INF3510 Information Security

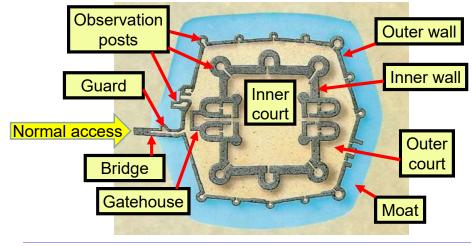
Lecture 11: Network Perimeter Security



L11: Perimeter Security

Nils Gruschka
University of Oslo
Spring 2018

Perimeter security analogy Medieval Castle Defences



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Outline

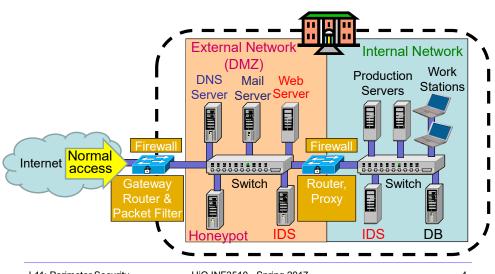
- Firewalls
 - Routers
 - Proxies
 - Architectures
- Intrusion Detection Systems
 - Host-based
 - Network based
 - Dealing with false alarms
- Wireless LAN Access Control
 - Evolution & history
 - WPA2: Robust Security Network architecture (RNS)

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Defending local networks Network Perimeter Security



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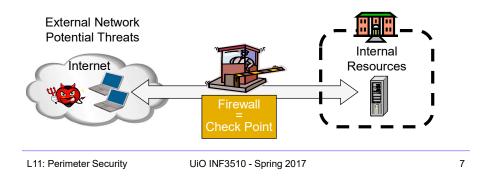
Firewalls

Firewalls: Overview 1

- If the risk of having a connection to the Internet is unacceptable, the most effective way of treating the risk is to avoid the risk altogether and disconnect completely.
- If disconnection from the Internet is not practical, then firewalls may provide an effective level of protection that can reduce the risk to an acceptable level.
- Firewalls are often the first line of defence against external attacks, but should not be the only defence.
- A firewall's purpose is to prevent unauthorized access to or from a private network.

Network perimeter security method: Firewalls

- A firewall is a check point that protects the internal networks against attack from outside networks
- The check point decides which traffic can pass in & out based on rules



Firewalls: Overview 2

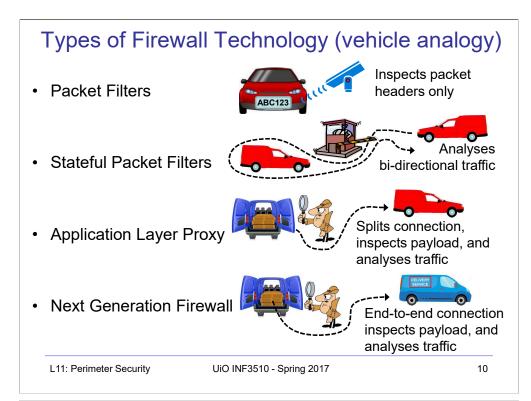
- All traffic entering or leaving must pass through firewall
- The network owner must define criteria for what is (un)authorized
- The effectiveness of firewalls depends on specifying authorized traffic in terms of rules
 - The rules defines what to let pass through;
 - The rules defines what to block.
- Firewalls must be effectively administered, updated with the latest patches and monitored.

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• Firewalls can be implemented in both hardware and software, or a combination of both.

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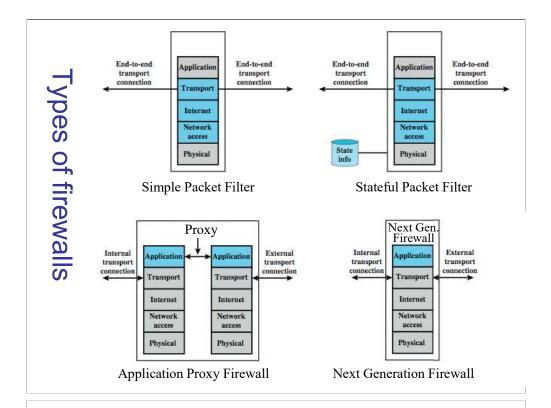


(Stateless) Packet Filter

- A packet filter is a network router that can accept/reject packets based on headers
- Packet filters examine each packet's headers and make decisions based on attributes such as:
 - Source or Destination IP Addresses
 - Source or Destination Port Numbers
 - Protocol (UDP, TCP or ICMP)
 - ICMP message type

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- And which interface the packet arrived on
- Unaware of session states at internal or external hosts
- High speed, but primitive filter



(Stateless) Packet Filters

- Widespread packet filter software (Linux):
 - iptables / netfilter
 - nft / nttables
- Examples (iptables)
- iptables -A FORWARD -s 131.234.142.33 -j ACCEPT
 - All packets from source IP Address 131.234.142.33 are accepted
- iptables -A FORWARD -p tcp -d 10.0.0.56 --dport 22 -j ACCEPT
 - All packets using transport protocol and destination address 10.0.0.56 and destination port 22 are accepted

Problems with Stateless Filtering

- Assume a typical "security policy":
 - Access from internal to external allowed
 - Access from external to internal prohibited
 - Example application: home network
- · Naive packet filter configuration:
 - outgoing packet → forward
 - incoming packet → reject

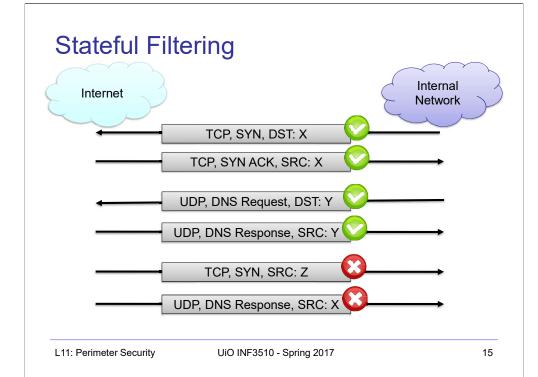
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Stateful Packet Filters

- · Stateful packet filters track current state of a connection
 - More 'intelligent' than simple packet filters.
- · Stateful packet filters keep track of sessions
 - Recognise if a particular packet is part of an established connection by 'remembering' recent traffic history.
 - Will add a temporary rule to allow the reply traffic back through the firewall.
 - When "session" is finished, the temporary rule is deleted.
- This makes the definition of filtering rules easier to accomplish and therefore potentially more secure.
- High speed, can use relatively advanced filter rules
- Requires memory
 - So can be subject to DOS (Denial of Service) attacks



Stateful Packet Filters

- Examples (iptables)
- iptables -A FORWARD -m state --state NEW -i eth0 -j ACCEPT
- Accept new connections (i.e. TCP SYN) from network interface eth0 ("from inside")
- iptables -A FORWARD -m state --state ESTABLISHED, RELATED -j ACCEPT
- Accept ALL packets which belong to an established TCP connection or are related to an existing UDP communication

(Stateful) Packet Filter: Evaluation

- Strengths:
 - Low overhead and high throughput
 - Supports almost any application
- Weaknesses:
 - Unable to interpret application layer data/commands
 - · may allow insecure operations to occur
 - Allows direct connection between hosts inside & outside firewall

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IPv4 Network Address Translation (NAT)

- NAT used to increase IPv4 address space
- Translates public IP addr. ↔ private IP addr. and ports
- Each local network can reuse private IP address ranges
 - Artificially increases the number of usable IP addresses
- · Possibilities:
 - Static mapping
 - permanent mapping of public to private address (no gain)
 - Dynamic mapping
 - · mapping of public to private address when needed
 - · unmapped when no longer needed
 - PAT (Port Address Translation)
 - multiple internal addresses mapped to same public address but with different port numbers

Personal Firewalls

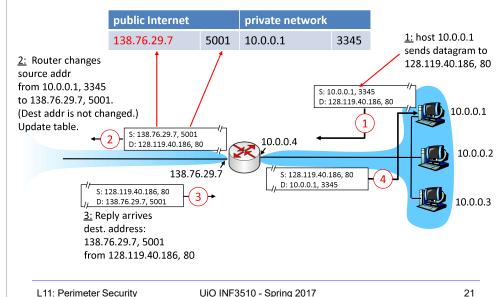
- A personal firewall is a program that is designed to protect the computer on which it is installed
- Personal firewalls are frequently used by home users to protect themselves from the Internet
- Nowadays for example included in Windows
- Advantage compared to network firewall: rules can take applications into account

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IPv4 Network Address Translation (NAT)



IPv4 NAT: + & -

- Advantages
 - Helps enforce control over outbound connections
 - Helps restrict incoming traffic
 - Helps conceal internal network configuration
 - Makes port scanning more difficult
- Can't be used with:
 - protocols that require a separate back-channel
 - protocols that encrypt TCP headers such as IPSec
 - embedded TCP address info
 - (Not recommended with) IPv6

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Next Generation Firewalls (NGFW)

- Inspects payload in end-to-end or proxy application connection
- Support specific application protocols
 - e.g. http, telnet, ftp, smtp etc.
 - each protocol supported by a specific proxy HW/SW module
- Can be configured to filter specific user applications
 - E.g. Facebook, Youtube, LinkedIn
 - Can filter detailed elements in each specific user application
- Can support TLS/SSL encrypted traffic inspection
- Can provide intrusion detection and intrusion prevention
- Very high processing load in firewall
 - High volume needs high performance hardware, or else will be slow



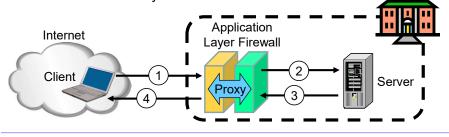




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Application Layer Proxy 1. External client sends a request to the

- 1. External client sends a request to the server, which is intercepted by the outwards-facing firewall proxy
- 2. Inwards-facing proxy sends request to server on behalf of client.
- 3. Server sends reply back to inwards-facing firewall proxy.
- 4. Outwards facing proxy sends reply to the client.
- Client and server both think they communicate directly with each other, not knowing that they actually talk with a proxy.
- The proxy can inspect the application data at any level of detail, and can even modify the data



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High performance NGFWs





High range model: *PA-7050* Up to 120 Gbps throughput

Prices starting from: US\$ 200,000





High range model: 61000 Security system

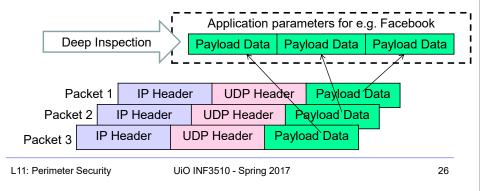
Up to 400 Gbps throughput

Prices starting from: US\$ 200,000

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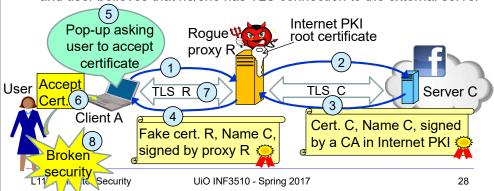
Inline Deep Packet Inspection

- Deep Packet Inspection looks at application content instead of individual or multiple packets.
- Deep inspection keeps track of application content across multiple packets.
- Potentially unlimited level of detail in traffic filtering



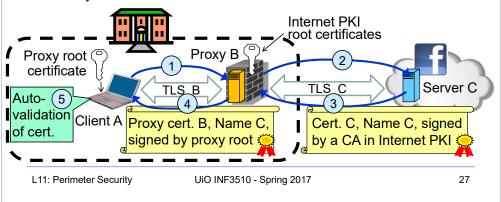
TLS inspection attack with rogue proxy server

- Depending on network, attackers may be able to install rogue proxy
- Rogue TLS inspect does not assume pre-installed proxy root certificate
- Proxy creates fake server certificate with the name of external server (e.g. facebook.com), that e.g. can be self-signed
- Fake server certificate is not validated, so browser asks user to accept it
- Fake certificate has (name = domain dame), so browser sets up TLS, and user believes that he/she has TLS connection to the external server



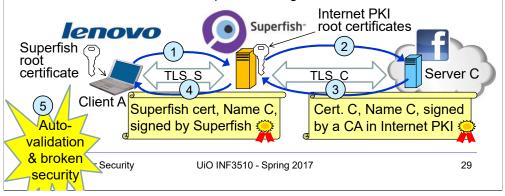
TLS/SSL encrypted traffic inspection in firewalls

- TLS designed for end-to-end encryption, normally impossible to inspect
- In order to inspect TLS, proxy must pretend to be external TLS server
- Proxy creates proxy server certificate with the name of external server (e.g. facebook.com), signed by local proxy root private key
- · Assumes that local proxy root certificate is installed on all local hosts
- The proxy server certificate is automatically validated by local client, so user may believe that he/she has TLS connection to the external server



Lenovo and the Superfish scam

- Superfish root certificate and diversion in shipped Lenovo models during 2014
- All https connections diverted to Superfish server to inject advertisements.
- Superfish created fake server certificates with names of web servers (e.g. facebook.com), signed by Superfish root private key.
- Fake server certificates were automatically validated, so users believed that he/she had secure end-to-end https connection to the web server.
- Scam discovered in 2015, Superfish cert. deleted and diversion removed.
- Embarrassment for Lenovo. Superfish changed name to JustVisual.



Application Proxy Firewalls + & -

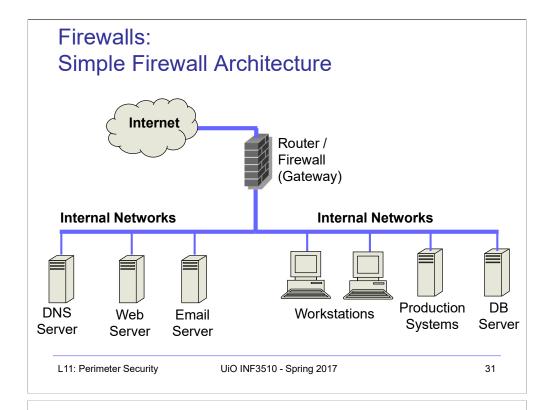
- · Strengths:
 - Easy logging and audit of all incoming traffic
 - Provides potential for best security through control of application layer data/commands
- Weaknesses:
 - May require some time for adapting to new applications
 - Much slower than packet filters
 - Much more expensive than packet filters

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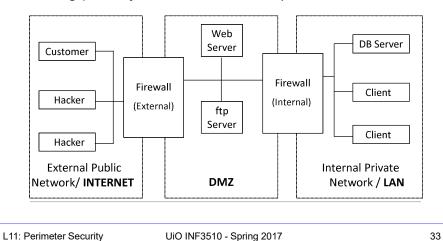
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Firewalls: **DMZ Firewall Architecture** Internet External Router / Firewall Production DB **DMZ (Demilitarized Zone)** Systems Server Workstations Internal Router / Firewall DNS Web Email Server Server Server **Internal Networks** L11: Perimeter Security UiO INF3510 - Spring 2017 32



DMZ Example

• DMZ = A part of your LAN with other restrictions, e.g. allowing publicly available services (web servers, mail etc.)



Intrusion Detection Systems

Intrusion Detection Systems:

- IDS are automated systems that detect suspicious activity
- IDS can be either host-based or network-based.
- A host based IDS is designed to detect intrusions only on the host it is installed on
 - monitor changes to host's OS files and traffic sent to the host
- Network based IDS (NIDS) detect intrusions on one or more network segments, to protect multiple hosts
 - monitor network/s looking for suspicious traffic
- What can be detected:
 - Attempted and successful misuse, both external and internal agents
 - Malware: Trojan programs, viruses and worms
 - DOS (Denial Of Service) attacks

Intrusion Detection and Prevention

Intrusion

 Actions aimed at compromising the security of a target network (confidentiality, integrity, availability of resources)

Intrusion detection

- The identification of possible intrusion through intrusion signatures and network activity analysis
- IDS: Intrusion Detection Systems

Intrusion prevention

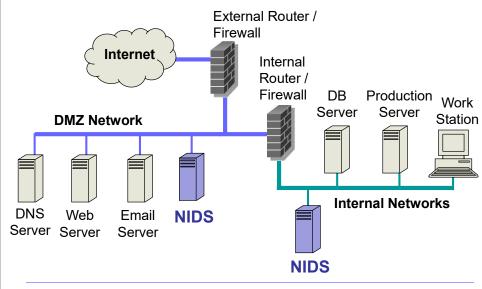
- The process of both detecting intrusion activities and managing automatic responsive actions throughout the network
- IPS: Intrusion Prevention Systems
- IDPS: Intrusion Detection and Prevention Systems

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Network IDS Deployment



Intrusion Detection Techniques

- Misuse detection
 - Use attack "signatures" (need a model of the attack)
 - · Sequences of system calls, patterns of network traffic, etc.
 - Must know in advance what attacker will do (how?)
 - Can only detect known attacks
 - Relatively few false positives
- Anomaly detection
 - Using a model of normal system behavior, try to detect deviations and abnormalities
 - E.g., raise an alarm when a statistically rare event(s) occurs
 - Can potentially detect unknown attacks
 - Many false positives

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Example: Vulnerability + Snort Rule

₩CVE-2017-0147 Detail

Current Description

The SMBv1 server in Microsoft Windows Vista SP2; Windows Server 2008 SP2 and R2 SP1; Windows 7 SP1; Windows 8.1; Windows Server 2012 Gold and R2; Windows RT 8.1; and Windows 10 Gold, 1511, and 1607; and Windows Server 2016 allows remote attackers to obtain sensitive information from process memory via a crafted packets, aka "Windows SMB Information Disclosure Vulnerability."

Source: MITRE

Description Last Modified: 03/16/2017

View Analysis Description

alert tcp \$HOME_NET 445 -> any any (msg:"OS-WINDOWS Microsoft Windows SMB possible leak of kernel heap memory"; flow:to_client,established; content:"Frag",fast_pattern; content:"Free"; content:"|FA FF FF|"; content:"|F8 FF FF|",within 3,distance 5; content:"|F8 FF FF|",within 3,distance 5; metadata:policy balanced-ips alert,policy security-ips drop,ruleset community; service:netbios-ssn; reference:cve,2017-0147; reference:url,technet.microsoft.com/en-us/security/bulletin/MS17-010; classtype:attempted-recon; sid:42339; rev:2;)

Popular NIDS



- Snort (popular open-source tool)
 - Large rule sets for known vulnerabilities, e.g.
 - 2009-03-31: A programming error in MySQL Server may allow a remote attacker to cause a Denial of Service (DoS) against a vulnerable machine.
 - 2009-03-27: Microsoft Windows GDI Buffer Overflow: A programming error in the Microsoft Windows kernel may allow a remote attacker to execute code with system level privileges. This may be exploited when specially crafted EMF files are viewed using Microsoft Internet Explorer.
- Bro (developed by Vern Paxson)

- Bro
- Separates data collection and security decisions
 - Event Engine distills the packet stream into high-level events describing what's happening on the network
 - Policy Script Interpeter uses a script defining the network's security policy to decide what to do in response

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Port Scanning

- Many vulnerabilities are OS-specific
 - Bugs in specific implementations, default configuration
- Port scan is often a prelude to an attack
 - Attacker tries many ports on many IP addresses
 - For example, looking for an old version of some daemon with an unpatched buffer overflow
 - If characteristic behavior detected, mount attack
 - "The Art of Intrusion": virtually every attack involves port scanning and password cracking

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Port Scanning

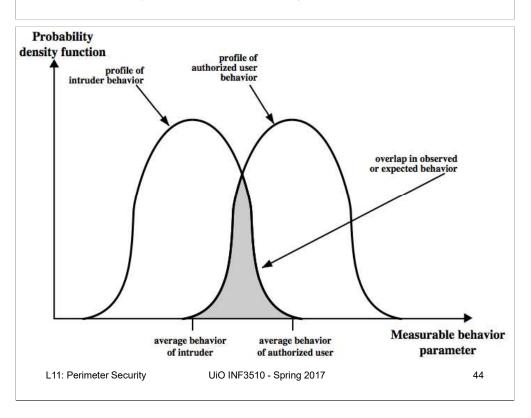
Example: network services on a Windows computer

Proto.	Local Address	Foreign Address	State
TCP	0.0.0.0:80	0.0.0.0:0	LISTEN
TCP	0.0.0.0:135	0.0.0.0:0	LISTEN
TCP	0.0.0.0:445	0.0.0.0:0	LISTEN
TCP	0.0.0.0:554	0.0.0.0:0	LISTEN
TCP	0.0.0.0:623	0.0.0.0:0	LISTEN
TCP	0.0.0.0:2869	0.0.0.0:0	LISTEN
TCP	0.0.0.0:5357	0.0.0.0:0	LISTEN
TCP	0.0.0.0:10243	0.0.0.0:0	LISTEN
TCP	0.0.0.0:16992	0.0.0.0:0	LISTEN
TCP	0.0.0.0:49152	0.0.0.0:0	LISTEN
TCP	0.0.0.0:49153	0.0.0.0:0	LISTEN
TCP	0.0.0.0:49154	0.0.0.0:0	LISTEN
TCP	0.0.0.0:49155	0.0.0.0:0	LISTEN
TCP	0.0.0.0:49157	0.0.0.0:0	LISTEN
TCP	0.0.0.0:56238	0.0.0.0:0	LISTEN

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Intrusion Detection Problems

- · Lack of training data with real attacks
 - But lots of "normal" network traffic, system call data
- Data drift
 - Statistical methods detect changes in behavior
 - Attacker can attack gradually and incrementally
- Discriminating characteristics hard to specify
 - Many attacks may be within bounds of "normal" range of activities
- False identifications are very costly
 - Sysadm will spend many hours examining evidence

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Intrusion Detection Errors

- · False negatives: attack is not detected
 - Big problem in signature-based misuse detection
- False positives: harmless behavior is classified as attack
 - Big problem in statistical anomaly detection
- · Both types of IDS suffer from both error types
- Both false positives and false negatives are problematic
 - Attacks are fairly rare events
 - IDS often suffer from "base-rate fallacy"

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Base Rate Fallacy

- Consider statements:
 - A: "attack occurs"
 - D: "detection occurs"
- We can measure/estimate:
 - P(D|A): probability of detection, given that attack occurs
 - P(D|!A): probability of detection, given that no attack occurs
 - *P*(*A*): probability of attack
- We want to know (false/true positives):
 - P(A|D): probability of attack, given that detection triggers
- · Bayes theorem:

-
$$P(A|D) = \frac{P(D|A) \cdot P(A)}{P(D)} = \frac{P(D|A) \cdot P(A)}{P(D|A) \cdot P(A) + P(D|A) \cdot P(A|A)}$$

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Remarks on Intrusion Detection

- Most alarms are false positives
 - Requires automated screening and filtering of alarms
- Most true positives are trivial incidents
 - can be ignored,
 - the attacks will never be able to penetrate any system
- Serious incidents need human attention
 - Can be dealt with locally
 - May require external expertise
- Potential for improvement through more intelligent IDS
 - Less false positives
 - Better detection of advanced attacks (APT)

Base Rate Fallacy

- Example:
 - Scanner is 99% correct: $P(D \mid A) = 0.99$, $P(D \mid A) = 0.01$
 - Attack probability: P(A) = 1 / 10000

•
$$P(A|D) = \frac{0.99 \cdot 0.0001}{0.99 \cdot 0.0001 + 0.01 \cdot 0.9999} = 0.010098$$

- Result:
 - 1% accuracy
 - 99 false positives per true positives

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Intrusion Prevention Systems

- Intrusion Prevention System (IPS) is a relatively new term that can mean different things
- Most commonly, an IPS is a combination of an IDS and a firewall
- A system that detects an attack and can stop it as well
- · Can be application specific
 - Deployed on a host to stop attacks on specific applications such as IIS
- Can be an extension of an NIDS
- False positives are problematic, because automated prevention measures can block services

Honeypots

- A honeypot:
 - is a computer configured to detect network attacks or malicious behavior,
 - appears to be part of a network, and seems to contain information or a resource of value to attackers.
- But honeypots are isolated, are never advertised and are continuously monitored
- All connections to honeypots are per definition malicious
- Can be used to extract attack signatures
- Honeynet is an international security club, see next slide

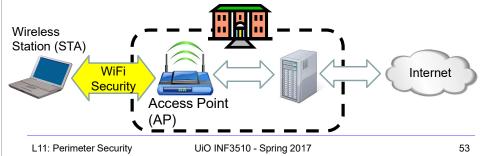
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IEEE 802.11 Standards for WLAN

- IEEE 802.11 formed in 1990's
 - charter to develop a protocol & transmission specifications for wireless LANs (WLANs)
- Since then the demand for WLANs, at different frequencies and data rates, has exploded
- New ever-expanding list of standards issued
 - from 10Mbps to 1Gbps transmission rate



WLAN Security



802.11 WiFi Security

- · Only authorized terminals (or users) may get access through Wireless LAN
- Should be impossible to set up rogue AP
- Interception of traffic by radios within range should be impossible

	WEP (1999) 801.11b	WPA (2003) 802.11i (subset)	WPA2 (2004) (aka. RSN) 802.11i (full set)
Auth. & key gen.	WEP	EAP	EAP
Encryption	RC4	RC4+TKIP	CCMP AES CTR (or TKIP)

- WEP: Wired Equivalent Privacy (broken)
- WPA: WiFi Protected Access
- EAP: Extensible Authentication Protocol
- RC4: Rivest Cipher 4 (a stream cipher)
- TKIP: Temporal-Key Integrity Protocol
- CCMP: Counter Mode with CBC Message Authentication Protocol
- RSN: Robust Security Network

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IEEE 802 Terminology

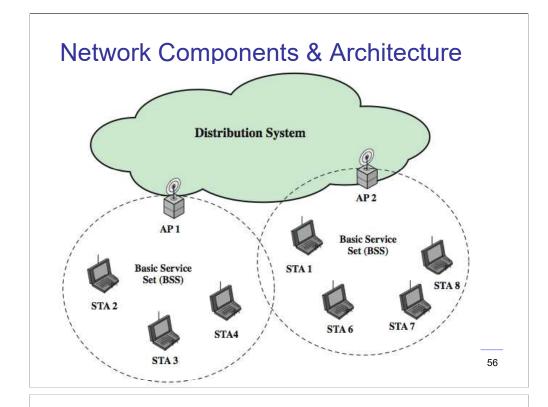
- Station (STA)
 - Wireless terminal that communicates with 802.11 functionality
- Access Point (AP)
 - Receives radio signals and controls access to network
- Basic Service Set (BSS)
 - Set of stations and one AP
- Extended Service Set (ESS)
 - Set of multiple BSSs
- Distribution System (DS)
 - Contains an Authentication Server (AS)
 - Integrates multiple BSSs into one ESS

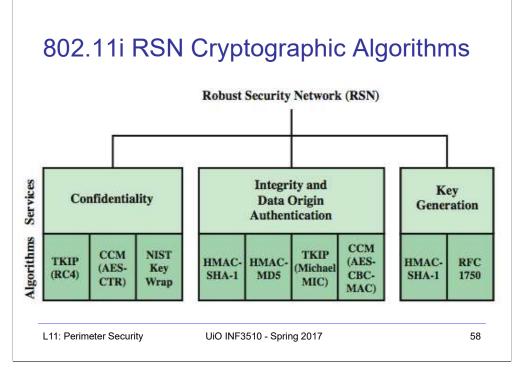
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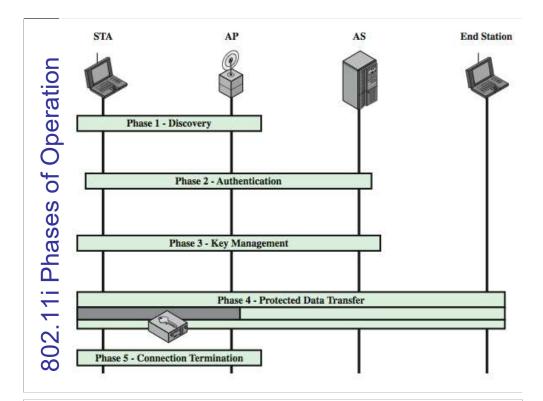
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802.11i RSN Services and Protocols Robust Security Network (RSN) Confidentialiy, Data Authentication Origin Authentication Access Control and Key and Integrity and Generation Replay Protection Protocols **IEEE 802.1** Extensible TKIP CCMP Port-based Authentication Access Control Protocol (EAP) UiO INF3510 - Spring 2017 57 L11: Perimeter Security





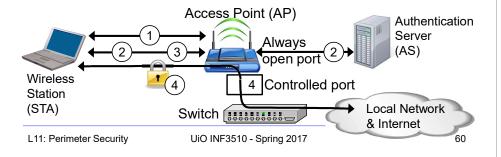


When you don't control the WLAN

- Often you want to connect to a wireless LAN over which you have no control, e.g. in café
- Options:
 - If you can, connect securely (WPA2, 802.11i, etc.)
 - · Beware of SSL-stripping
 - If unsecured, connect to online resources securely:
 - Use a VPN (Virtual Private Network)
 - IPSEC connection to home gateway
 - TLS/SSL connections to secure web server (with HSTS)
 - Be careful not to expose passwords
 - Watch for direct attacks on untrusted networks

802.11i WiFi Access Control

- 1. Mutual identity request between STA and AP
- 2. Mutual authentication between STA and AS.
- 3. Derive pairwise master key (PMK) between STA and AP.
- 4. Encrypt radio link and open port (connect) to network access
- Controlled port from AP to network
 - is closed (disconnected) before authentication
 - is open (connected) after successful authentication



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End of Lecture

This lecture presented:

- Firewall techniques
- Intrusion detection techniques
- WLAN Access