## UNIVERSITETET I OSLO Institutt for Informatikk



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## INF 5110: Compiler construction

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Topic: Chapter 4: grammars

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Exercise 1 (LL(1)) Check if the following grammar is LL(1)?

$$S \rightarrow (S)S \mid \epsilon$$

Exercise 2 (Ambiguity) Given the following grammar.

$$exp \rightarrow exp + exp \mid (exp) \mid if exp then exp else exp \mid var var \rightarrow \dots$$

- 1. Try to come up with an unambiguous grammar for the language of the given grammar, where
  - (a) addition is left-associative, and where
  - (b) if x then y else z + y is meant to mean if x then y else (z + y).
- 2. Why don't we have a dangling else problem here?

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Exercise 3 (Ambiguity) Given the following grammar.

$$exp \rightarrow exp \ op \ exp \mid (exp) \mid number$$
  
 $op \rightarrow + \mid - \mid * \mid / \mid \uparrow \mid < \mid =$ 

Do the following things.<sup>1</sup>

1. The grammar is pretty ambiguous. Make an unambiguous grammar capturing the same language, under the following side conditions

	precedence	assoc
$\uparrow$	highest (3)	right
*, /	level 2	left
+, -	level 1	left
<, =	0	non-associative

- 2. Give recursive-descent procedures for each non-terminal to check the grammar (using also loops, if advisable). Divide the terminals representing *op* in an appropriate manner
- 3. Based on the previous point: add tree-building code into the procedures in such a way that sequences of exponentiations \(\gamma\) are treated appropriately in the sense that the tree reflects the intended right-associativity.
- 4. Take the unambiguous grammar done in the first point, remove left-recursion, and do left-factorization (without destroying unambiguity).
- 5. Check whether the resulting grammar is LL(1).

<sup>&</sup>lt;sup>1</sup>There's a certain amount of repetition here, we won't go through everything during class-time, but a proposal for solution will be available.