

# Multimedia Network Processor Examples

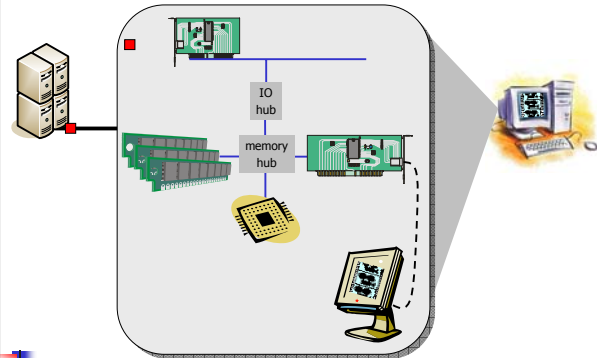
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## Overview

- Video Client Operations
- Multicast Video-Quality Adjustment
- Booster Boxes

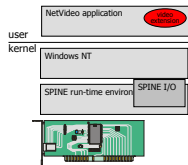
## Example: Video Client Operations

## Video Client Operations



## SPINE: Video Client Operations

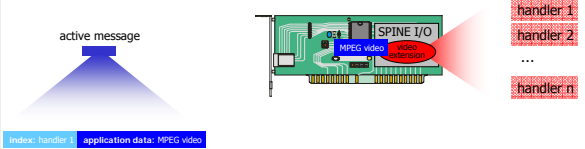
- Fiuczynski et. al. 1998:**
  - use an extensible execution environment to enable applications to compute on the network card
  - SPINE extends SPIN (an extensible operating system) to the network interface
    - define I/O extensions in Modula-3 (type-safe, Pascal-like)
    - these I/O modules may be dynamically loaded onto the NIC, or into the kernel (as in SPIN)
- perform video client operations on-board a *Myrinet* network card (33 Mhz LANai CPU, 256 KB SRAM)



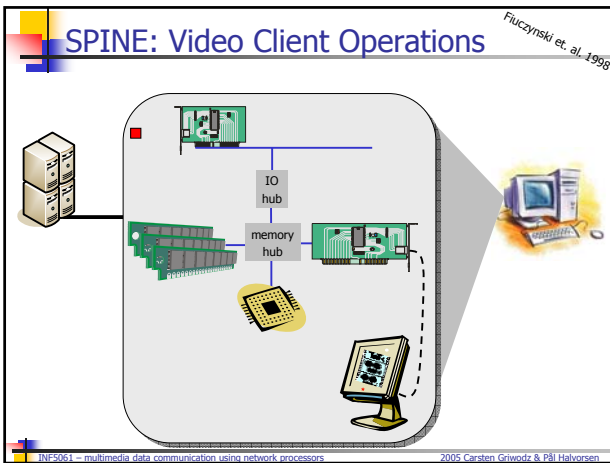
## SPINE: Video Client Operations

Fiuczynski et. al. 1998

- A message-driven architecture



- Application creates the framing window and informs the SPINE extension about the coordinates
- SPINE puts video data in corresponding frame buffer memory according to window placement on screen

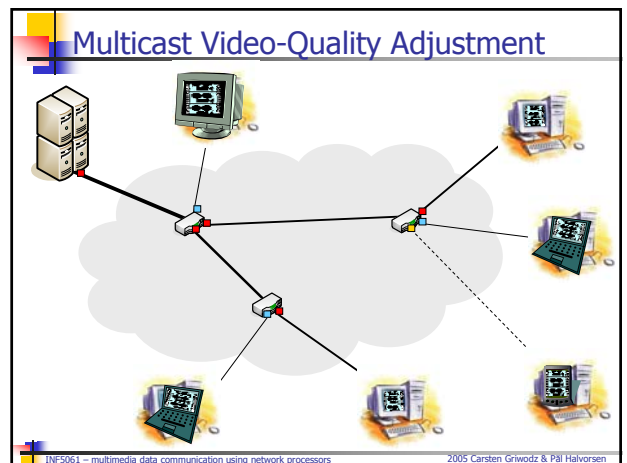
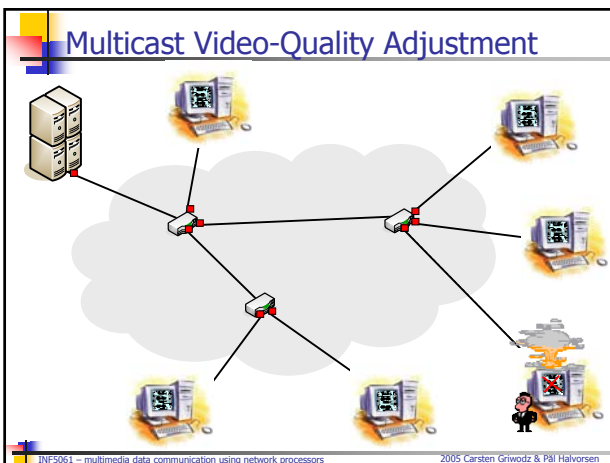


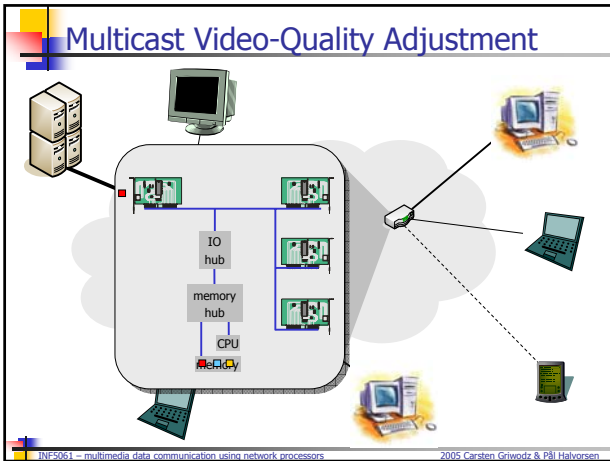
- ### SPINE: Video Client Operations
- Fiuczynski et. al. 1998
- Evaluation
    - managed to support several clients in different windows
    - data DMA'ed to frame buffer
      - zero host CPU requirement for video client(s)
    - a 33 Mhz LANai CPU too slow to do large video decoding operations
      - server converted MPEG to raw bitmap before sending
      - only I/O processing and data movement offloading
    - frequent synchronization between host and device-based component is expensive
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- ### SPINE: Internet Protocol Routing
- A SPINE router extension on the network processor
  - Able to fully offload host CPU
  - Forwarding latency 6% slower compared to host (but 33MHz embedded CPU vs. 200MHz Pentium Pro)
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### Example:

### Multicast Video-Quality Adjustment





- ### Multicast Video-Quality Adjustment
- Several ways to do video-quality adjustments
    - frame dropping
    - re-quantization
    - scalable video codec
    - ...
  - Yamada et. al. 2002:
    - use **low-pass filter** to eliminate high-frequency components of the MPEG-2 video signal and thus reduce data rate
      - determine a low-pass parameter for each GOP
      - use low-pass parameter to calculate how many DCT coefficients to remove from each macro block in a picture
      - by eliminating the specified number of DCT coefficients the video data rate is reduced
    - implemented the low-pass filter on an IXP1200
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### Multicast Video-Quality Adjustment

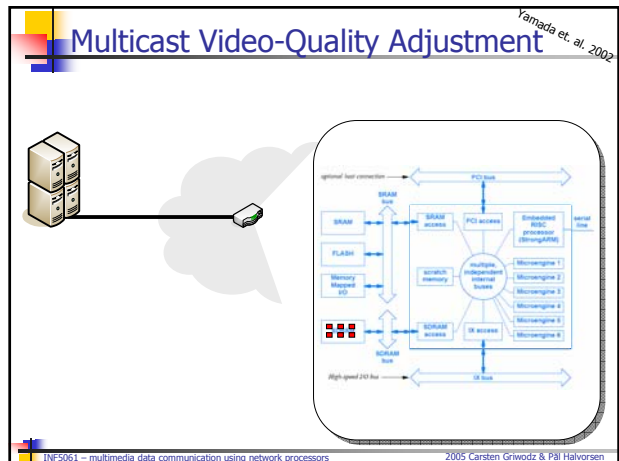
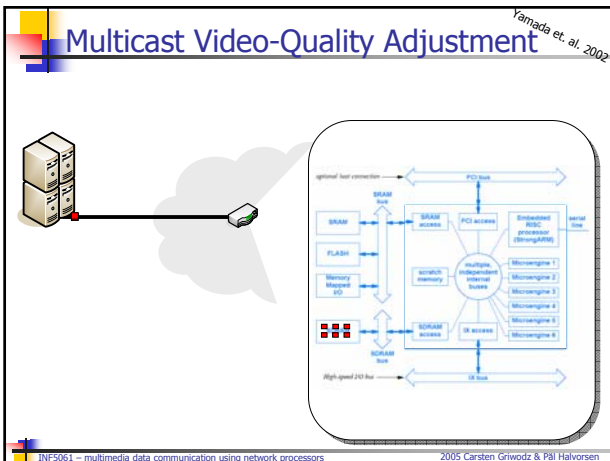
Yamada et. al. 2002

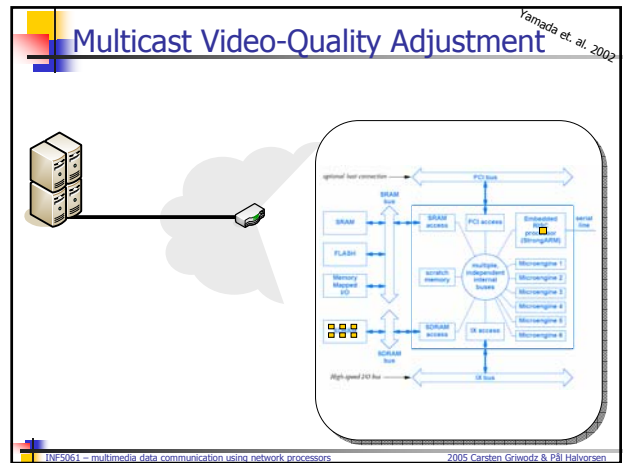
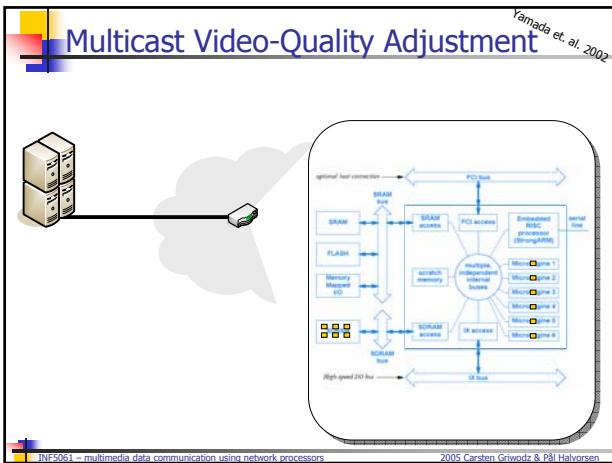
- Segmentation of MPEG-2 data
  - slice = 16 bit high stripes
  - macroblock = 16 x 16 bit square
    - four 8 x 8 luminance
    - two 8 x 8 chrominance
  - DCT transformed with coefficients sorted in ascending order
- Data packetization for video filtering
  - 720 x 576 pixels frames and 30 fps
  - 36 "slices" with 45 macroblocks per frame
  - Each slice = one packet
  - 8 Mbps stream → ~7Kb per packet

Sequence layer: GoP, GoP, ..., GoP  
GoP layer: I B P B P B B P B B  
Picture layer: Slice  
Slice Layer: Y1 Y2 Cb Cr  
MacroBlock Layer: Y1 Y2 Cb Cr  
Block Layer: DCT coefficient

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- ### Multicast Video-Quality Adjustment
- Yamada et. al. 2002
- Low-pass filter on IXP1200
    - parallel execution on 200MHz StrongARM and microengines
    - 24 MB DRAM devoted to StrongARM only
    - 8 MB DRAM and 8 MB SRAM shared
  - test-filtering program on a regular PC determined work-distribution
    - 75% of data from the block layer
    - 56% of the processing overhead is due to DCT
  - five step algorithm:
    - StrongArm receives packet → copy to shared memory area
    - StrongARM process headers and generate macroblocks (in shared memory)
    - microengines read data and information from shared memory and perform quality adjustments on each block
    - StrongARM checks if the last macroblock is processed (if not, go to 2)
    - StrongARM rebuilds packet
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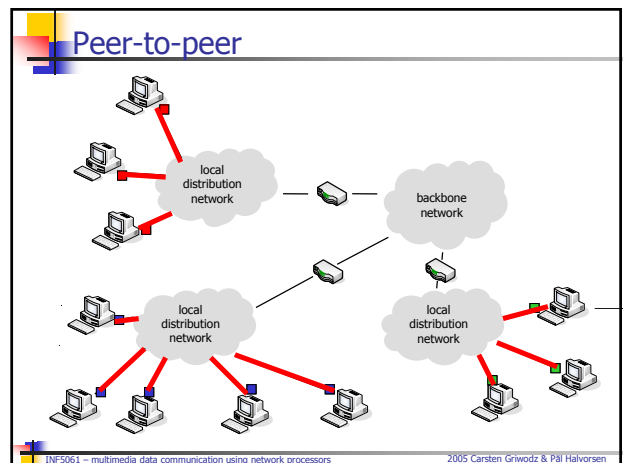
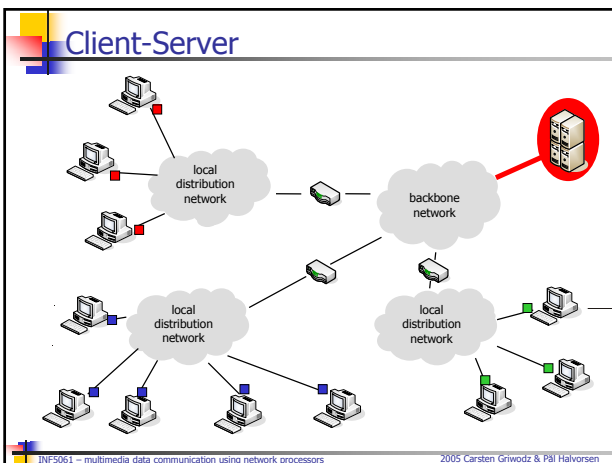




- ### Multicast Video-Quality Adjustment
- Yamada et. al. 2002
- Evaluation – three scenarios tested
    - StrongARM only → 550 kbps
    - StrongARM + 1 microengine → 350 kbps
    - StrongARM + all microengines → 1350 kbps
  - achieved **real-time** transcoding not enough for practical purposes, but **distribution of workload is nice**
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### Example: Booster Boxes

slide content and structure mainly from the NetGames 2002 presentation by Bauer, Rooney and Scotton



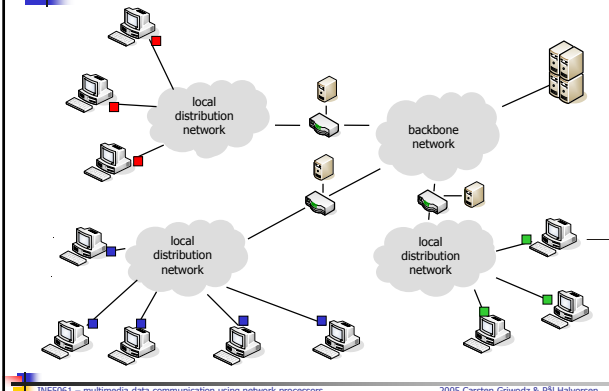
## IETF's Middleboxes

- Middlebox
  - network intermediate device that implements middlebox services
  - a middlebox function requires application specific intelligence
- Examples
  - policy based packet filtering (a.k.a. firewall)
  - network address translation (NAT)
  - intrusion detection
  - load balancing
  - policy based tunneling
  - IPsec security
  - ...
- RFC3303 and RFC3304
  - From traditional middleboxes
    - Embed application intelligence within the device
  - To middleboxes supporting the MIDCOM protocol
    - Externalize application intelligence into MIDCOM agents

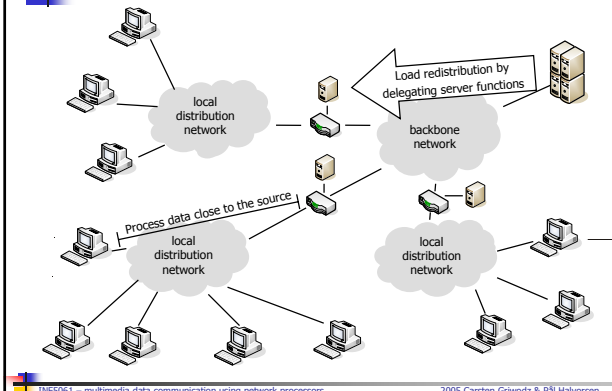
## Booster Boxes

- **Booster Boxes  $\approx$  Middleboxes**
  - attached directly to ISPs' access routers
  - less generic than, e.g., firewalls or NAT
- Assist distributed event-driven applications
  - improve scalability of client-server and P2P applications
- Application-specific code: **"Boosters"**

## Booster boxes



## Booster boxes

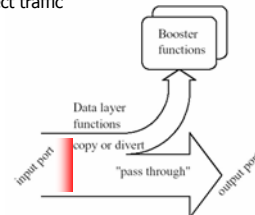


## Booster Box

- Application-specific code
  - Caching on behalf of a server
    - Non-real time information is cached
    - Booster boxes answer on behalf of servers
  - Aggregation of events
    - Information from two or more clients within a time window is aggregated into one packet
  - Intelligent filtering
    - Outdated or redundant information is dropped
  - Application-level routing
    - Packets are forward based on
      - Packet content
      - Application state
      - Destination address

## Architecture

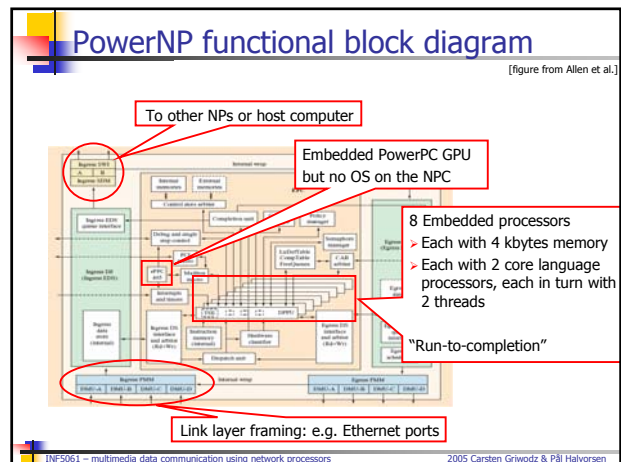
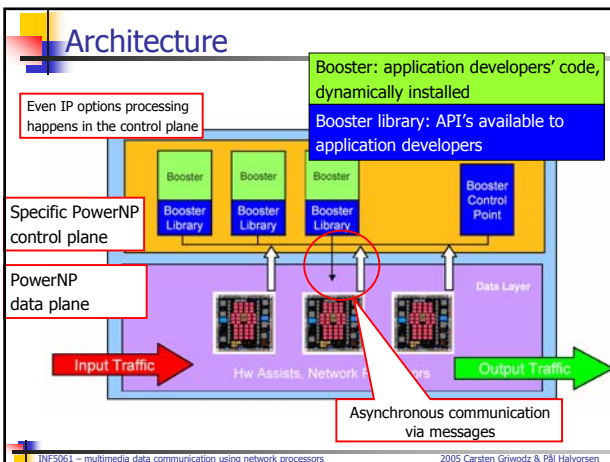
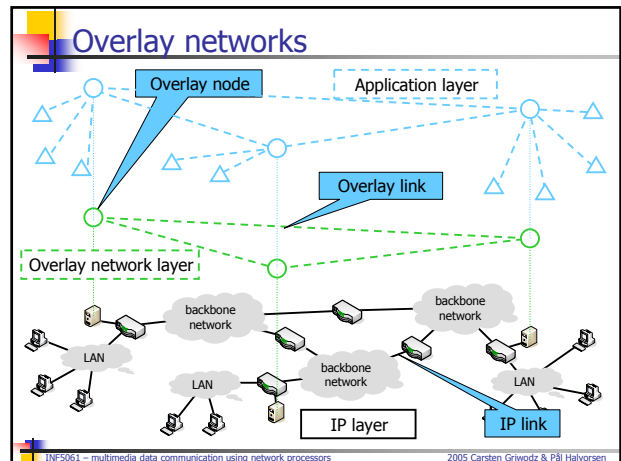
- Data Layer
  - behaves like a layer-2 switch for the bulk of the traffic
  - copies or diverts selected traffic
  - IBM's booster boxes use the packet capture library ("pcap") filter specification to select traffic



## Architecture

- Booster layer
  - Booster
    - Application-specific code
    - Executed either on the host CPU or the network processor
  - Library
    - Boosters can call the data-layer operation
- Generates a QoS-aware Overlay Network (Booster Overlay Network - BON)

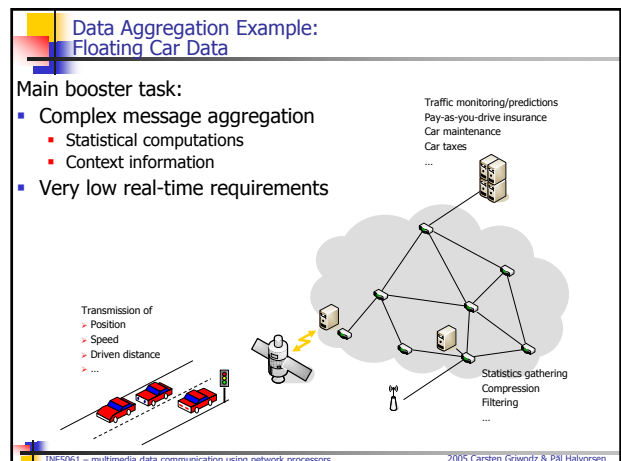
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## Intel IXP vs. IBM NP

- Difference between IBM NPs and IXP
  - IXP advantage
    - General purpose processor on the card
    - Operating system on the card
  - IXP disadvantage
    - Higher memory consumption for pipelining
    - Larger overhead for communication with host machine

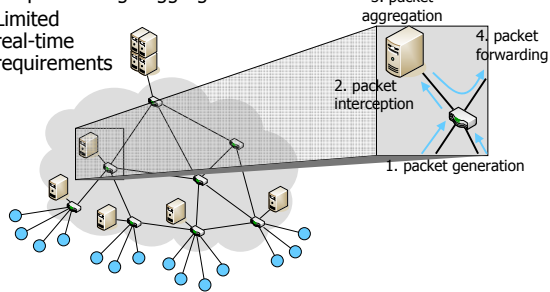
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## Interactive TV Game Show

Main booster task:

- Simple message aggregation
- Limited real-time requirements



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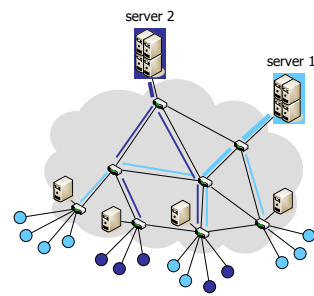
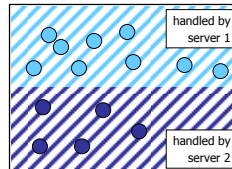
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## Game with large virtual space

Main booster task:

- Dynamic server selection
  - based on current in-game location
  - Require application-specific processing

Virtual space



- High real-time requirements

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## Summary

- Scalability
  - by application-specific knowledge
  - by network awareness
- Main mechanisms
  - Caching on behalf of a server
  - Aggregation of events
  - Attenuation
  - Intelligent filtering
  - Application-level routing
- Application of mechanism depends on
  - Workload
  - Real-time requirements

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## Some References

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