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## **INF 4140: Models of Concurrency**

Høst2015

## Series 6

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## **Topic: Program Analysis II**

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**Exercise 1 (While program)** Consider the following program S:

Prove that the following triple is a theorem in PL, i.e., that it is *derivable*:

$$\{b \ge 0\} S \{x = 2 * b\}.$$
(1)

You may use the following predicate I as loop invariant:

 $I: x \le y \land x = 2(y-b) . \tag{2}$ 

## **Exercise 2 (Factorial function)** Consider the following program S:

Prove the following triple using PL:

$$\{n \ge 0\} S \{x = n!\}$$
(3)

As a loop invariant I, you may use:

$$I: x = i! \land i \le n \tag{4}$$

You may assume the following when reasoning about the factorial function:

- 1) 0! = 1
- 2) (j+1)! = j! \* (j+1) for any integer  $j \ge 0$

**Exercise 3 (Monitor verification)** Consider the monitor for Shortest-Job-Next allocation in the book (section 5.2.3). Use Programming Logic, extended with rules for signal and wait (lecture slides, week 6), to prove that this monitor satisfies the second part of the SJN invariant:

$$\texttt{free} \Rightarrow (\#\texttt{turn} = 0) \tag{5}$$

(You may use the rule for wait(cv) to reason about wait(cv,rank)).

*Hint.* When arriving at an implication, it is enough to argue for the truth of it. However, we may use the following rules when reasoning about implications.

$$(A \land B) \Rightarrow C$$
 $(A \land B) \Rightarrow C$  $((A \land B) \Rightarrow C$  $(\neg A) \lor B$ false  $\Rightarrow A$  $A \Rightarrow (B \Rightarrow C)$  $((A \Rightarrow B) \land A) \Rightarrow C$  $A \Rightarrow B$ 

Exercise 4 Further exercises from the textbook:

2.22, 2.16, 2.24, 2.31, (2.28a, 2.29a)

Exercise 5 (For-loop) Design a rule for *for-loops*. ([?, Exercise 2.22])

**Exercise 6 (Verification of a parallel program ([?, Exercise 2.16]))** Consider the following parallel program.

```
int x := 0; { x = 0 }
1
2
   co
       Process_1: < await (x \neq 0) x := x - 2 >
3
   4
       Process<sub>2</sub>: < await (x \neq 0) x := x - 3>
\mathbf{5}
   6
       Process_3: < await (x=0) x := x + 5>
\overline{7}
8
   oc
```

Prove that the final value is 0.

**Exercise 7 (Interference freedom ([?, Exercise 2.24]))** Consider the following statement together with a pre-condition:

$$\{x \ge 4\} < x := x - 4 >$$
(6)

Then consider, whether the given statements *interfere* with it.

1)	$\{ x \ge 0 \}$	$\langle x := x + 5 \rangle$	$\{x \ge 5\}$
2)	$\{ x \ge 0 \}$	$\langle x := x + 5 \rangle$	$\{ x \ge 0 \}$
3)	$\{ x \ge 10 \}$	$\langle x := x + 5 \rangle$	$\{ x \ge 11 \}$
4)	$\{ x \ge 10 \}$	$\langle x := x + 5 \rangle$	$\{ x \ge 12 \}$
5)	$\{x \text{ is odd }\}$	$\langle x := x + 5 \rangle$	$\{x \text{ is even }\}$
6)	$\{x \text{ is odd }\}$	$\langle y := x + 1 \rangle$	$\{x \text{ is even }\}$
7)	$\{x \text{ is odd }\}$	$\langle y := y + 1 \rangle$	$\{x \text{ is even }\}$
8)	$\{x \text{ is a multiple of } 3\}$	$\langle y := x \rangle$	$\{ y \text{ is a multiple of } 3 \}$ .

Exercise 8 (Interference freedom (Exercise [?, 2.31])) Assume two triples

 $\{P_1\} S_1 \{Q_1\} \text{ and } \{P_2\} S_2 \{Q_2\}.$ 

Assume they are *interference* free (according to the definition). Assume that  $S_1$  contains an await-statement  $\langle \texttt{await}(B) T \rangle$ . Let then  $S'_1$  be the same as  $S_1$ , except that the awaitstatement is replaced by "corresponding" while-loop

- 1. Assume the triple {  $P_1$  }  $S_1$  {  $Q_1$  } holds. Then: Is {  $P_1$  }  $S'_1$  {  $Q_1$  } still true.
- 2. Are {  $P_1$  }  $S_1$  {  $Q_1$  } and {  $P_2$  }  $S_2$  {  $Q_2$  } still interference-free?

**Exercise 9 (Parallel boolean check (2.28a))** Check whether all elements in an array are set 0. See [?, Exercise 2.28(a), page 89].

Exercise 10 (Maximum (2.29a)) Determine the max from an integer array, searching for the even and odd numbers in parallel. See [?, Exercise 2.29, page 89].