

INF3510 Information Security

Lecture 10: Communications Security

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Outline

- Network security concepts
 - Communication security
 - Perimeter security
- Protocol architecture and security services
- Example security protocols
 - Transport Layer Security (TLS)
 - IP Layer Security (IPSec)
- VPN – Virtual Private Network

Network Security Concepts

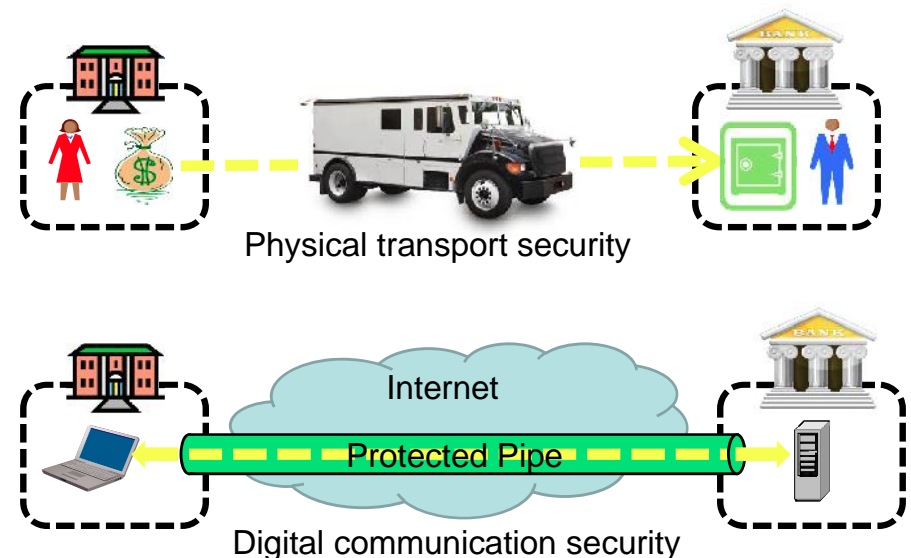
Assumes that each organisation owns a network

- Wants to protect own local network
- Wants to protect communication with other networks

Network Security: two main areas

- **Communication Security:** Protection of data transmitted across networks between organisations and end users
 - Topic for this lecture
- **Perimeter Security:** Protection of an organization's network from unauthorized access
 - Topic for next lecture

Communication Security Analogy



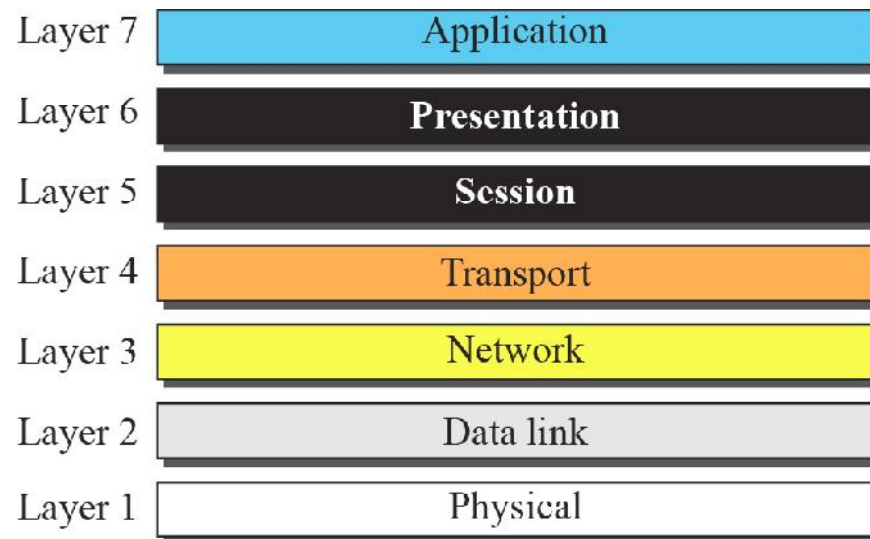
Communication Protocol Architecture

- Layered structure of hardware and software that supports the exchange of data between systems
- Each protocol consists of a set of rules for exchanging messages, i.e. “the protocol”.
- Two standards:
 - OSI Reference model
 - Never lived up to early promises
 - TCP/IP protocol suite
 - Most widely used

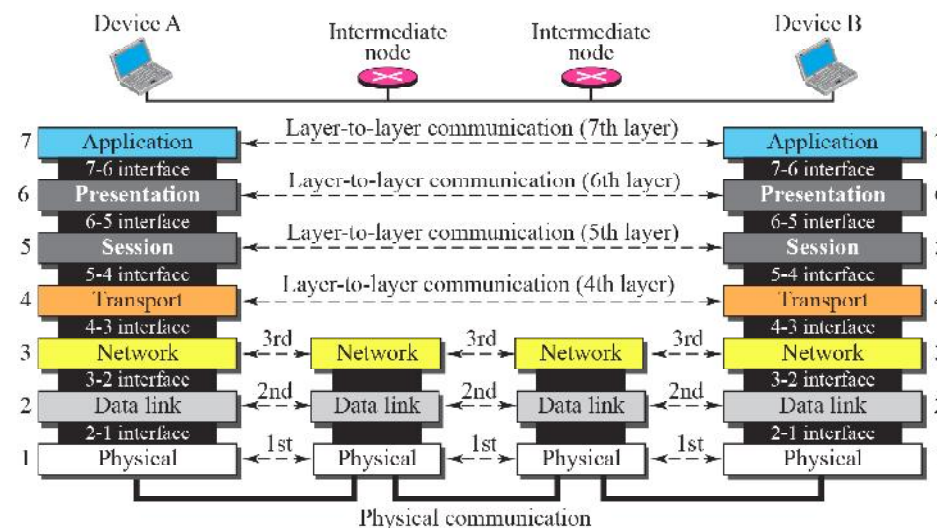
OSI – Open Systems Interconnection

- Developed by the International Organization for Standardization (ISO)
- A layer model of 7 layers
- Each layer performs a subset of the required communication functions
- Each layer relies on the next lower layer to perform more primitive functions
- Each layer provides services to the next higher layer
- Changes in one layer should not require changes in other layers

The OSI Protocol Stack



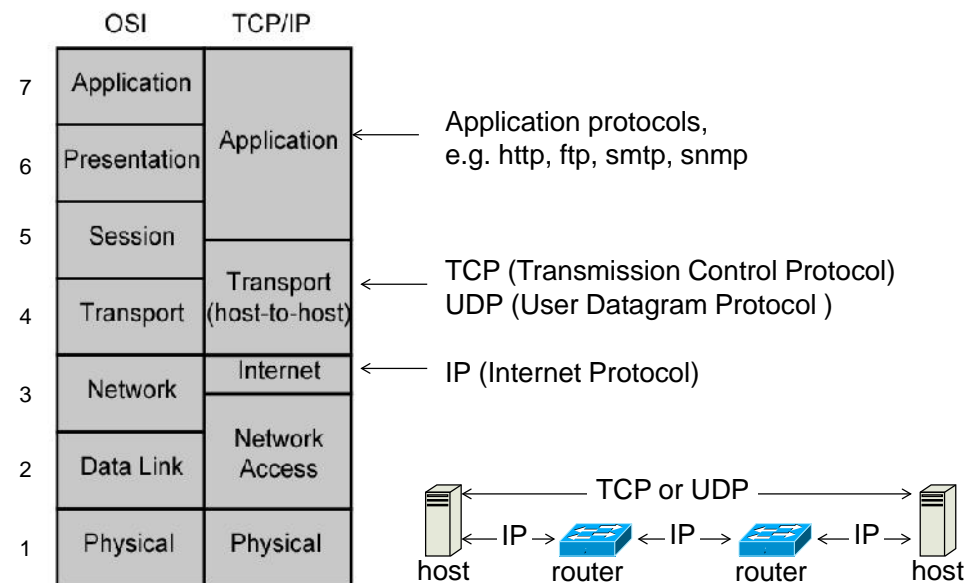
Communication across OSI



TCP/IP Protocol Architecture

- Developed by the US Defense Advanced Research Project Agency (DARPA) for its packet switched network (ARPANET)
- Used by the global Internet
- No official model, but it's a working one.
 - Application layer
 - Host to host or transport layer
 - Internet layer
 - Network access layer
 - Physical layer

OSI model vs. TCP/IP model (The Internet)



Security Protocols

- Many different security protocols have been specified and implemented for different purposes
 - Authentication, integrity, confidentiality
 - Key establishment/exchange
 - E-Voting
 - Secret sharing
 - etc.
- Protocols are surprisingly difficult to get right!
 - Many vulnerabilities are discovered years later
 - ... some are never discovered (or maybe only by the attackers)

Security Protocols Overview

- This lecture discusses the operation of two network-related protocols that are in common use.
 - **Transport Layer Security (TLS):**
Used extensively on the web and is often referred to in privacy policies as a means of providing confidential web connections.
 - **IP Security (IPSec):**
Provides security services at the IP level and is used to provide Virtual Private Network (VPN) services.

Transport Layer Security

TLS/SSL

SSL/TLS: History

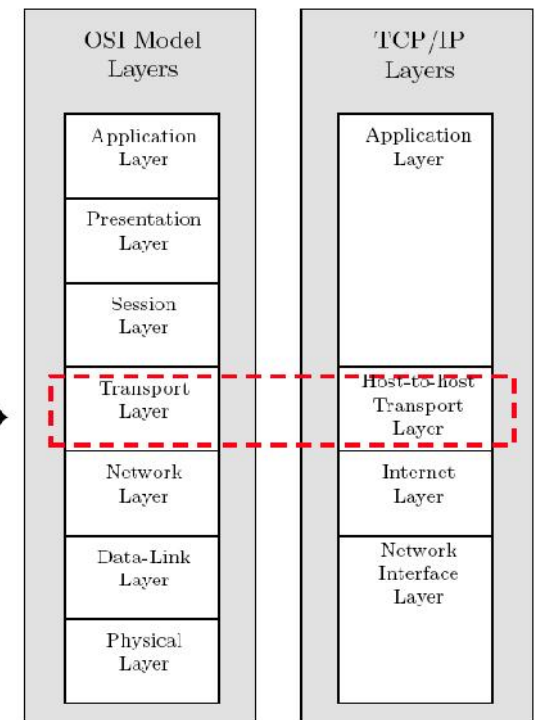
- 1994: Netscape Communications developed the network authentication protocol Secure Sockets Layer, SSLv2.
 - Badly broken
- 1995: Netscape release their own improvements SSLv3.
 - Widely used for many years.
- 1996: SSLv3 was submitted to the IETF as an Internet draft, and an IETF working group was formed to develop a recommendation.
- In January 1999, [RFC 2246](#) was issued by the IETF, Transport Layer Security Protocol: TLS 1.0
 - Similar to, but incompatible with SSLv3
 - Currently TLS 1.2 (2008) (allows backwards compatibility with SSL)
 - Draft TLS 1.3 (2016) (totally bans SSL)
 - Firefox browser enabled TLS 1.3 by default in February 2017¹

TLS: Overview

- TLS is a cryptographic services protocol based on the Browser PKI, and is commonly used on the Internet.
 - Each server has a server certificate and private key installed
 - Allows browsers to establish secure sessions with web servers.
- Port 443 is reserved for HTTP over TLS/SSL and the protocol https is used with this port.
 - `http://www.xxx.com` implies using standard HTTP using port 80.
 - `https://www.xxx.com` implies HTTP over TLS/SSL with port 443.

TLS: Layer 4 Security

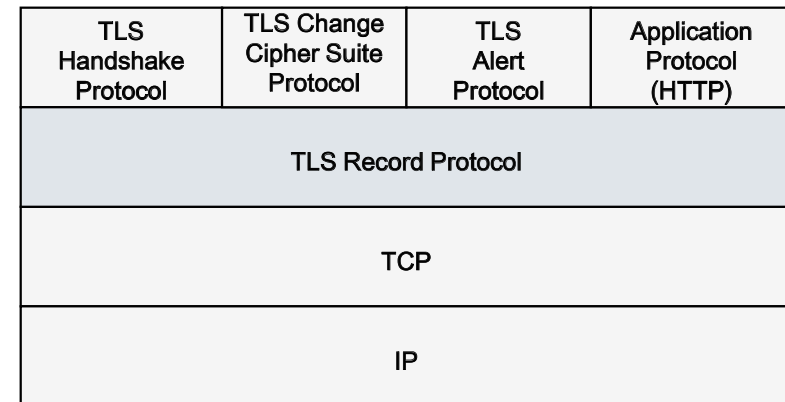
TLS operates at Layer 4



TLS: Architecture Overview

- Designed to provide secure reliable end-to-end services over TCP.
- Consists of 3 higher level protocols:
 - TLS Handshake Protocol
 - TLS Alert Protocol
 - TLS Change Cipher Spec Protocol
- The TLS Record Protocol provides the practical encryption and integrity services to various application protocols.

TLS: Protocol Stack



TLS: Handshake Protocol

- The handshake protocol
 - Negotiates the encryption to be used
 - Establishes a shared session key
 - Authenticates the server
 - Authenticates the client (optional)
 - Completes the session establishment
- After the handshake, application data is transmitted securely
- Several variations of the handshake exist
 - RSA variants
 - Diffie-Hellman variants

TLS: Handshake Four phases

- Phase 1: Initiates the logical connection and establishes its security capabilities
- Phases 2 and 3: Performs key exchange. The messages and message content used in this phase depends on the handshake variant negotiated in phase 1.
- Phase 4: Completes the setting up of a secure connection.

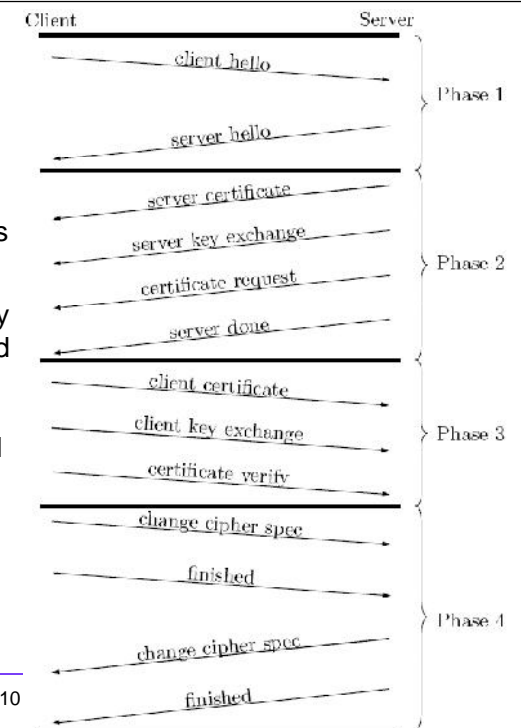
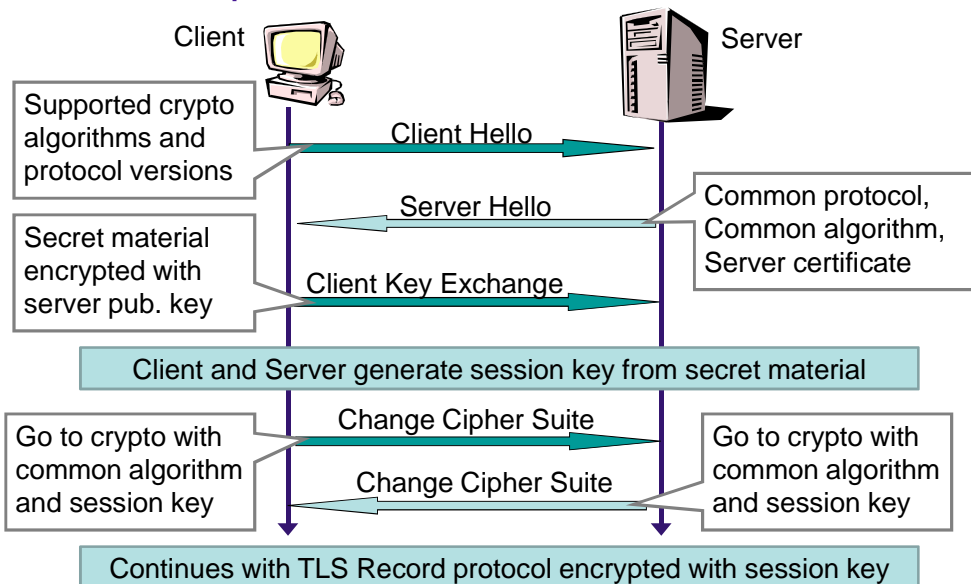


Diagram TLS: Simplified RSA-based Handshake



TLS: Elements of Handshake

- **Client hello**
 - Advertises available cipher suites (e.g. RSA, AES, SHA256)
- **Server hello**
 - Returns the selected cipher suite
 - Server adapts to client capabilities
- **RSA and Server Certificate**
 - X.509 digital certificate sent to client, assumes RSA algorithm
 - Client verifies the certificate including that the certificate signer is in its acceptable Certificate Authority (CA) list. Now the client has the server's certified public key.
- **RSA and Client Certificate**
 - Optionally, the client can send its X.509 certificate to server, in order to provide mutual authentication, assumes RSA algorithm
- **Anonymous Diffie-Hellman**
 - Optionally, the client and server can establish session key using the Diffie-Hellman algorithm

TLS: Record Protocol Overview

- Provides two services for SSL connections.
 - **Message Confidentiality:**
 - Ensure that the message contents cannot be read in transit.
 - The Handshake Protocol establishes a symmetric key used to encrypt SSL payloads.
 - **Message Integrity:**
 - Ensure that the receiver can detect if a message is modified in transmission.
 - The Handshake Protocol establishes a shared secret key used to construct a MAC.

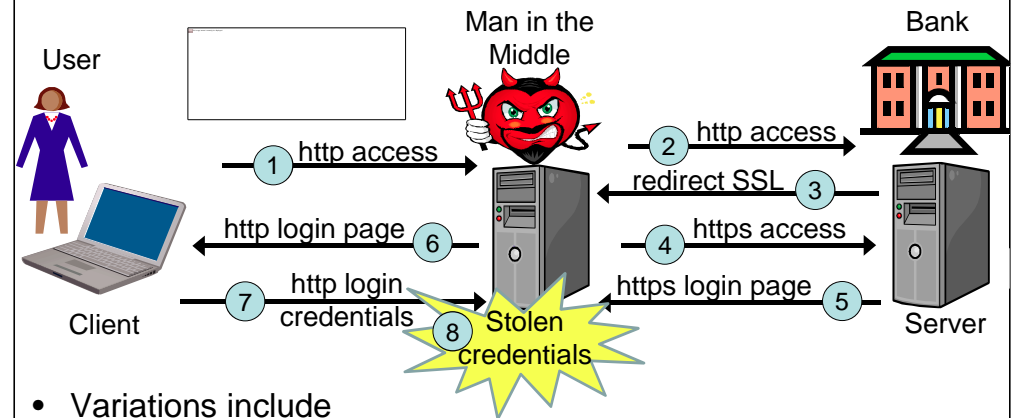
TLS: Record Protocol Operation

- **Fragmentation:**
 - Each application layer message is fragmented into blocks of 214 bytes or less.
- **Compression:**
 - Optionally applied.
 - SSL v3 & TLS – default compression algorithm is null
- **Add MAC:**
 - Calculates a MAC over the compressed data using a MAC secret from the connection state.
- **Encrypt:**
 - Compressed data plus MAC are encrypted with symmetric cipher.
 - Permitted ciphers include AES, IDEA, DES, 3DES, RC4
 - For block ciphers, padding is applied after the MAC to make a multiple of the cipher's block size.

SSL/TLS Challenges

- Higher layers should not be overly reliant on SSL/TLS.
- Many vulnerabilities exist for SSL/TLS.
 - People are easily tricked
 - Changing between http and https causes vulnerability to SSL stripping attacks
 - SSL/TLS only as secure as the cryptographic algorithms used in handshake protocol: hashing, symmetric and asymmetric crypto.
- Relies on Browser PKI which has many security issues
 - Fake server certificates difficult to detect
 - Fake root server certificates can be embedded in platform, see e.g. Lenovo Komodia advare scam

SSL Stripping Attack

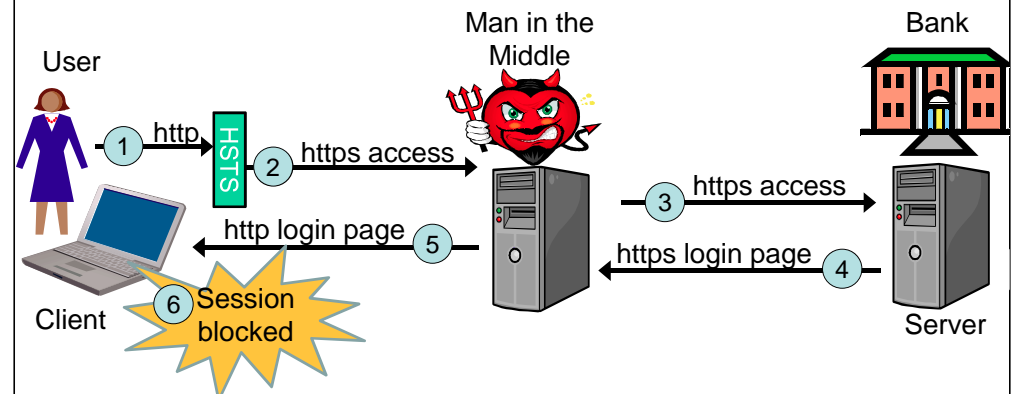


- Variations include
 - MitM server can connect to client over https in msg (6) with server certificate that has similar domain name as real server.
 - Attacker can leave the connection after stealing credentials, then the client connects directly to real server with https

HSTS – HTTP Strict Transport Security Preventing SSL Stripping

- A secure server can instruct browsers to only use https
- When requesting website that uses HSTS, the browser automatically forces connect with https.
- Users are not able to override policy
- Two ways of specifying HSTS websites
 - List of HSTS websites can be preloaded into browsers
 - HSTS policy initially specified over a https connection
 - HSTS policy can be changed over a https connection
- Disadvantages
 - HSTS websites can not use both http and https
 - Difficult for a website to stop using https
 - Can cause denial of service, e.g. no fallback to http in case of expired server certificate

Preventing SSL Stripping with HSTS



- Limitation of HSTS:
 - Requires first visit to secure website to set HSTS policy in browser
- Can be solved by browser having preloaded list of HSTS websites
- Browsers would be vulnerable if attacker could delete HSTS cache

Server Authentication Modalities

Syntactic entity authentication:

- Verification that the identity of the remote entity is as claimed.
- Does not provide any meaningful security because of indifference to the owner entity of authenticated identity.

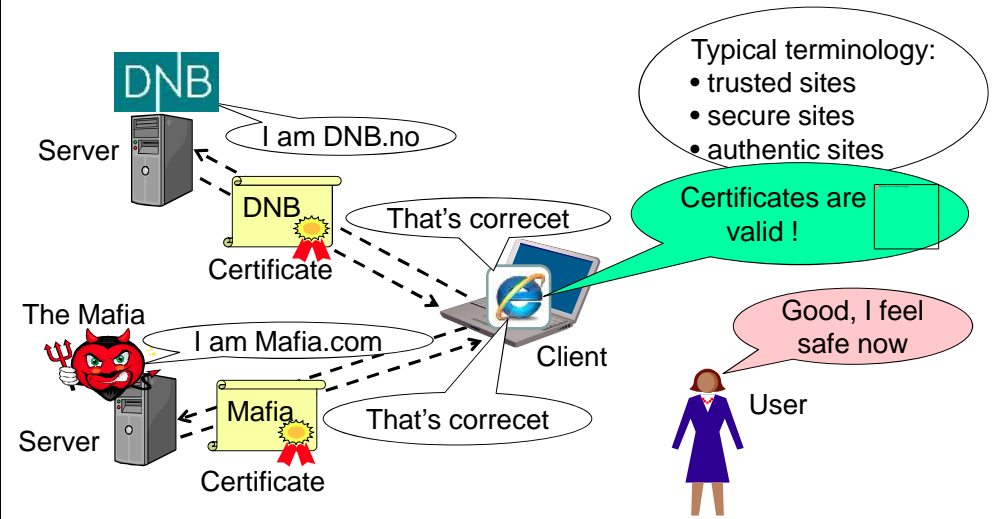
Semantic entity authentication:

- Verification that the identity of the remote entity is as claimed, combined with a policy for authenticated entities.
- Automated control of policy, e.g. with white-list of domains.

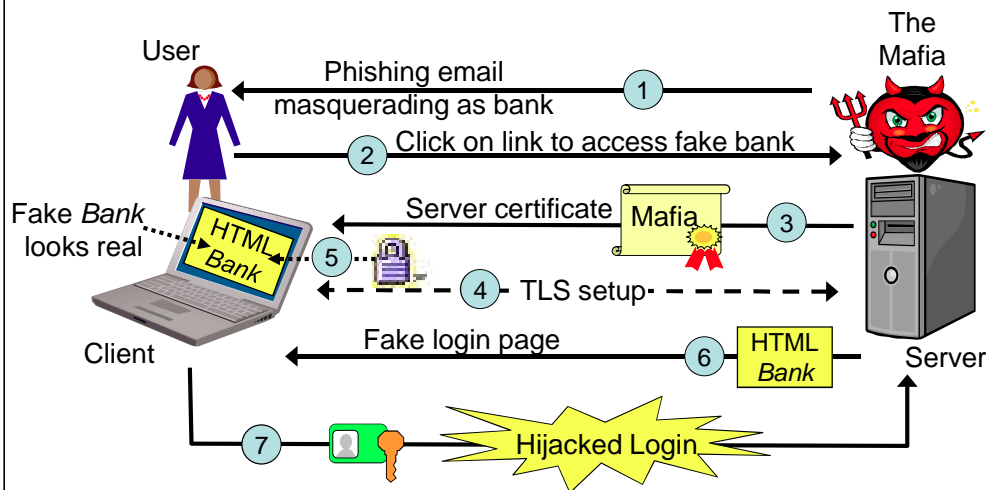
Cognitive entity authentication:

- Verification by a cognitive entity (human) that the identity of the remote entity is as claimed, and a conscious decision that the identity is acceptable and as expected.
- Manual control of policy, e.g. by inspecting certificates.

Meaningless Server Authentication when only applying syntactic authentication



Phishing and failed authentication



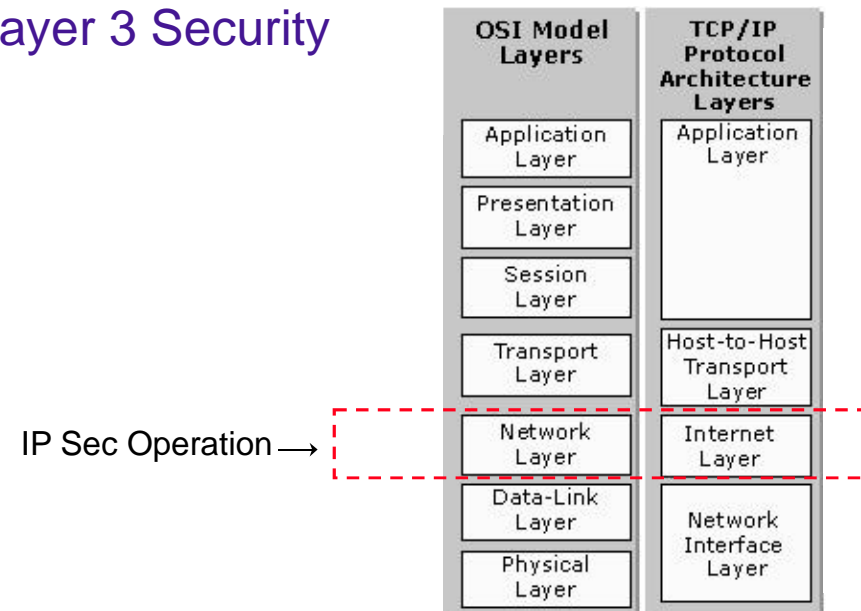
IP Layer Security

IPSec & Virtual Private Networks

IPSec: Introduction

- Internet Protocol security (IPSec) is standard for secure communications over Internet Protocol (IP) networks, through the use of cryptographic security services.
- Uses encryption, authentication and key management algorithms
- Based on an end-to-end security model at the IP level
- Provides a security architecture for both IPv4 and IPv6
 - Mandatory for IPv6
 - Optional for IPv4
- Requires operating system support, not application support.

Layer 3 Security



IPSec: Security Services

- **Message Confidentiality.**
 - Protects against unauthorized data disclosure.
 - Accomplished by the use of encryption mechanisms.
- **Message Integrity.**
 - IPSec can determine if data has been changed (intentionally or unintentionally) during transit.
 - Integrity of data can be assured by using a MAC.
- **Traffic Analysis Protection.**
 - A person monitoring network traffic cannot know which parties are communicating, how often, or how much data is being sent.
 - Provided by concealing IP datagram details such as source and destination address.

IPSec: Security Services

- **Message Replay Protection.**
 - The same data is not delivered multiple times, and data is not delivered grossly out of order.
 - However, IPsec does not ensure that data is delivered in the exact order in which it is sent.
- **Peer Authentication.**
 - Each IPsec endpoint confirms the identity of the other IPsec endpoint with which it wishes to communicate.
 - Ensures that network traffic is being sent from the expected host.
- **Network Access Control.**
 - Filtering can ensure users only have access to certain network resources and can only use certain types of network traffic.

IPSec: Common Architectures

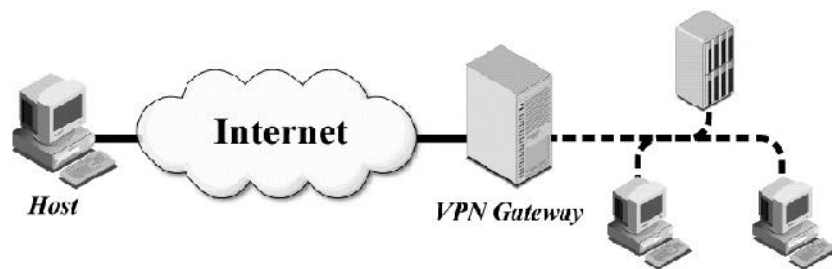
- Gateway-to-Gateway Architecture
- Host-to-Gateway Architecture
- Host-to-Host Architecture

IPSec: Gateway-to-Gateway Architecture



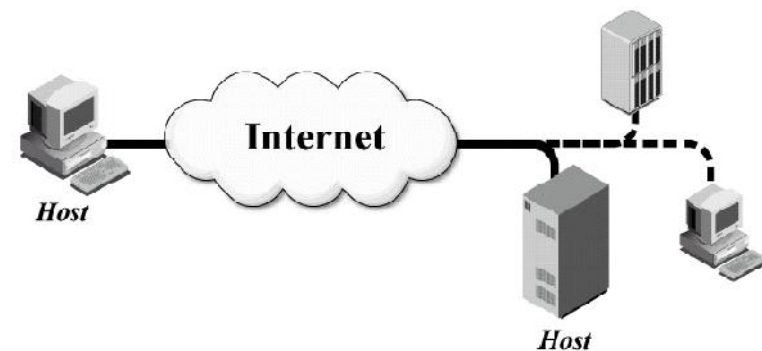
Source: NIST Special Publication 800-77

IPSec: Host-to-Gateway Architecture



Source: NIST Special Publication 800-77

IPSec: Host-to-Host Architecture



Source: NIST Special Publication 800-77

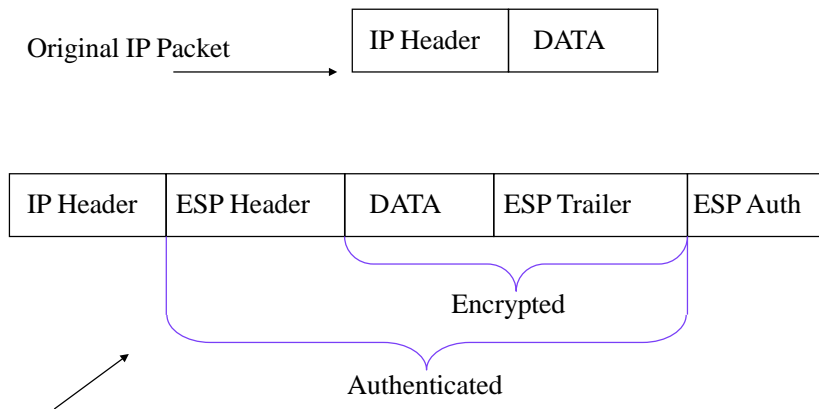
IPSec: Protocols Types

- Encapsulating Security Payload (ESP)
 - Confidentiality, authentication, integrity and replay protection
- Authentication Header (AH)
 - Authentication, integrity and replay protection. However there is no confidentiality
- Internet Key Exchange (IKE)
 - negotiate, create, and manage security associations
- A connection consists of two SA (Security Associations)
 - One SA for each directions
 - Each SA is described by a set of parameters

IPSec: Modes of operation

- Each protocol (ESP or AH) can operate in transport or tunnel mode.
- **Transport mode:**
 - Operates primarily on the payload (data) of the original packet.
 - Generally only used in host-to-host architectures.
- **Tunnel mode:**
 - Original packet encapsulated into a new one, payload is original packet.
 - Typical use is gateway-to-gateway and host-to-gateway architectures.

Transport Mode ESP

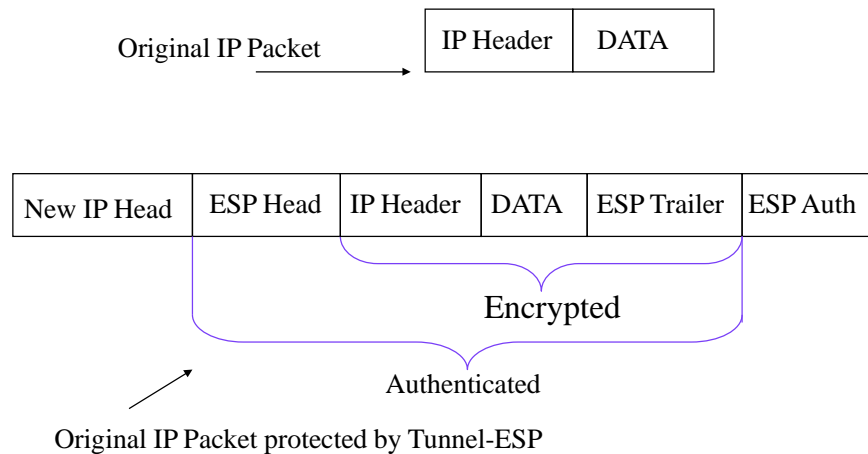


Original IP Packet protected by Transport-ESP

IPSec - ESP in Transport Mode: Outbound Packet Processing

- The data after the original IP header is padded by adding an ESP trailer and the result is then encrypted using the symmetric cipher and key in the SA.
- An ESP header is prepended.
- If an SA uses the authentication service, an ESP MAC is calculated over the data prepared so far and appended.
- The original IP header is prepended.
- However, some fields in the original IP header must be changed. For example,
 - Protocol field changes from TCP to ESP.
 - Total Length field must be changed to reflect the addition of the AH header.
 - Checksums must be recalculated.

Tunnel Mode ESP



IPSec - ESP in Tunnel Mode: Outbound Packet Processing

- The entire original packet is padded by adding an ESP trailer and the result is then encrypted using the symmetric cipher and key agreed in the SA.
- An ESP header is prepended.
- If an SA uses the authentication service, an ESP MAC is calculated over the data prepared so far and appended.
- A new 'outer' IP header is prepended.
 - The 'inner' IP header of the original IP packet carries the ultimate source and destination addresses.
 - The 'outer' IP header may contain distinct IP addresses such as addresses of security gateways.
 - The 'outer' IP header Protocol field is set to ESP.

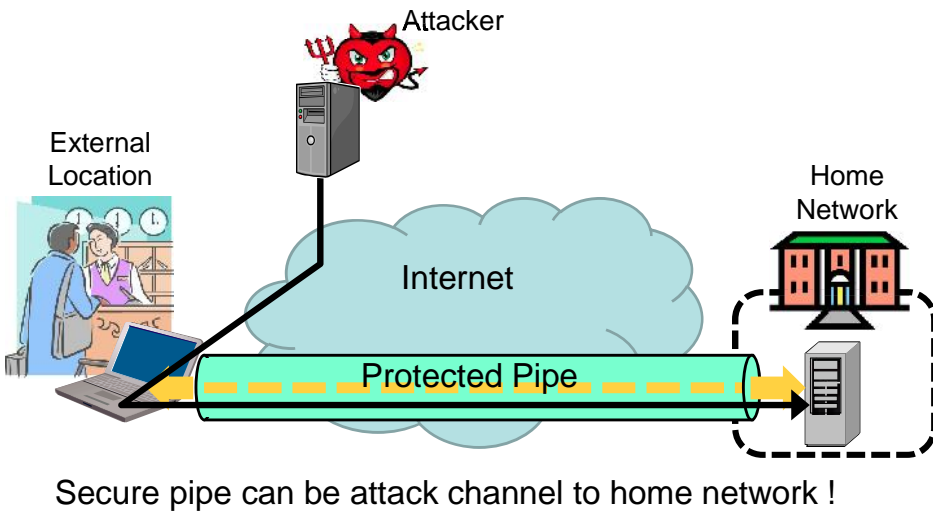
Security Associations

- A security association (SA) contains info needed by an IPSec endpoint to support one end of an IPSec connection.
- Can include cryptographic keys and algorithms, key lifetimes, security parameter index (SPI), and security protocol identifier (ESP or AH).
- The SPI is included in the IPSec header to associate a packet with the appropriate SA.
- Security Associations are simplex
 - need one for each direction of connection
 - stored in a security association database (SAD).
- Key exchange is largely automated after initial manual configuration by administrator prior to connection setup.
- (See ISAKMP, IKE, Oakley, Skeme and SAs)

Risks of using IPSec for VPN

- IPSec typically used for VPN (Virtual Private Networks)
- A VPN client at external location may be connected to the Internet (e.g. from hotel room or café) while at the same time being connected to home network via VPN.
 - VPN gives direct access to resources in home network.
- Internet access from external location may give high exposure to cyber threats
 - No network firewall, no network IDS
- Attacks against the VPN client at external location can directly access the home network through VPN tunnel

Risk of using VPN

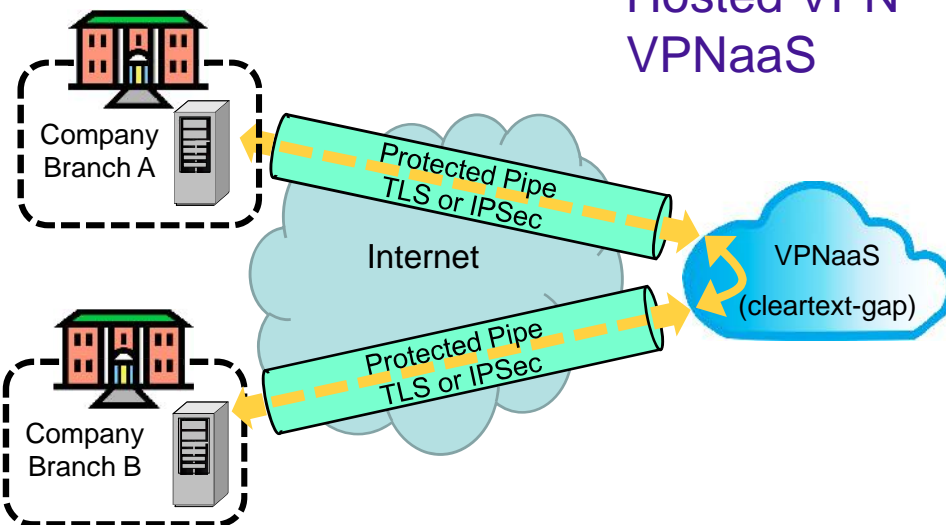


Cloud VPN



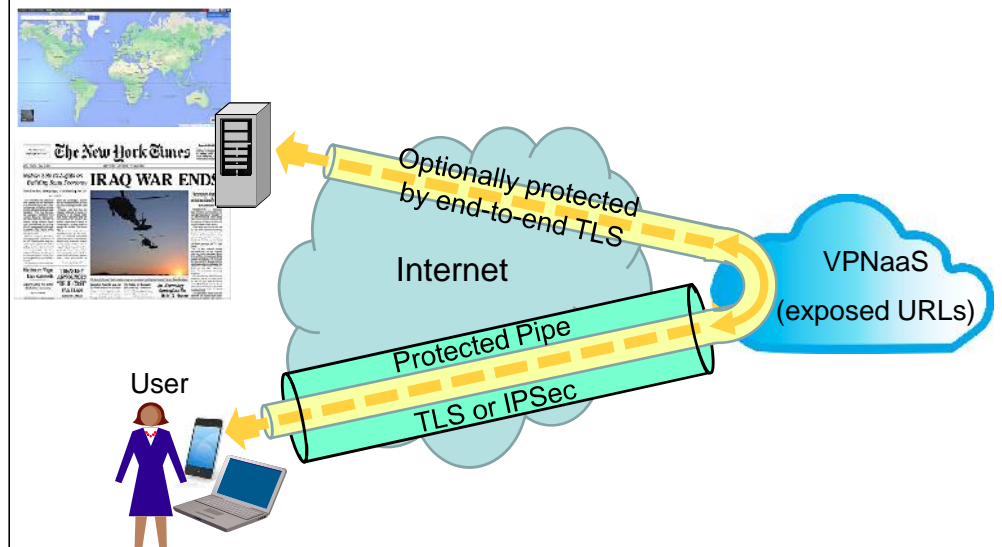
- A cloud-based infrastructure for VPN.
- A.k.a.:
 - Hosted VPN
 - VPNaaS (Virtual Private Network as a Service)
- Cloud VPNs provide security and globally accessible VPN service access without the need for any VPN infrastructure on the user's end.
- The user connects to the cloud VPN through the provider's website or a desktop/mobile app.
- The pricing of cloud VPN is based on pay-per-usage or a flat-fee subscription.
- Disadvantages /risks
 - Cleartext-gap at the VPN provider
 - VPN provider knows Internet usage profile
 - Malicious VPN service?

Cloud VPN Hosted VPN VPNaaS



Internet services

VPN Browsing – via VPN Proxy



Tor – The Onion Router



Image courtesy indymedia.de

- An anonymizing routing protocol
- Originally sponsored by the US Naval Research Laboratory
- From 2004 to 2006 was supported by EFF
- Since 2006 independent nonprofit organisation
- Creates a multi-hop proxy circuit through the Internet from client to destination.
- Each hop “wraps” another encryption layer with symmetric keys, thereby hiding the next destination.
- No cleartext-gap, except at the exit-node.
- No node knows end-to-end client-server association
- Full technical details: <http://www.torproject.org/tor-design.pdf>

How Tor Works: 1

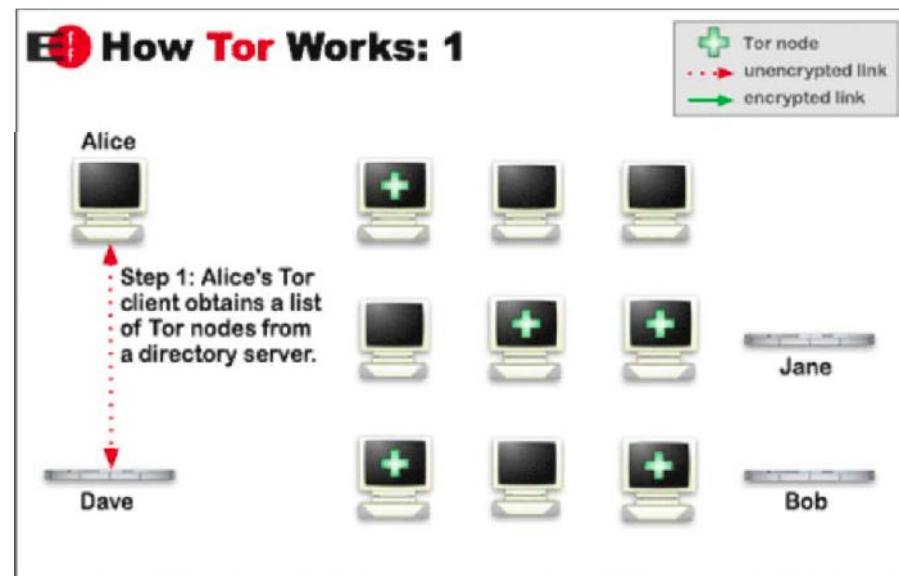


Image courtesy torproject.org

How Tor Works: 2

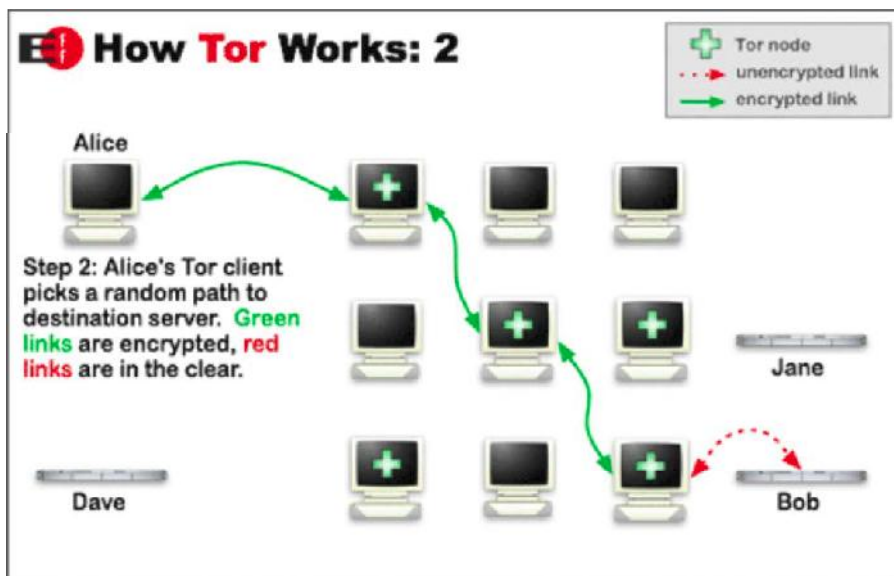


Image courtesy torproject.org

How Tor Works: 3

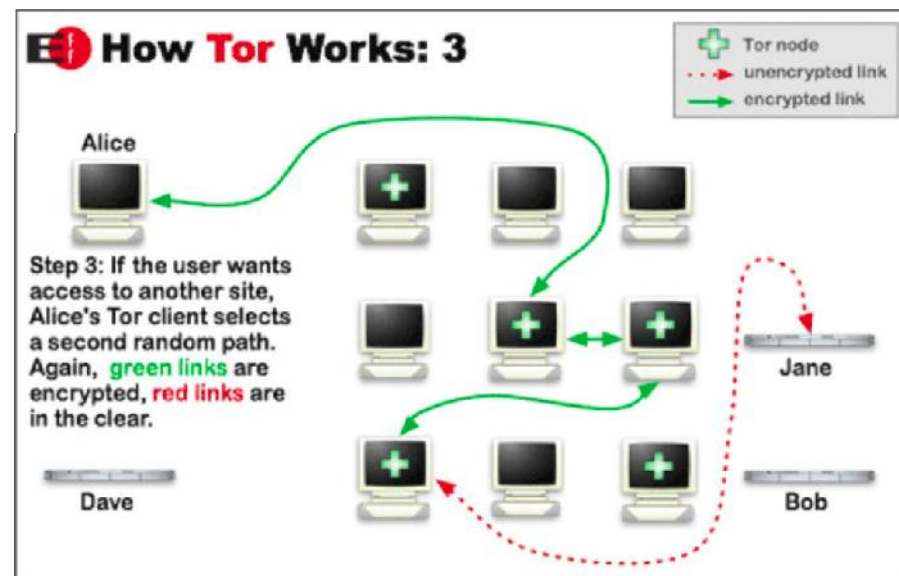


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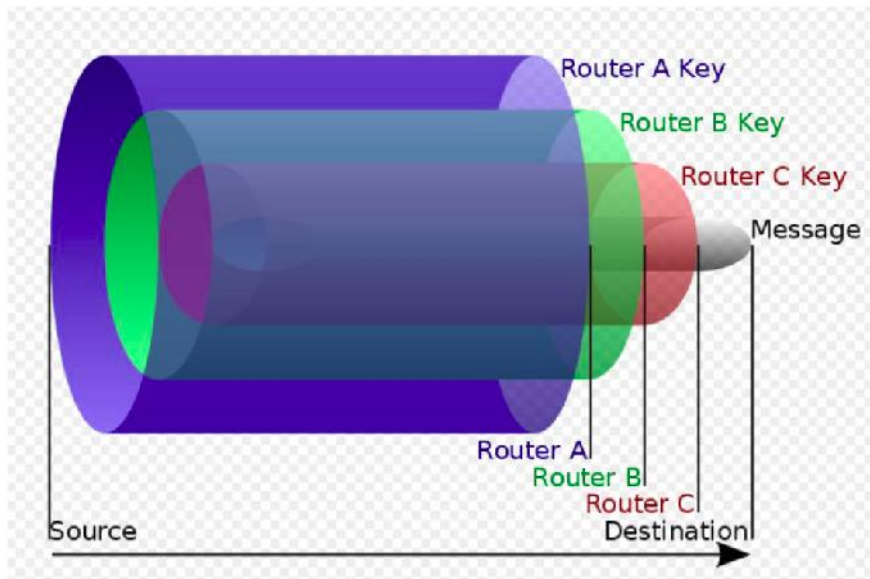


Diagram courtesy Wikimedia Commons

End of lecture