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INF 5110: Compiler construction

Spring 2021 Series 7 12. 3. 2021

Topic: Run-time environments (Chapter 7)

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Exercise 1 (Run-time environment) Draw a possible organization for the runtime environment of the following C program, for the following two situations. See corresponding figures from the lecture as inspiration (for example, the slide entitled "Stack gcd", approximately at slide 8.18):

- 1. after entry into block A in function f.
- 2. after entry into block B in function g.

```
int a [10];
   char * s = "hello";
2
   int f(int i, int b[])
4
   {\bf int} j=i;
5
    A: \{ int i=j;
6
        char c = b[i];
7
8
        // . . . ;
9
      return 0;
10
11
12
   void g(char * s)
13
   \{ char c = s[0]; 
14
    B:{ int a[5];
15
16
17
18
19
   main ()
20
   \{ int x=1; 
21
      x = f(x,a);
22
23
      g(s);
      return 0;
24
25
```

Exercise 2 (Activation records (Pascal)) Draw the stack of activation reords for the following *Pascal* program, showing the *control* and *access* links, after the second call to procedure c. Describe how the variable x is accessed from within c.

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```
program env;
2
   procedure a;
3
   var x: integer;
4
5
       procedure b;
6
           procedure c;
           begin
8
              x := 2;
9
              b;
10
           end;
11
       begin (* b *)
12
           c:
13
       end;
14
15
   begin (* a *)
16
       b;
17
   end;
18
19
   begin (* main *)
20
21
       a;
   end.
```

Exercise 3 (Access chaining vs. display) An alternative to access chaining in a language with local procedures is to keep the access links in an array *outside* the stack, *indexed* by the *nesting level*. This array is called the *display*. For example, the run-time stacks of the program chain and the corresponding stack picture on the slide entitled "access chaining" at approx. 8-36 from the lecture would now look as Figure 1 resp. Figure 2.

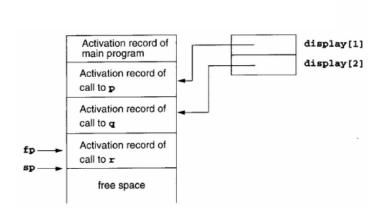


Figure 1: RTE with display (1)

- 1. Describe how a display can improve *efficiency* of nonlocal references from deeply nested procedures.
- 2. Redo Exercise 2 from this sheet, using a display.

Exercise 4 (Virtual function tables and memory layout for classes) Draw the memory layout of objects of the following C^{++} classes, together with the *virtual function tables*.

```
class A { public:
```

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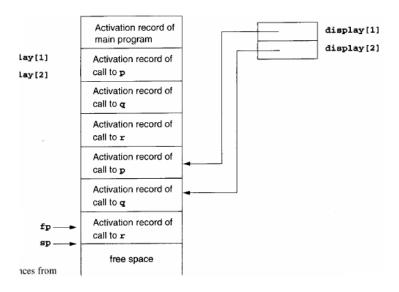


Figure 2: RTE with display (2)

```
int a;
3
     virtual void f();
4
     virtual void g();
5
   };
6
   class B: public A
8
   { public:
9
     int b;
10
     virtual void f();
11
     void h();
12
   };
13
14
15
   class C: public B
16
   \{ public:
17
     int c;
     virtual void g();
19
20
```

Exercise 5 (Parameter passing) Give the output of the following program (written in C syntax) using the 4 parameter passing methods discussed in in the lecture.

```
#include <stdio.h>
  int i = 0;
2
3
  void p(int x, int y)
4
  \{ x += 1;
5
    i += 1;
    y += 1;
  }
8
9
  main ()
10
  \{ int a[2] = \{1,1\};
11
    p(a[i],a[i]);
12
    13
    return 0;
14
15
```

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Exercise 6 (Parameter passing) Give the output of the following program (written in C syntax) using the 4 parameter passing methods discussed in the lecture.

```
#include <stdio.h>
   int i = 0;
2
   void swap (int x, int y)
5
      x = x + y;
6
      y = x - y;
7
      x = x - y;
8
9
10
   main ()
11
   \{ \text{ int } a[3] = \{1,2,0\}; 
12
      swap\left(\,i\,\,,a\left[\,i\,\,\right]\,\right);
13
      printf("%d\_%d\_%d\_%d\n", i, a[0], a[1], a[2]);
14
15
16
```