs)
 - Issuer distinguished name (Name of the CA)
 - Subject distinguished name (Name of certificate owner)
 - Public key
 - Validity period (certificate should not be used outside this time)
 - Digital signature (of the certificate, signed by the CA)

Example of X.509 Certificate

Certificate	Certificate		2 🛛
General Details Certification Path	General	Details Certification Path	
Certificate Information	Show: <	(All>	
 This certificate is intended for the following purpose(s) Protects e-mail messages Proves your identity to a remote computer Ensures software came from software publisher Protects software from alteration after publication 2.16.840.1.113733.1.7.1.1 	Field	Value Il number 52 c8 20 13 7c ature algorithm md5RSA er Class 2 Public P from Tuesday, May to Wednesday, 12	85 a7 ed f2 17 Primary Certifica 12, 1998 10:00
* Refer to the certification authority's statement for details. Issued to: VeriSign Class 2 CA - Individual Subscriber	Subje	ect VeriSign Class 2 c key RSA (1024 Bits) cape Cert Type SSL CA, SMIME	2 CA - Individua) : CA (06)
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; Valid from 5/12/1998 to 1/7/2004	5f a2 cf a5 22 99 48 99 e2 7a	23 1C 69 4d b8 3b db /d d 58 3a b6 4b c4 d4 db d8 a 22 01 28 60 a5 db d5 30 d ad 21 54 91 d1 de 5f fb 3 53 58 c5 0d 5d 13 07 b3 5	la 8f C1 09 ke 75 fa 86 lf 21 70 5e 18 29 53 1b 50 c4 06 4b 🕑
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Using certificates to send confidential messages

Alice sends confidential message C = (M, K_{pub}(A)) to Bob

- 1. Alice is the relying party and must get Bob's public key
 - Alice can obtain Cert_B
 - Alice validates Cert_B
 - Alice obtains K_{pub}(B) from Cert_B
- 2. Alice uses K_{pub}(B) to encrypt message M
- If Alice
 - <u>trusts</u> the CA that issued $Cert_B$ (i.e. to be competent and honest)
 - and is certain of CA's public key and unique identifier
 - and is certain of Bob's unique identifier
- then Alice can be sure that only Bob will be able to decrypt the message

Using certificates to verify signatures

Bob sends a signed message {M, Sig_B, Cert_B} to Alice

- 1. Alice is the relying party and must first validate Cert_B
 - Alice uses CA's public key K_{pub}(CA) to verify CA's signature on the binding between the public key and Bob's unique identifier.
- 2. Alice obtains $K_{pub}(B)$ from the certificate Cert_B
- 3. Alice uses $K_{pub}(B)$ to verify signature Sig_B on M
 - If Alice
 - <u>trusts</u> the CA that issued Cert_B (i.e. to be competent and honest)
 - and is certain of CA's public key and unique identifier
 - and is certain of Bob's unique identifier
- then Alice is certain that message M came from Bob

Digital certificates Some questions

- 1. What advantage is there for Bob in having digital certificate $Cert_B$?
- 2. Who can have access to $Cert_B$?
- 3. Why would someone want to verify the signature in Cert_B?
- 4. What does Alice need for verifying $Cert_B$?
- 5. After someone has verified Cert_B, of what can they be assured?

PKI trust models

- Q: Can Alice trust Bob's digital certificate?
- Q: Can Alice trust the CA that issued the cert. and how does she get the CA's public key?
 - To answer these questions we need to
- Establish trust relationships
 - between different Certificate Authorities, and
 - between Certificate Authorities and end users,

Define PKI trust models

PKI Certification Path



 For Alice to validate CertB and be assured of the authenticity of Bob's public key, Alice needs to create a path of trust from Alice to Bob's key

PKI Certification Path Explanation

- Alice receives Bob's digital certificate Cert_B with Bob's public key issued by CA_Z
- Alice needs to validate Bob's certificate
 - Cert_B can be verified with Cert_Z
 - Cert_z can be verified with Cert_Y
 - Cert_Y can be verified with Cert_X
- Alice trusts CA_X and has an authentic copy of CA_X's public key K_{pub}(CA_X)
- A certification path is established !
- The relying party Alice can validate Cert_B

Root CA's and Self Signed Certificates

- A root CA is at the top of a hiearchy
- No other CA can certify the root CA's public key
- Self-signed certificate is created by signing the root CA public key with the corresponding private key
 - Semantically <u>not</u> a digital signature and <u>not</u> a certificate
 - Syntactically looks like a certificate
 - Enables processing in PKI software
 - <u>Must be sent securely out-of-band</u>



PKI Trust Models



PKI trust models Strict hierarchical model



PKI trust models Strict hierarchical model

- Highly regulated
 - each CA must follow rules regarding to whom they may issue certificates
- Tree structure
 - Single root CA
 - Users are leaves of the tree
 - Each node is certified by its immediate parent CA
 - Root CA
 - Starting point for trust
 - All users trust the root CA, and must receive its public key through a secure out-of-band channel

PKI trust models Strict hierarchical model

- Advantages:
 - works well in highly-structured setting such as military and government
 - unique certification path between two entities (so finding certification paths is trivial)
 - scales well to larger systems
- Disadvantages:
 - need a trusted third party (root CA)
 - 'single point-of-failure' target
 - If any node is compromised, trust impact on all entities stemming from that node
 - Does not work well for global implementation (who is root TTP?)

PKI trust models User-centric model





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- Each user is completely responsible for deciding which public keys to trust
- Example: *Pretty Good Privacy (PGP)*
 - 'Web of Trust'
 - Each user may act as a CA, signing public keys that they will trust
 - Public keys can be distributed by key servers and verified by fingerprints
 - OpenPGP Public Key Server: http://pgpkeys.mit.edu:11371/
- PGP or GPG What is the difference?

PKI trust models User-centric model



Advantages:

- Simple and free
- Works well for a small number of users
- Does not require expensive infrastructure to operate
- User-driven grass roots operation

Disadvantages:

- It relies on human judgment
 - Works well with technical users who are aware of the issues, but not the general public
- Not appropriate for more trust-sensitive areas such as finance and government

PKI Trust Models Browser PKI

 The browser PKI model consists of isolated strict hierarchies where the (root) CA certificates are installed together with the browser



Browser PKI root certificate installation

- Distribution of root certificates which should happen securely out-of-band, is now done through online downloading of browser SW
- Users are in fact trusting the browser vendor who supplied the installed certificates, rather than a root CA
- Example: used by Mozilla Firefox and Microsoft Internet Explorer
- CAs pay browser vendors to carry CA certs

Browser PKI limitations and weaknesses

- Certification path processing is limited
 - Certificate validation only by stored certificates
- List of trusted certificates controlled by user not well protected from modification attacks
 - Users habitually accept incoming certificates that the browser cannot automatically verify
 - Malware can install fake server and root certificates
- Cross certification not supported
 - Certificate Revocation not supported
- No legal agreement between users and CAs

 Liability rests with the users and not with the CAs

Authenticating a Web Site

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000		General	Forms Links	Media Security		
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Google Scholar	ISI Webma	Web Site Identity		General Details		
Mail :: INBOX The web site online		This certificate has been	verified for the following use	25:		
Online Banking	Sig	Trust Network, a c	SSL Server Certificate			
Apply Online	1. En Usi		SSL Server with Step-up			
Home Loans Personal Loans	Se					
Credit cards	C		Issued To			
Accounts Investment & Superannuation Insurance	2. En Usi		Common Name (CN) Organization (O) Organizational Unit (OU) Serial Number	online.westpac.com.au Westpac Banking Corporation Terms of use at www.verisign.c 2E:E8:A1:C0:C2:5E:1C:6A:CE:6	com/rpa (c)00 3:91:0E:14:E4:80:CB	
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		information traveli	Fingerprints			1.00
		anyone read this p	SHA1 Fingerprint MD5 Fingerprint	F8:DE:04:AE:35:68:E0:D5:18:73 05:5C:C1:2A:6B:20:03:46:B1:51	3:1E:46:4F:86:9D:8A:FE:F7:17:C3 F:C8:F2:63:C1:8F:C7	
	General ad situation or <u>documents</u> , when decid					
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Phishing and fake certificates Hawaii Federal Credit Union



Authentic bank login

https://hcd.usersonInet.com/asp/USERS /Common/Login/NettLogin.asp

Fake bank login

https://hawaiiusafcuhb.com/cgibin/mcw00.cgi?MCWSTART

Authentic and Fake Certificates

Certificate ?X	Certificate ?X
General Details Certification Path	General Details Certification Path
Certificate Information This certificate is intended for the following purpose(s): •Ensures the identity of a remote computer	Certificate Information This certificate is intended for the following purpose(s): •Ensures the identity of a remote computer
* Refer to the certification authority's statement for details. Issued to: hcd.usersonInet.com	* Refer to the certification authority's statement for details. Issued to: www.hawaiiusafcuhb.com
Issued by: Class 3 Open Financial Exchange CA - G2	Issued by: VeriSign Class 3 Secure Server CA
Valid from 19/08/2006 to 13/09/2007	Valid from 29/11/2006 to 15/12/2009
Install Certificate Issuer Statement OK	Install Certificate Issuer Statement OK
Authentic certificate	Fake certificate

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Browser PKI and Fake Certificates

- The browser validates server certificates by
 - using public key of root certificate stored in browser
 - checking that the server certificate distinguished name and the domain name of web server are the same
- Attackers buy legitimate certificates which are automatically validated by browsers
 - Fake certificates are legitimate certificates !!!
- Server certificate validation is not authentication
 - Users who don't know the server domain name cannot distinguish between authentic and fake server certificates

PKI components

- Common PKI components include:
 - Certification authorities (CA)
 - Certification Practice Statement (CPS)
 - Aka. Certification Policy (CP)
 - Registration authorities (RA)
 - Validation authorities (VA)
 - Certificate Revocation Lists (CRL)
PKI components

- Certification authorities
 - Are the primary building elements of a PKI
 - Create, issue, and revoke certificates for subscribers and other CAs
 - Have a Certification Practice Statement (CPS)
 - legally binding document
 - must be approved before CA can operate
 - outlines practices and procedures that the CA will follow in issuing certificates

PKI components CPS and CP

- Certification Practice Statement may specify:
 - ID checks performed before certificate issue
 - physical, personnel and procedural security controls for the CA
 - technical and key pair protection and management controls
 - certificate revocation management procedures
 - audit procedures for the CA
 - accreditation information
 - legal and privacy issues and liability limitations
 - profiles and descriptions covering issued certificates including naming conventions

PKI components

Registration Authorities

- act as interface between user and CA
 - · collects details from certificate requesters
 - authenticates identity, and
 - submits certificate request to CA
- Use of RA is optional, CA may employ an RA
 - can provide cost benefits for CA
 - quality of RA determines level of trust in issued certificates

PKI components Validation Authority

- Q: What if the relying party needs to validate a certificate but can't establish a trust path to it?
- A: The relying party can use a validation service from a VA (Validation Authority) to validate the certificate
 - The relying party needs to trust the VA
 - The relying party needs an authentic copy of the VA's public key, e.g. received securely out-of-band
- The VAs must establish trust relationships and get authentic public keys from Root CAs of many different PKIs

PKI Components Validation Authorities



PKI components: CRL

Certificate Revocation

- Q: When might a certificate need to be revoked ?
- A: When certificate becomes outdated <u>before</u> it expires, due to:
 - private key being disclosed
 - subscriber name change
 - change in authorisations, etc
- Revocation may be checked online against a certificate revocation list (CRL)
- Checking the CRL creates a huge overhead which threatens to make PKI impractical

Standardisation Activities



- X.509 specifies the the overall structure of certificates
- IETF has specified a series of standards for implementing PKIs based on X.509 certificates

PKI services

Several organisations operate PKI services

- Private sector
- Public sector
- Military sector
- Mutual recognition and cross certification between PKIs is difficult
- Can be expensive to operate a robust PKI
- The Browser PKI is the most widely deployed PKI thanks to piggy-backing on browsers and the lax security requirements

PKI Summary

- Public key cryptography needs a PKI to work
 - Digital certificates used to provide integrity for public keys
 - Acceptance of certificates requires trust
 - Trust relationships between entities in a PKI can be modelled in different ways
 - Establishing trust has a cost, e.g. because secure out-of-band channels are expensive