# INF3510 Information Security University of Oslo Spring 2016

# Lecture 9

Identity Management and Access Control

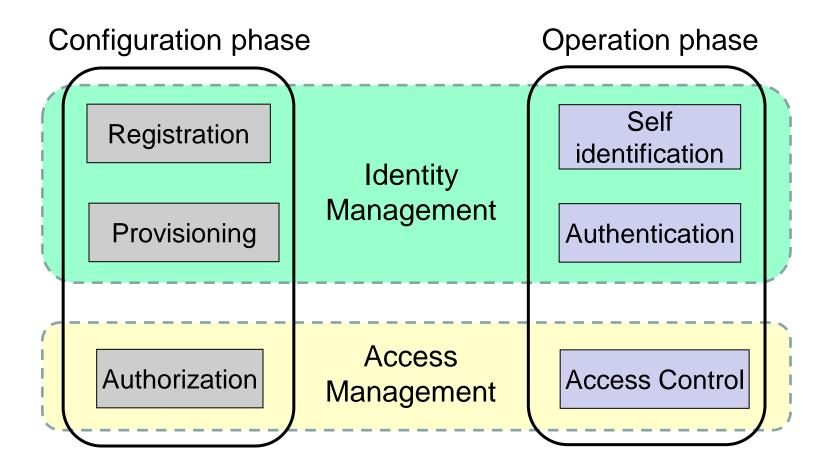


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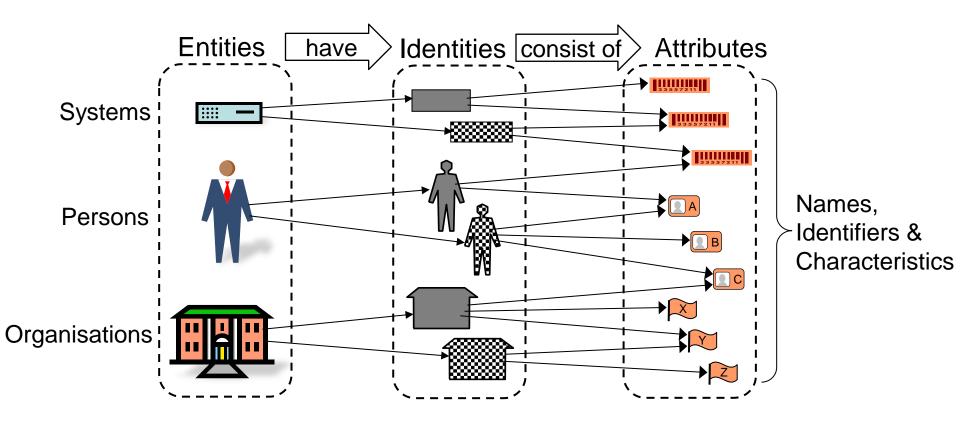
#### **Outline**

- Identity and access management concepts
- Identity management models
- Access control models (security models)

# IAM Identity and Access Management



# The concept of identity



# Concepts related to identity

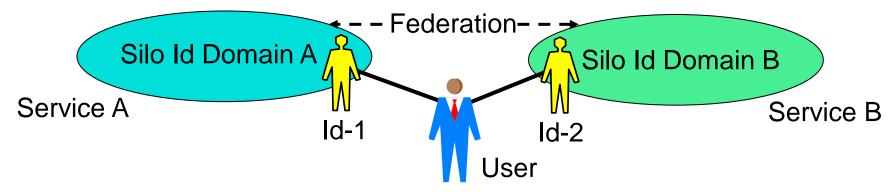
- Entity
  - A person, organisation, agent, system, etc.
- Identity
  - A set of names / attributes of entity in a specific domain
  - An entity may have identities in multiple domains
  - An entity may have multiple identities in one domain
- Digital identity
  - Digital representation of names / attributes in a way that is suitable for processing by computers
- Names and attributes of entity
  - Can be unique or ambiguous within a domain
  - Transient or permanent, self defined or by authority, interpretation by humans and/or computers, etc

# **Identity**

- Etymology (original meaning of words)
  - "identity" = "same one as last time".
- "First-time" authentication is not meaningful
  - because there is no "previous time"
- Authentication requires a first time registration of identity in the form of a name within a domain
- Registration can be take two forms:
  - pre-authentication, from previous identity, e.g. passport
  - creation of new identity, e.g. New born baby

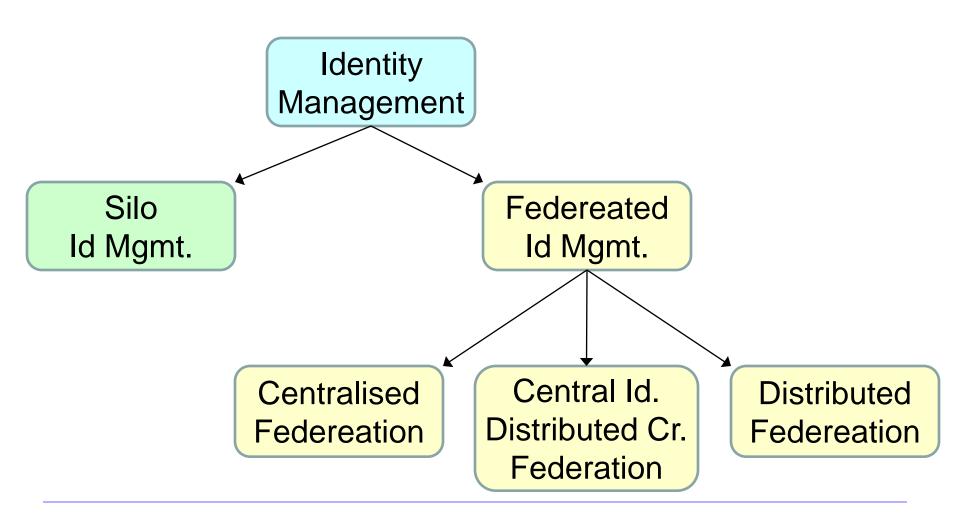
# **Identity Domains**

- An identity domain has a name space of unique names
  - Same user has separate identities in different domains

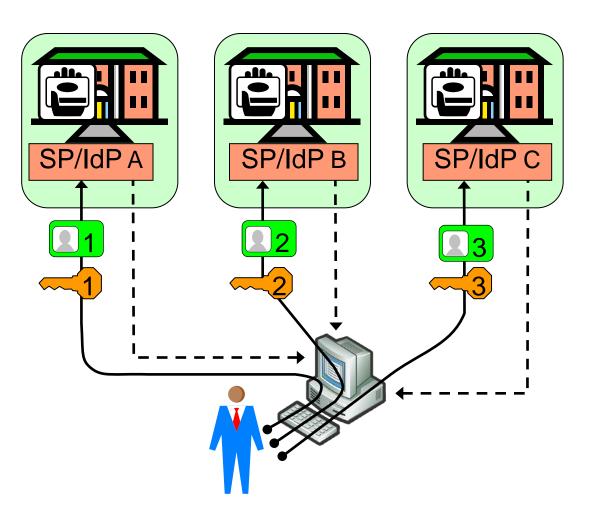


- Identity domain structure options:
  - Silo domain with single authority, e.g. User Ids in company network
  - Distributed hierarchic domain: e.g. DNS (Domain Name System)
- Federation of identity domains
  - Multiple domains have single identity for same user
  - Requires alignment of identity policy between domains

# Taxonomy of Identity Management Architectures



# Silo identity management model



#### Legend:





IdP



Identity domain



User identifier for silo domain



Authentication token for silo domain

→ Service logon

--- Service provision

#### Silo Id domains

- SP (Service Provider) = IdP (Identity Provider):
   SP controls name space and provides access credentials
- Unique identifier assigned to each entity
- Advantages
  - Simple to deploy, low cost for SPs
- Disadvantages
  - Identity overload for users, poor usability, lost business

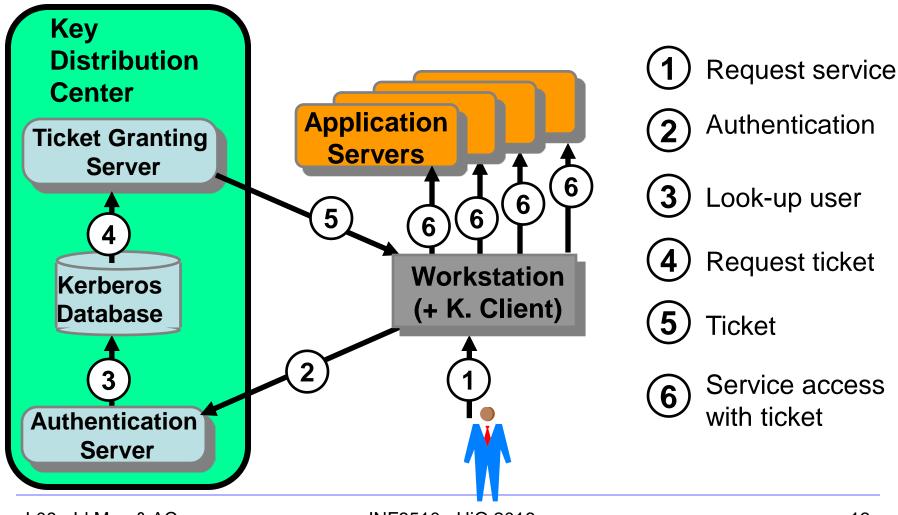
# Single Id and SSO (Single Sign-On)

- Users don't want more identifiers and credentials
- Low acceptance of new services that require separate user authentication
- Silo model requires users to provide same information to many service providers
- Silo model makes it difficult to offer bundled services, i.e. from different service providers
- Service providers want to bundle and collect user information

#### Kerberos SSO

- Part of project Athena (MIT) in 1983.
- User must authenticate once at the beginning of a workstation session (login session).
- Server then authenticates Kerberos client on user's workstation instead of authenticating the user
  - So user does not need to enter password every time a service is requested!
- Every user shares a password with the AS (Authentication Server)
- Every SP (service provider) shares a secret key with the TGS (Ticket Granting Server)
- Tickets are sealed (encrypted) by TGS proves to SPs that the user has been authenticated

#### Kerberos – simplified protocol



# Kerberos – Advantages and limitations

- First practical SSO solution
- Centralized TTP (Trusted Third Party) model
- Uses only symmetric cryptography
- Requires Kerberos clients and servers + KDC
- Only suitable for organisations under common management (single domain)
- Does not scale to very large domains
- Not suitable for open environments (Internet)

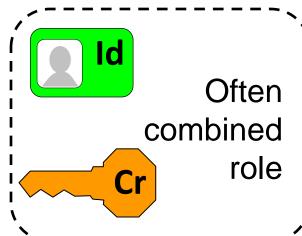
### **Identity Federation Roles**

#### User

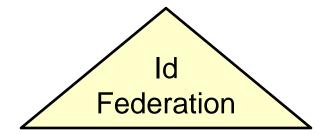
- Needs identities and credentials to access multiple SPs.
- Service Provider (SP)
  - Needs to know identity of users, and needs assurance of authenticity.
- Identity Provider (IdP)
  - Controls name space of identities.
     Issues/registers identities for users.
- Credentials Provider (CrP)
  - Issues/registers credentials for users.
     Performs authentication of users.







#### **Identity Federation**



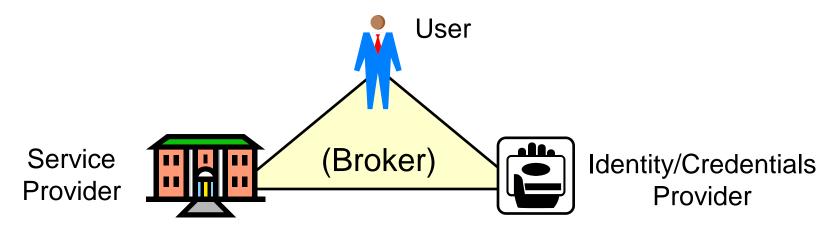
- Identity Federation
  - A set of agreements, standards and technologies that enable a group of SPs to recognise and trust user identities and credentials from different IdPs, CrPs and SPs.
- Three main architectures:
  - **1.Centralized Federation:** Centralised mgmt of name space and credentials by single IdP/CrP.
  - **2. Distributed Federation:** Distributed mgmt of name space and credentials by multiple IdPs and CrPs. Normally combined IdP/CrP.
  - 3. Centralised Identity with Distributed Authentication: Centralised mgmt of name space by single IdP. Distributed mgmt. of credentials and authentication by multiple CrPs.

# Examples of Identity Federations

Federation types	Centralised Identity	Distributed Identity
Centralised Authentication	Centralised	Distributed Id Central Cr
	AADHAAR	
	Central Id Distributed Cr	Distributed
Distributed Authentication	ID-porten ∷⊡≡ altinn	eduroam OpenID®

#### Federation protocols

- Authentication by one IdP/CrP/SP is communicated as a security assertions (cryptographic token) to other SPs that trust and accept the assertion of authenticity.
- Usually based on SAML protocol
  - Security Assertions Markup Language
- Involves multiple entities
  - User, IdP, CrP, SP, and sometimes broker entity



#### Advantage/Disadvantage of Federation

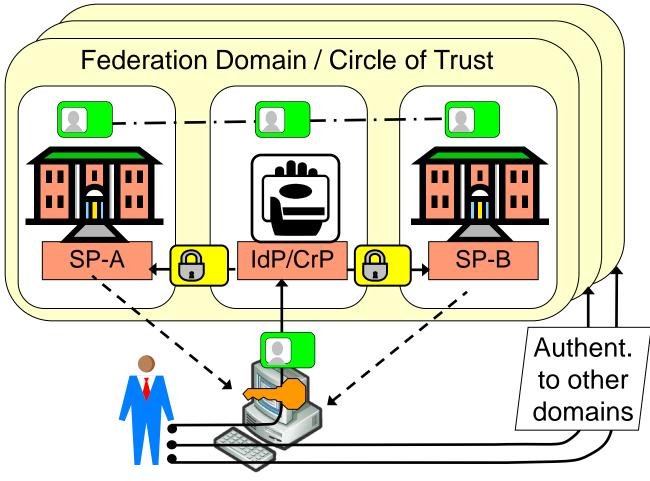
#### Advantages

- Improved usability
- Allows SPs to bundle services and collect user info
- Strengthen privacy through pseudonym identities

#### Disadvantages

- High technical and legal complexity
- High trust requirements
  - E.g. IdP is technically able to access SP on user's behalf
- Privacy issues,
  - IdP collects info about user habits wrt. which SPs are used
- Limited scalability,
  - Limited by political and economical constraints
  - An Identity federation becomes a new form of silo

#### **Centralised Federation**



#### Legend:



SP



IdP/CrP



Identity domain



User identifier issued by IdP



Authentication cred. managed by IdP



Security assertion issued by IdP



→ Service logon



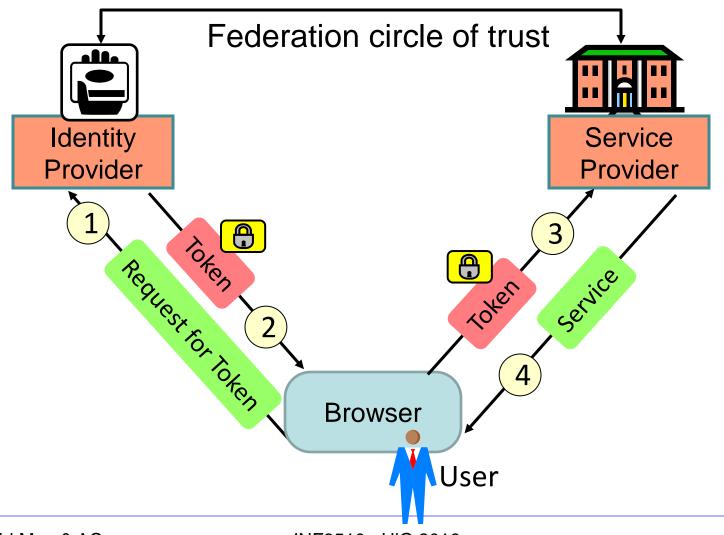
Service provision



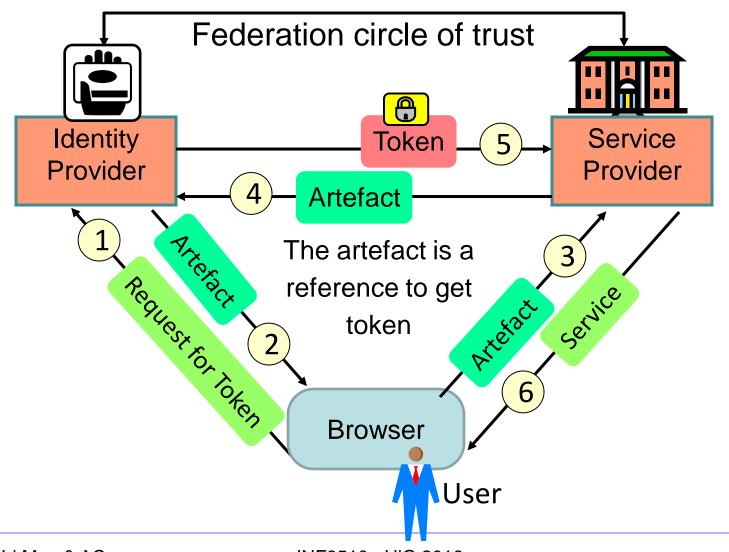
· Identifier mapping

Examples: Facebook connect

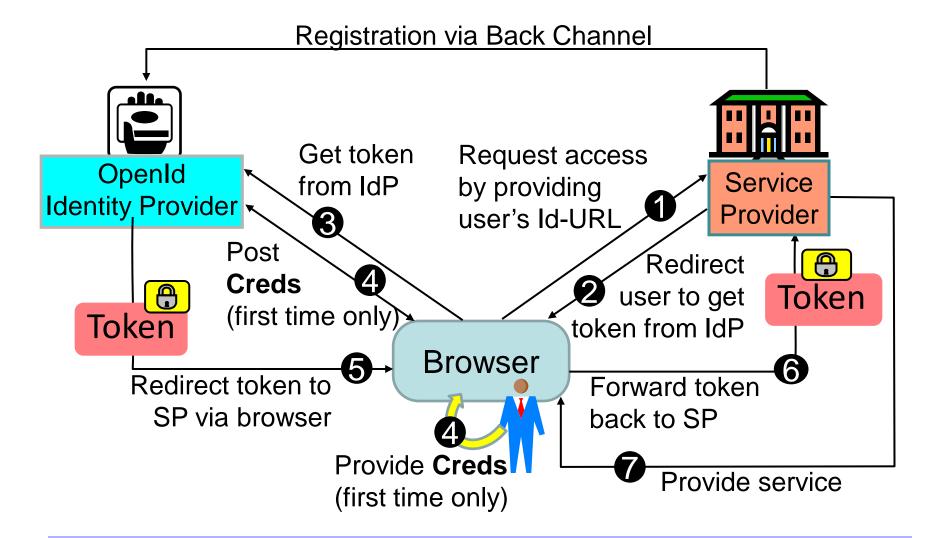
# SAML protocol profile: Browser Post Security token via front-channel



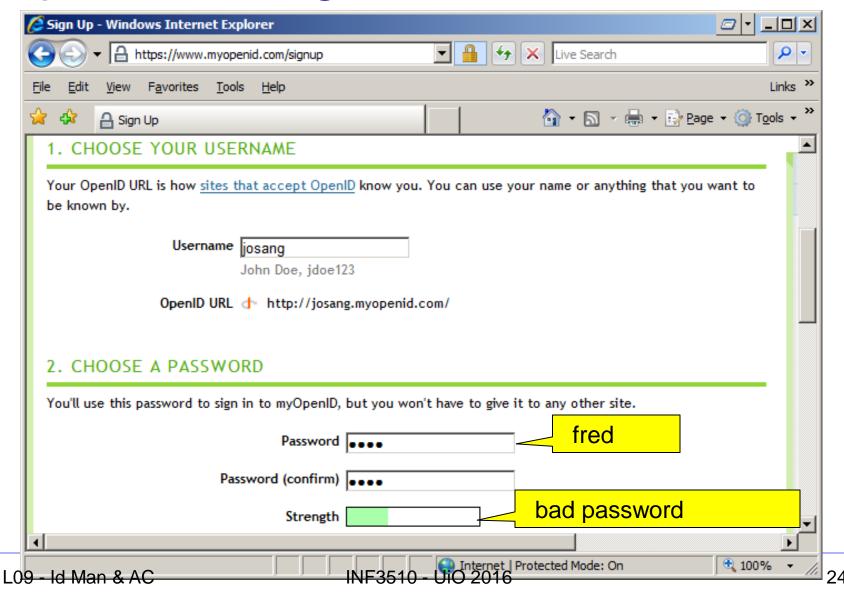
# SAML protocol profile: Browser Artefact Security token via back-channel



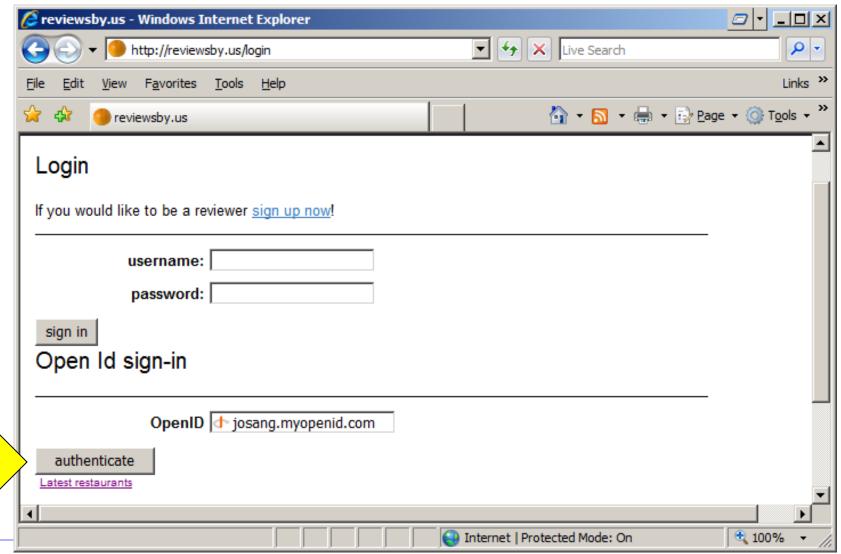
# OpenID Distributed Federation



# OpenID self registration



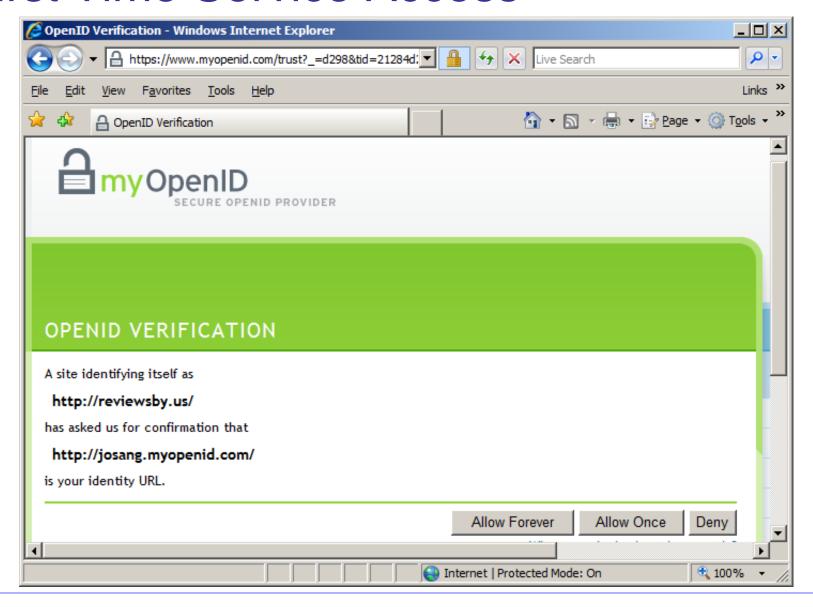
#### Service Access Without Password



L09 - Id Man & AC

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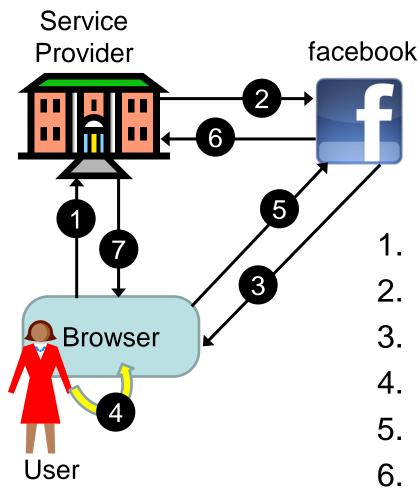
#### First Time Service Access



# **OpenID Characteristics**

- Self registration
- Anybody can be IdProvider and Server, also you
- Not all IdProviders are recognised as "authorities"
- A SP can specify which IdPs it accepts
- Not suitable for sensitive services
- Typically for services that only require low authentication assurance
- Vulnerable to multiple forms of abuse

#### Facebook Centralised Federation



Authentication with Facebook Connect

- 1. User requests service
- 2. Redirect to facebook authentication
- 3. Present facebook login form
- 4. User provides Id + credential
- 5. Credentials forwarded to facebook
- 6. Confirm authenticated user
- 7. Provide service

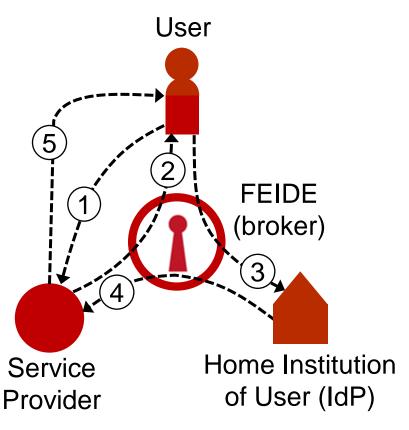


- FEIDE is a distributed federation with centralised broker for the Norwegian national education sector.
- Users register username and password with own home organisation
- Users authenticate to web-services via FEIDE's centralized login service
- The Service Provider receives user attributes from the user's Home Institution
- The Service Providers never sees the user's password/credential, it only receives user attributes that it need to know in order to provide the service.



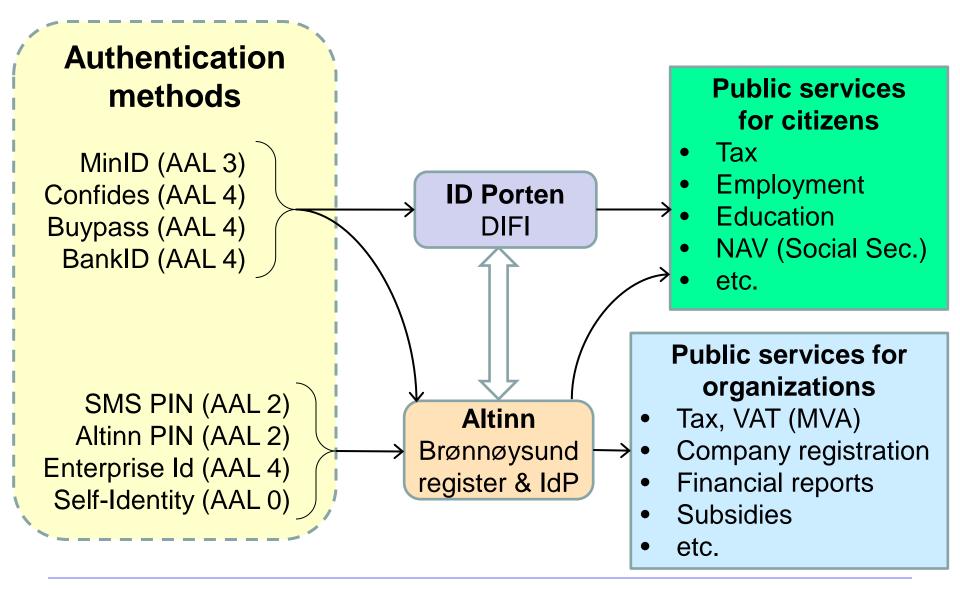
- FEIDE has formal agreements with the universities and schools before they are connected
- Home Institutions (universities and schools) are responsible for keeping user data correct and up-to-date
- Service Providers decide themselves what services their own users and other users should be able to access via FEIDE's central log-in service.





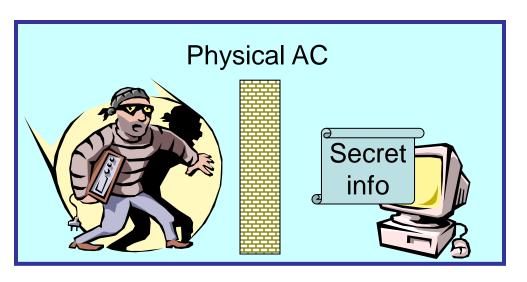
- 1. User requests access to service
- 2. Service Provider sends authentication request to FEIDE, and displays FEIDE login form to user.
- 3. User enters name and password in FEIDE login form, which are sent for validation to Home Institution of user.
- Home Institution confirms authentic user and provides user attributes to FEIDE which forwards these to SP
- 5. Service Provider analyses user attributes and provides service according to policy

#### Norw, e-Gov. Distributed Fed. with Broker

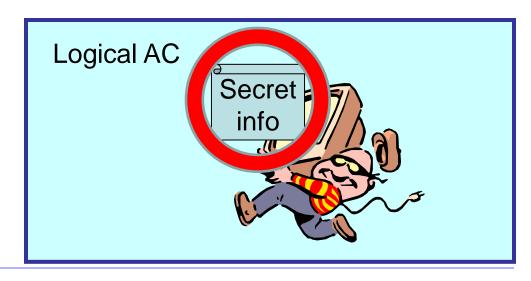


### Introduction to Logical Access Control

Physical Access Control: (not the theme today)



Logical Access Control: (this lecture)



### Basic concepts

- Access control security models:
  - How to define which subjects can access which objects with which access modes?
- Three classical approaches
  - Discretionary Access Control (DAC)
  - Mandatory access control (MAC)
  - Role-Based Access Control (RBAC)
- Advanced approach for distributed environments:
  - Attribute-Based Access Control (ABAC)
    - Generalisation of DAC, MAC and RBAC

#### Access modes

- Modes of access:
  - Authorizations specify the access permissions of subjects (users) when accessing objects (resources)
- If you are authorized to access a resource, what are you allowed to do to the resource?
  - Example: possible access permissions include
    - read observe
    - write observe and alter
    - execute neither observe nor alter
    - append alter

# DAC / MAC According to the Orange Book (TCSEC)

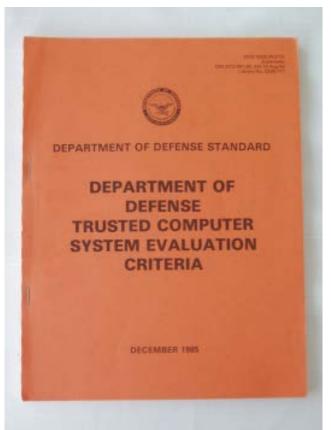
#### TCSEC (1985) specifies two AC security models

- Discretionary AC (DAC)
  - AC policy based on user identities
  - e.g. John has (r,w) access to HR-files

	HR	Sales
John	r,w	
Mary		r,w

- Mandatory AC (MAC)
  - AC policy based on security labels
  - e.g. secret clearance needed for access





Orange Book, 1985

# DAC - Discretionary Access Control

- Access authorization is specified and enforced based on the identity of the user.
- DAC is typically implemented with ACL (Access Control Lists)
- DAC is discretionary in the sense that the owner of the resource can decide at his/her discretion who is authorized
- Operating systems using DAC:
  - Windows and Linux

# DAC principles

- AC Matrix
  - General list of authorizations
  - Impractical, too many empty cells
- Access Control Lists (ACL)
  - Associated with an object
  - Represent columns from AC Matrix
  - Tells who can access the object

Columns→		Objects			
↓Rows		01	O2	O3	O4
	S1	r,w	-	Х	r
Subject names	S2	r	•	r	r,w
	S3	I	Х	ı	-
ľ	S4	r,w	Х	X	X

**AC Matrix** 

AC	lists	$\rightarrow$
	AC	<b>AC</b> lists

	01
S1	r,w
S2	r
S3	-
S4	r,w

	O2
S1	I
S2	ı
S3	X
S4	Х

	O3
S1	X
S2	r
S3	-
S4	Х

	04
S1	r
S2	r,w
S3	-
S4	Х

### **ACL** in Unix

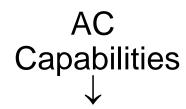
#### Each file and directory has an associated ACL

- ◆Three access operations:
  - -read: from a file
  - -write: to a file
  - -execute: a file

- Access applied to a directory:
  - read: list contents of dir
  - write: create or rename files in dir
  - execute: search directory
- Permission bits are grouped in three triples that define read,
   write, and execute access for owner, group and others.
- •A '-' indicates that the specific access right is not granted.
- •rw-r--r-- means: read and write access for the owner, read access for group, and for others (world).
- rwx----- means: read, write, and execute access for the owner, no rights for group and no rights for others

# Capabilities

- Focus on the subjects:
  - access rights stored with subjects
  - Represents rows of AC Matrix
- Must be impossible for users to create fake capabilities
- Subjects may grant own capabilities to other subjects.
   Subjects may grant the right to grant rights.
- Challenges:
  - How to check who may access a specific object?
  - How to revoke a capability?
- Similar to SAML security token



	01	O2	O3	04
S1	r,w	•	X	r

	01	O2	O3	04
S2	r	•	r	r,w

	01	O2	O3	04
S3	-	Х	-	_

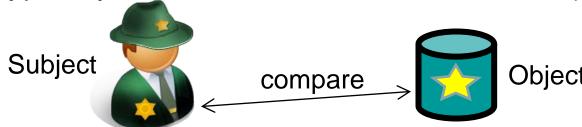
	01	02	O3	O4
S4	r,w	Х	X	X

## MAC - Mandatory Access Control

- Access authorization is specified and enforced with security labels
  - Security clearance for subjects
  - Classification levels for objects
- MAC compares subject and object labels
- MAC is mandatory in the sense that users do not control access to the resources they create.
- A system-wide set of AC policy rules for subjects and objects determine modes of access
- OS with MAC:
  - SE Linux supports MAC

# MAC principles: Labels

- Security Labels can be assigned to subjects and objects
  - Can be strictly ordered security levels, e.g. "Confidential" or "Secret"
  - Can also be partially ordered categories, e.g. {Sales-dep, HR-dep}
- Dominance relationship between labels
  - (  $L_A \ge L_B$  ) means that label  $L_A$  dominates label  $L_B$
- Object labels are assigned according to sensitivity
- Subject labels are determined by security clearance
- Access control decisions are made by comparing the subject label with the object label according to specific model
- MAC is typically based on Bell-LaPadula model (see later)



#### Bell-LaPadula: The classical MAC model

#### SS-property (Simple Security): No Read Up

- A subject should not be able to read files with a higher label than its own label, because otherwise it could cause unauthorized disclosure of sensitive information.
- So you should only be able to read documents with an equal or lower label as your security clearance level.

#### \*-Property (Star Property): No Write Down

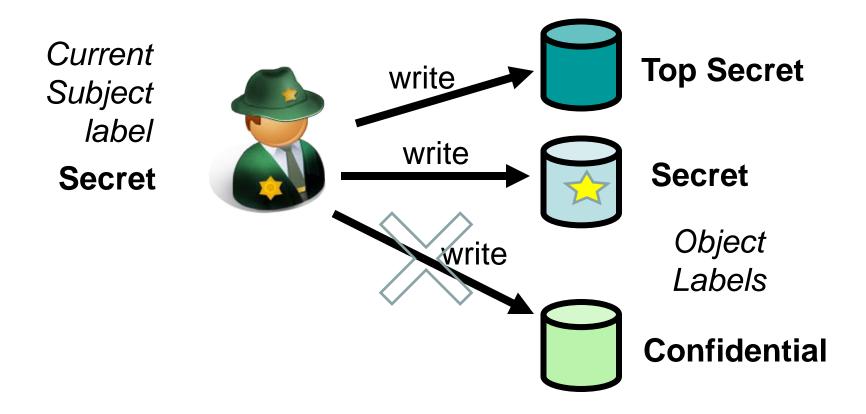
- Subjects working on information/tasks at a given level should not be allowed to write to a lower level, because otherwise it could create unauthorized information flow.
- So you should only be able write to files with an equal or higher label as your security clearance level.

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# Bell-LaPadula (MAC model) SS-Property: No Read Up

**Top Secret** Current read Object Subject Labels Label read Secret **Secret** read **Confidential**  Diagram\_\_\_

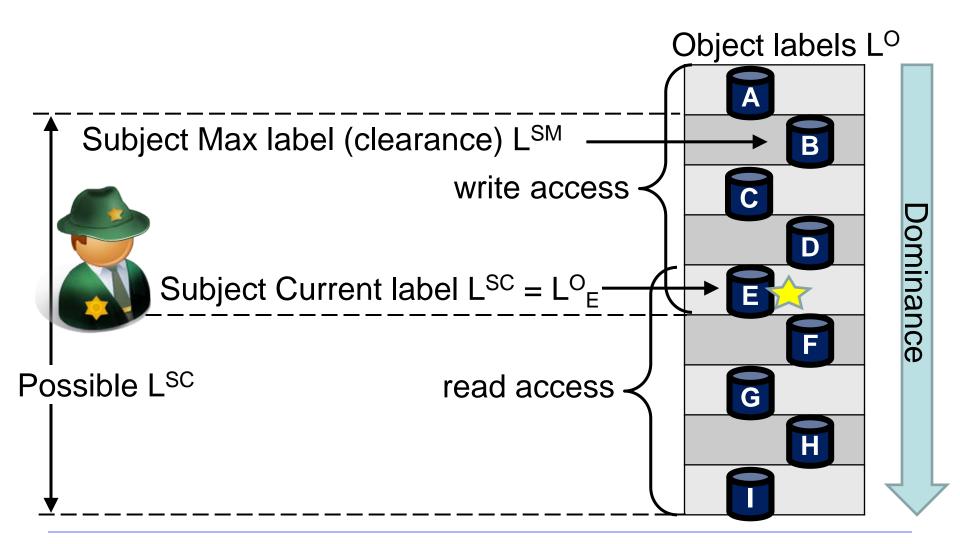
# Bell-LaPadula (MAC model) \*-Property: No Write Down



## Labels in Bell La Padula

- Users have a clearance level L<sup>SM</sup> (Subject Max level)
- Users log on with a current clearance level  $L^{SC}$  (Subject Current level) where  $L^{SC} \leq L^{SM}$
- Objects have a sensitivity level L<sup>O</sup> (Object)
- SS-property allows read access when L<sup>SC</sup> ≥ L<sup>O</sup>
- \*-property allows write access when L<sup>SC</sup> ≤ L<sup>O</sup>

# Bell-LaPadula label relationships



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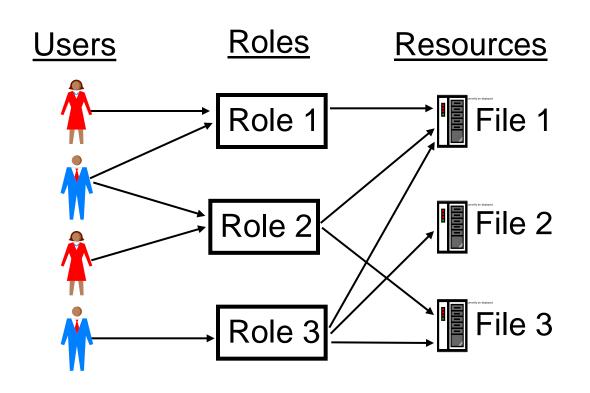
## Combined MAC & DAC

- Combining access control approaches:
  - A combination of mandatory and discretionary access control approaches is often used
    - MAC is applied first,
    - DAC applied second after positive MAC
    - Access granted only if both MAC and DAC positive
  - Combined MAC/DAC ensures that
    - no owner can make sensitive information available to unauthorized users, and
    - 'need to know' can be applied to limit access that would otherwise be granted under mandatory rules

## RBAC: Role Based Access Control

- A user has access to an object based on the assigned role.
- Roles are defined based on job functions.
- Permissions are defined based on job authority and responsibilities within a job function.
- Operations on an object are invocated based on the permissions.
- The object is concerned with the user's role and not the user.

## **RBAC Flexibility**



User's change frequently, roles don't

RBAC can be configured to do MAC and/or DAC

# RBAC Privilege Principles

- Roles are engineered based on the principle of least privilege.
- A role contains the minimum amount of permissions to instantiate an object.
- A user is assigned to a role that allows her to perform only what's required for that role.
- All users with the same role have the same permissions.

## ABAC and XACML

#### ABAC = Attribute Based Access Control

- ABAC specifies access authorizations and approves access through policies combined with attributes. The policy rules can apply to any type of attributes (user attributes, resource attribute, context attributed etc.).
- XACML used to express ABAC attributes and policies.

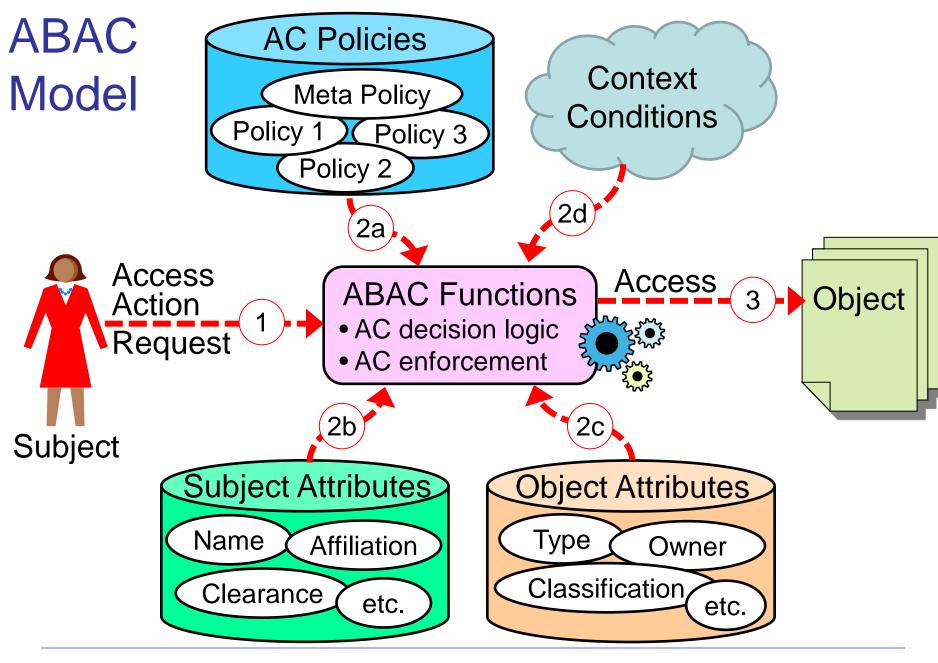
#### **XACML** = eXtensible Access Control Markup Language

- The XACML standard defines a language for expressing access control attributes and policies implemented in XML, and a processing model describing how to evaluate access requests according to the rules defined in policies.
- XACML attributes are typically structured in ontologies

### **Attribute Based Access Control**

- ABAC makes AC decisions based on Boolean conditions on attribute values.
- Subject, Object, Context, and Action consist of attributes
  - Subject attributes could be: Name, Sex, DOB, Role, etc.
  - Each attributes has a value, e.g.:
  - (Name (subject) = Alice), (Sex(subject) = F), (Role(subject) = HR-staff), (AccessType(action) = {read, write}),
     (Owner(object) = HR), (Type(object) = salary)
- The AC logic analyses all (attribute = value) tuples that are required by the relevant policy.
  - E.g. permit if:

```
[Role(subject) = HR-staff) and (AccessType(action) = read) and (Owner(object) = HR) ] and (Time(query) = office-hours) ]
```



## Global Consistence

- ABAC systems require an XML terminology to express all possible attributes and their values,
- Must be consistent across the entire domain,
  - e.g. the attribute Role and all its possible values, e.g.
     (Role(subject) = HR-staff), must be known and interpreted by all systems in the AC security domain.
- Requires standardization:
  - e.g. for access to medical journals, medical terms
     must be interpreted in a consistent way by all systems
  - current international work on XML of medical terms
- Consistent interpretation of attributes and values is a major challenge for implementing ABAC.

## ABAC: + and -

#### On the positive side:

- •ABAC is much more flexible than DAC, MAC or RBAC
  - DAC, MAC and RBAC can be implemented with ABAC
- Can use any type of access policies combined with an unlimited number of attributes
- Suitable for access control in distributed environments
  - e.g. national e-health networks

#### On the negative side:

- •Requires defining business concepts in terms of XML and ontologies which is much more complex than what is required in traditional DAC, MAC or RBAC systems.
- Political alignment and legal agreements required for ABAC in distributed environments

## Meta-policies i.c.o. inconsistent policies

- Sub-domain authorities defined their own policies
- Potential for conflicting policies
  - E.g. two policies dictate different access decisions
- Meta-policy rules needed in case the ABAC logic detects policy rules that lead to opposite decisions
- Meta-policy takes priority over all other policies, e.g.
  - Meta-Policy Deny Overrides: If one policy denies access, but another policy approves access, then access is denied.
     This is a conservative meta-policy.
  - Meta-Policy Approve Overrides: If one policy denies access, but another policy approves access, then access is approved.
  - This is a lenient meta-policy.

## End of lecture