

6.5 Rewrite the attribute grammar of Table 6.2 to compute a *postfix* string attribute instead of *val*, containing the postfix form for the simple integer expression. For example, the *postfix* attribute for **(34 - 3) * 42** is “34 3 – 42 + *.” You may assume a concatenation operator || and the existence of a **number.strval** attribute.

6.7 Consider the following grammar for simple Pascal-style declarations:

$$\begin{aligned} decl &\rightarrow var-list : type \\ var-list &\rightarrow var-list , id \mid id \\ type &\rightarrow \mathbf{integer} \mid \mathbf{real} \end{aligned}$$

Write an attribute grammar for the type of a variable.

6.13 Consider the following attribute grammar:

| Grammar Rule | Semantic Rules |
|---------------------------|---|
| $S \rightarrow A \ B \ C$ | $B.u = S.u$ $A.u = B.v + C.v$ $S.v = A.v$ |
| $A \rightarrow a$ | $A.v = 2 * A.u$ |
| $B \rightarrow b$ | $B.v = B.u$ |
| $C \rightarrow c$ | $C.v = 1$ |

- Draw the parse tree for the string abc (the only string in the language), and draw the dependency graph for the associated attributes. Describe a correct order for the evaluation of the attributes.
- Suppose that $S.u$ is assigned the value 3 before attribute evaluation begins. What is the value of $S.v$ when evaluation has finished?

- c. Suppose the attribute equations are modified as follows:

| Grammar Rule | Semantic Rules |
|---------------------------|---|
| $S \rightarrow A \ B \ C$ | $B.u = S.u$ $C.u = A.v$ $A.u = B.v + C.v$ |
| $A \rightarrow a$ | $S.v = A.v$ |
| $B \rightarrow b$ | $A.v = 2 * A.u$ |
| $C \rightarrow c$ | $B.v = B.u$ $C.v = C.u - 2$ |

What value does $S.v$ have after attribute evaluation, if $S.u = 3$ before evaluation begins?

```
class → class name superclass { decls }
decls → decls ; decl / decl
decl → variable-decl
decl → method-decl
method-decl → type name ( params ) body
type → int /bool / void
superclass → name
```

Ord i kursiv er non-terminaler, ord og tegn i fet skrift er terminal-symboler, mens **name** representerer et navn som scanneren leverer. Det kan antas at **name** har attributtet 'name'.

Metoder med samme navn som klassen er 'konstruktører', og det gjelder følgende regel: Konstruktører må være spesifisert med typen **void**.

Lag semantiske regler for denne regel i følgende fragment av en attributtgrammatikk.

| Grammar Rule | Semantic Rule |
|--|----------------------------|
| <i>class</i> → <i>class name { decls }</i> | |
| <i>decls</i> → <i>decls ; decl</i> | |
| <i>decls</i> → <i>decl</i> | |
| <i>decl</i> → <i>variable-decl</i> | <i>Skal ikke fylles ut</i> |
| <i>decl</i> → <i>method-decl</i> | |
| <i>method-decl</i> → <i>type name (params) body</i> | |
| <i>type</i> → int | |
| <i>type</i> → bool | |
| <i>type</i> → void | |