
Least Squares Problems

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Overdetermined Equations

- Given $A^{m,n}$ and $b \in \mathbb{R}^m$.
- The system $Ax = b$ is **over-determined** if $m > n$.
- This system has a solution if $b \in \text{Span}(A)$, the column space of A , but normally this is not the case and we can only find an approximate solution.
- A general approach is to choose a vector norm $\|\cdot\|$ and find x which minimizes $\|Ax - b\|$.
- We will only consider the Euclidian norm here.

The Least Squares Problem

- Given $A^{m,n}$ and $\mathbf{b} \in \mathbb{R}^m$ with $m \geq n \geq 1$. The problem to find $\mathbf{x} \in \mathbb{R}^n$ that minimizes $\|A\mathbf{x} - \mathbf{b}\|_2$ is called the **least squares problem**.
- A minimizing vector \mathbf{x} is called a **least squares solution** of $A\mathbf{x} = \mathbf{b}$.

Example 1



$$\begin{array}{l} x_1 = 1 \\ x_1 = 1, \\ x_1 = 2 \end{array} \quad A = \begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix}, \quad \mathbf{x} = [x_1], \quad \mathbf{b} = \begin{bmatrix} 1 \\ 1 \\ 2 \end{bmatrix},$$

- Our approach is to minimize the least squares sum

$$\|A\mathbf{x} - \mathbf{b}\|_2^2 = (x_1 - 1)^2 + (x_1 - 1)^2 + (x_1 - 2)^2.$$

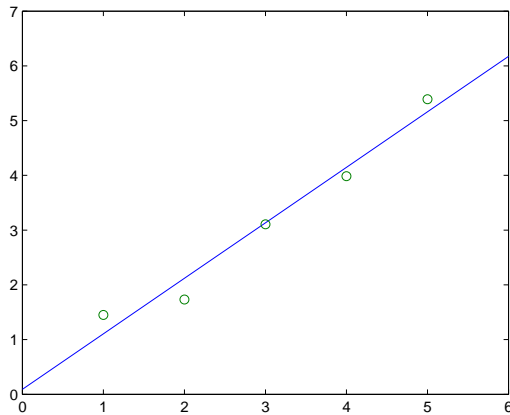
- Setting the first derivative with respect to x_1 equal to zero we obtain $2(x_1 - 1) + 2(x_1 - 1) + 2(x_1 - 2) = 0$ or $6x_1 - 8 = 0$ or $x_1 = 4/3$
- The second derivative is positive (it is equal to 6) and $x = 4/3$ is a global minimum.

Linear regression

• Given m points $(t_i, y_i)_{i=1}^m$ in the (t, y) plane.

• Example:

$$(t_i, y_i)_{i=1}^5 = [(1, 1.4501), (2, 1.7311), (3, 3.1068), (4, 3.9860), (5, 5.3913)]$$



• Find a straight line $y(t) = x_1 + tx_2$ such that $y_i = x_1 + t_i x_2$ all i .

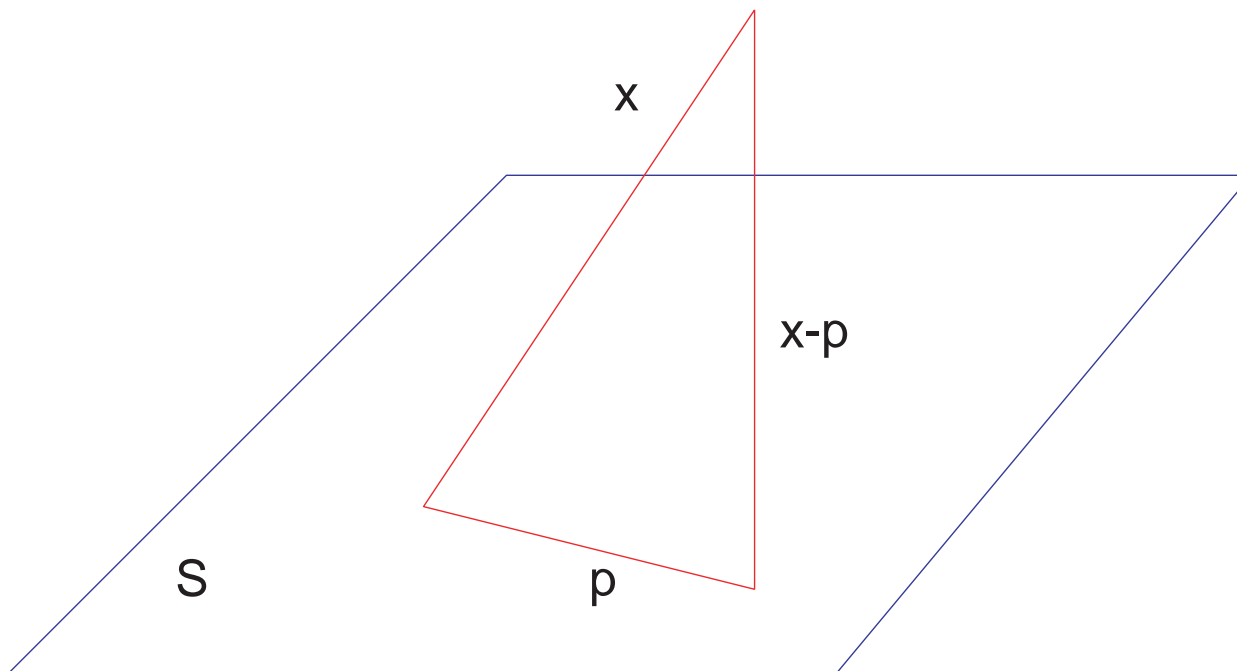
• We obtain a least squares problem with

$$A = \begin{bmatrix} 1 & t_1 \\ 1 & t_2 \\ \vdots & \vdots \\ 1 & t_m \end{bmatrix} \quad \mathbf{b} = \begin{bmatrix} y_1 \\ y_2 \\ \vdots \\ y_m \end{bmatrix}, \quad \|A\mathbf{x} - \mathbf{b}\|_2^2 = \sum_{i=1}^m (x_1 + t_i x_2 - y_i)^2.$$

Recall result on orthogonal projection

Theorem(Best Approximation) Let \mathcal{S} be a subspace in a real or complex inner product space $(\mathcal{V}, \mathbb{F}, \langle \cdot, \cdot \rangle)$. Let $x \in \mathcal{V}$, and $p \in \mathcal{S}$. The following statements are equivalent

1. $\langle x - p, s \rangle = 0$, for all $s \in \mathcal{S}$.
2. $\|x - s\| > \|x - p\|$ for all $s \in \mathcal{S}$ with $s \neq p$.



Existence and Uniqueness

Theorem 1. *The least squares problem always has a solution. The solution is unique if and only if A has linearly independent columns.*

Proof. ● We apply the inner product setup with $\mathcal{V} = \mathbb{R}^n$, the usual inner product in \mathbb{R}^n , \mathcal{S} equals $\text{Span}(A) := \{Ax : x \in \mathbb{R}^n\}$, the column space of A , and $x = b$.

● The inner product norm is the Euclidian norm $\|\cdot\|_2$.

● Let p be the orthogonal projection of b into $C(A)$. Since $p \in \text{Span}(A)$ there is an $x^* \in \mathbb{R}^n$ such that $Ax^* = p$.

● If $x \in \mathbb{R}^n$ and $x \neq x^*$ then $s := Ax \in \text{Span}(A)$ and by the best approximation theorem $\|b - Ax^*\|_2 = \|b - p\|_2 < \|b - s\|_2$.

Since p is unique and $Ax^* = p$ the least squares problem has a unique solution if and only if A has linearly independent columns. □

Characterization of any least squares solution

By the best approximation theorem $\langle \mathbf{b} - \mathbf{p}, \mathbf{s} \rangle = 0$ for all $\mathbf{s} \in \text{Span}(A)$ and we see that $\mathbf{b} - \mathbf{p}$ belongs to the orthogonal complement $\text{Span}(A)^\perp$ of A . We now recall:

Theorem 2. *The orthogonal complement of the column space of a matrix $A \in \mathbb{R}^{m,n}$ is the null space of A^T . In symbols*

$$\text{Span}(A)^\perp = \text{Ker}(A^T) := \{\mathbf{y} \in \mathbb{R}^m : A^T \mathbf{y} = \mathbf{0}\}. \quad (1)$$

It follows that \mathbf{x}^* minimizes $\|A\mathbf{x} - \mathbf{b}\|_2$ if and only if $\mathbf{b} - A\mathbf{x}^*$ belongs to $\text{Span}(A)^\perp$ or $A^T(\mathbf{b} - A\mathbf{x}^*) = \mathbf{0}$.

We obtain the linear system $A^T A\mathbf{x}^* = A^T \mathbf{b}$

The Normal Equations

The least squares solution can be found by solving a linear system.

Theorem 3. Suppose $A \in \mathbb{R}^{m,n}$ with $m > n$ and $\mathbf{b} \in \mathbb{R}^m$. The following is equivalent

1. \mathbf{x}^* minimizes $\|A\mathbf{x} - \mathbf{b}\|_2$.
2. $A^T A\mathbf{x}^* = A^T \mathbf{b}$. *(normal equations)*

The matrix $A^T A$ is nonsingular if and only if A has linearly independent columns.

Proof. The previous discussion shows that $1 \Leftrightarrow 2$. From an example in Chapter 8 we know that $A^T A$ is positive semidefinite and positive definite (and hence nonsingular) if and only if A has linearly independent columns. □

The linear system $A^T A\mathbf{x} = A^T \mathbf{b}$ is called the **normal equations**.

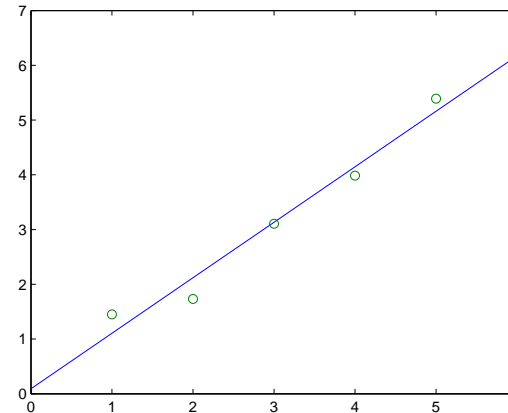
An observation

- We show that minimizing $\|Ax - \mathbf{b}\|_2$ and $\|Ax - \mathbf{b}\|_2^2$ are equivalent. Define functions $f : \mathbb{R}^n \rightarrow [0, \infty)$, $\phi : [0, \infty) \rightarrow [0, \infty)$, and $g : \mathbb{R}^n \rightarrow [0, \infty)$ by
- $f(\mathbf{x}) := \|Ax - \mathbf{b}\|_2$
- $\phi(t) = t^2$
- $g(\mathbf{x}) := \phi(f(\mathbf{x})) = \|Ax - \mathbf{b}\|_2^2$.
- Then f has a global minimum at \mathbf{x}^* if and only if g has a global minimum at \mathbf{x}^* .
- This follows since ϕ is strictly increasing.
- For if $h \neq 0$ then $g(\mathbf{x}^* + h) - g(\mathbf{x}^*) = \phi(f(\mathbf{x}^* + h)) - \phi(f(\mathbf{x}^*))$ so one difference is positive if and only the other is positive.

Linear Regression

$$A = \begin{bmatrix} 1 & t_1 \\ 1 & t_2 \\ \vdots & \vdots \\ 1 & t_m \end{bmatrix} \quad \mathbf{b} = \begin{bmatrix} y_1 \\ y_2 \\ \vdots \\ y_m \end{bmatrix}, \quad \min_{x_1 x_2} \sum_{i=1}^m (x_1 + t_i x_2 - y_i)^2.$$

$$A = \begin{bmatrix} 1 & 1 \\ 1 & 2 \\ 1 & 3 \\ 1 & 4 \\ 1 & 5 \end{bmatrix} \quad \mathbf{y} = \begin{bmatrix} 1.4501 \\ 1.7311 \\ 3.1068 \\ 3.9860 \\ 5.3913 \end{bmatrix}$$



$$A^T A = \begin{bmatrix} 5 & 15 \\ 15 & 55 \end{bmatrix} \quad \mathbf{c} = A^T \mathbf{y} = \begin{bmatrix} 16.0620 \\ 58.6367 \end{bmatrix} \quad \begin{array}{l} 5x_1 + 15x_2 = 16.0620 \\ 15x_1 + 55x_2 = 58.6367 \end{array}$$

$$\begin{array}{l} x_1 = 0.0772 \\ x_2 = 1.0451 \end{array}, \quad \mathbf{x} = A \setminus \mathbf{y}$$

Cholesky Factorization

1. Form the normal equations $B = A^T A$ and $c = A^T b$.
2. Find the Cholesky factorization $B = LL^T$ and find x .

1. for i=1:n
 for j=i:n
 B(i,j)=A(:,i)' \ast A(:,j);
 end
 c(i)=A(:,i)' \ast b;
end
 % Compute only one half of B
 % mn(n+1) flops
 % mn flops
2. $n^3/3$ flops

Forming the normal equations requires $O(mn^2)$ flops and this represents more work than solving them.

Why not always use the normal equations?

- Forming $A^T A$ squares the condition number of A .
- If A is ill conditioned then $A^T A$ will be severely ill conditioned
- Sometimes in applications A does not have full rank.
- Consider two other methods
- QR factorization
- Singular value decomposition

Squaring the Condition Number

The difference between the computed solution y of $Bx = c$ and the exact solution x satisfies

$$\frac{1}{\text{cond}_2(B)} \frac{\|r\|_2}{\|c\|_2} \leq \frac{\|x - y\|_2}{\|x\|_2} \leq \text{cond}_2(B) \frac{\|r\|_2}{\|c\|_2},$$

where $r = c - By$. If $\sigma_1 \geq \sigma_2 \geq \dots \geq \sigma_n > 0$ are the singular values of A then $\sigma_1^2, \sigma_2^2, \dots, \sigma_n^2$ are the singular values of $B = A^T A$ and

$$\text{cond}_2(B) = \frac{\sigma_1^2}{\sigma_n^2} = (\text{cond}_2(A))^2$$

For this reason one should not use Cholesky factorization to solve the least squares problem when A has a large condition number.

QR -factorization

Suppose A has full rank and that $A = Q_1 R_1$ is a reduced QR -factorization of A , that is $Q_1 \in \mathbb{R}^{m,n}$ has orthonormal columns and $R_1 \in \mathbb{R}^{n,n}$ is nonsingular and upper triangular. Then

$$A^T A x = A^T b \Rightarrow R_1^T Q_1^T Q_1 R_1 x = R_1^T Q_1^T b \Rightarrow R_1^T R_1 x = R_1^T Q_1^T b.$$

Since R_1^T is nonsingular we find the least squares solution by solving the triangular system

$$R_1 x = Q_1^T b.$$

Example

$$A = \begin{bmatrix} 1 & 1 \\ 1 & 2 \\ 1 & 3 \\ 1 & 4 \\ 1 & 5 \end{bmatrix} = \begin{bmatrix} -0.4472 & -0.6325 \\ -0.4472 & -0.3162 \\ -0.4472 & 0.0000 \\ -0.4472 & 0.3162 \\ -0.4472 & 0.6325 \end{bmatrix} * \begin{bmatrix} -2.2361 & -6.7082 \\ 0 & 3.1623 \end{bmatrix} = Q_1 R_1,$$

$$b = \begin{bmatrix} 1.4501 \\ 1.7311 \\ 3.1068 \\ 3.9860 \\ 5.3913 \end{bmatrix}, \quad c := Q_1^T b = \begin{bmatrix} -7.1831 \\ 3.3048 \end{bmatrix}$$

$$R_1 x = c \Leftrightarrow \begin{array}{l} -2.2361x_1 - 6.7082x_2 = -7.1831 \\ 3.1623x_2 = 3.3048 \end{array} \Leftrightarrow \begin{array}{l} x_1 = 0.0772 \\ x_2 = 1.0451 \end{array}.$$

Cholesky and Householder

Cholesky

- $O(mn^2)$ flops
- works only for full rank problems
- squares the condition number of A

Householder

- $O(2mn^2) - n^3/6$ flops.
- Can be used for rank deficient problems

Discussion:

- Cholesky more economical
- QR factorization should be used for problems with large condition numbers
- Cholesky cannot be used for rank deficient problems

The Singular Value Decomposition (SVD)

Suppose $A = U\Sigma V^T = U_1\Sigma_1V_1^T$ is the SVD of $A \in \mathbb{R}^{m,n}$. Recall that if $\text{rank}(A) = r$ then

- $\Sigma = \begin{bmatrix} \Sigma_1 & 0 \\ 0 & 0 \end{bmatrix}$

- $\Sigma_1 = \text{diag}(\sigma_1, \dots, \sigma_r), \sigma_1 \geq \dots \geq \sigma_r > 0,$



$$U \in \mathbb{R}^{m,m}, U^T U = I, U = (U_1, U_2), U_1 \in \mathbb{R}^{m,r}, U_2 \in \mathbb{R}^{m,m-r}$$

- $V \in \mathbb{R}^{n,n}, V^T V = I, V = (V_1, V_2), V_1 \in \mathbb{R}^{n,r}, V_2 \in \mathbb{R}^{n,n-r}.$

- $AV_2 = 0$ and V_2 is an orthonormal basis for $N(A)$

The Pseudo-inverse of $A = U\Sigma V^T$

The **pseudo-inverse** of $A \in \mathbb{R}^{m,n}$ is a matrix $A^\dagger \in \mathbb{R}^{n,m}$ given by

$$A^\dagger := V\Sigma^\dagger U^T = V_1\Sigma_1^{-1}U_1^T, \text{ where } \Sigma^\dagger := \begin{bmatrix} \Sigma_1^{-1} & 0 \\ 0 & 0 \end{bmatrix}, \quad (2)$$

Example:

$$A = \begin{bmatrix} 1 & 1 \\ 1 & 1 \\ 0 & 0 \end{bmatrix} = U\Sigma V^T = \frac{1}{\sqrt{2}} \begin{bmatrix} -1 & -1 & 0 \\ -1 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 2 & 0 \\ 0 & 0 \\ 0 & 0 \end{bmatrix} \frac{1}{\sqrt{2}} \begin{bmatrix} -1 & 1 \\ -1 & -1 \end{bmatrix}$$

$$\Sigma^\dagger = \begin{bmatrix} \frac{1}{2} & 0 & 0 \\ 0 & 0 & 0 \end{bmatrix}, \quad A^\dagger = V\Sigma^\dagger U^T = \frac{1}{4} \begin{bmatrix} 1 & 1 & 0 \\ 1 & 1 & 0 \end{bmatrix}$$

$$A^\dagger := V\Sigma^\dagger U^T$$

- Since Σ_1 is nonsingular A^\dagger is always defined.
- generalization of usual inverse. To see this note that
- $AA^\dagger = U_1\Sigma_1V_1^T V_1\Sigma_1^{-1}U_1^T = U_1U_1^T$.
- $A^\dagger A = V_1\Sigma_1^{-1}U_1^T U_1\Sigma_1V_1^T = V_1V_1^T$.
- If $A \in \mathbb{R}^{n,n}$ is square and nonsingular then $U_1 = U$ and $V_1 = V$ and $AA^\dagger = A^\dagger A = I$. Thus A^\dagger is the usual inverse.

Analysis of LSQ using SVD

Define

$$\mathbf{c} := U^T \mathbf{b} = \begin{bmatrix} U_1^T \mathbf{b} \\ U_2^T \mathbf{b} \end{bmatrix} = \begin{bmatrix} \mathbf{c}_1 \\ \mathbf{c}_2 \end{bmatrix}, \quad \mathbf{y} := V^T \mathbf{x} = \begin{bmatrix} V_1^T \mathbf{x} \\ V_2^T \mathbf{x} \end{bmatrix} = \begin{bmatrix} \mathbf{y}_1 \\ \mathbf{y}_2 \end{bmatrix}$$

$$\begin{aligned} \|\mathbf{b} - A\mathbf{x}\|_2^2 &= \|UU^T \mathbf{b} - U\Sigma V^T \mathbf{x}\|_2^2 = \|\mathbf{c} - \Sigma \mathbf{y}\|_2^2 = \left\| \begin{bmatrix} \mathbf{c}_1 \\ \mathbf{c}_2 \end{bmatrix} - \begin{bmatrix} \Sigma_1 & 0 \\ 0 & 0 \end{bmatrix} \begin{bmatrix} \mathbf{y}_1 \\ \mathbf{y}_2 \end{bmatrix} \right\|_2^2 \\ &= \left\| \begin{bmatrix} \mathbf{c}_1 - \Sigma_1 \mathbf{y}_1 \\ \mathbf{c}_2 \end{bmatrix} \right\|_2^2 = \|\mathbf{c}_1 - \Sigma_1 \mathbf{y}_1\|_2^2 + \|\mathbf{c}_2\|_2^2. \end{aligned}$$

We have $\|\mathbf{b} - A\mathbf{x}\|_2 \geq \|\mathbf{c}_2\|_2$ for all $\mathbf{x} \in \mathbb{R}^n$ with equality iff

$$\mathbf{x} = V\mathbf{y} = \begin{bmatrix} V_1 & V_2 \end{bmatrix} \begin{bmatrix} \Sigma_1^{-1} \mathbf{c}_1 \\ \mathbf{y}_2 \end{bmatrix} = A^\dagger \mathbf{b} + V_2 \mathbf{y}_2, \text{ for all } \mathbf{y}_2 \in \mathbb{R}^{n-r}. \quad (3)$$

Projections

The general solution of $\min \|A\mathbf{x} - \mathbf{b}\|_2$ is given by

$$\mathbf{x} = V\mathbf{y} = [V_1 \ V_2] \begin{bmatrix} \Sigma_1^{-1} U_1^T \mathbf{b} \\ \mathbf{y}_2 \end{bmatrix} = A^\dagger \mathbf{b} + V_2 \mathbf{y}_2, \text{ for all } \mathbf{y}_2 \in \mathbb{R}^{n-r}. \quad (4)$$

- The solution is unique if and only if $V_1 = V$ that is if and only if $r = \text{rank}(A) = n$.
- Since $AV_2 = 0$ we have $A\mathbf{x} = AA^\dagger \mathbf{b} + AV_2 \mathbf{y}_2 = AA^\dagger \mathbf{b}$
- $\mathbf{b}_1 := AA^\dagger \mathbf{b} = U_1 U_1^T \mathbf{b}$ is the vector projection of \mathbf{b} onto the column space $\text{Span}(A)$ of A
- $\mathbf{b}_2 := \mathbf{b} - \mathbf{b}_1 = (I - U_1 U_1^T) \mathbf{b} = U_2 U_2^T \mathbf{b}$ is the vector projection of \mathbf{b} onto $N(A^T)$

The Minimal Norm Solution

Suppose A is rank deficient ($r < n$) and let

$$\mathbf{z} := V\mathbf{y} = [V_1 \ V_2] \begin{bmatrix} \Sigma_1^{-1}U_1^T\mathbf{b} \\ \mathbf{y}_2 \end{bmatrix} = A^\dagger\mathbf{b} + V_2\mathbf{y}_2 \text{ with } \mathbf{y}_2 \in \mathbb{R}^{n-r} \text{ and}$$

nonzero. Let $\mathbf{x} = A^\dagger\mathbf{b}$ be the solution corresponding to

$\mathbf{y}_2 = 0$. Then $\|\mathbf{z}\|_2^2 = \|A^\dagger\mathbf{b}\|_2^2 + \|\mathbf{y}_2\|_2^2 > \|\mathbf{x}\|_2^2$ The vector

$\mathbf{x} = A^\dagger\mathbf{b}$ is the **minimal norm solution** to the LSQ problem.

Three Algorithms for solving LSQ

- Solving the normal equations by Cholesky factorization
- Using the QR -factorization of A
- Using the SVD of A

QR with column pivoting or SVD is used for rank deficient problems. (Non unique solutions)

Minimal norm solution by SVD

The minimal norm solution is $\mathbf{x} = A^\dagger \mathbf{b} = V_1 \Sigma_1^{-1} U_1^T \mathbf{b}$.

1. Compute $A = U \Sigma V^T = U_1 \Sigma_1 V_1^T$, the SVD of A .

2. $\mathbf{c} = U_1^T \mathbf{b}$

rm flops

3. $\mathbf{d} = \Sigma_1^{-1} \mathbf{c}$

r flops

4. $\mathbf{x} = V_1 \mathbf{d}$

rn flops

This is a great method if the SVD of A is known.

Computing the SVD of $A \in \mathbb{R}^{m,n}$

1. Transform A to **bidiagonal** form using Householder reflections
 2. Use the QR -method to find the singular values.
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