# UNIVERSITY OF OSLO <br> Faculty of mathematics and natural sciences 

Exam in: $\quad$ MAT3110/MAT4110 - Introduction to numerical analysis
Day of examination: 19 January 2021
Examination hours: 09:00-13:00
This problem set consists of 7 pages.
Appendices: None
Permitted aids: All written aids

Please make sure that your copy of the problem set is complete before you attempt to answer anything.

## Note:

- There are in total 11 subproblems (1, 2a, 2b,...), and you can get 5-10 points for each sub-problem, for a total of 100 points.
- All answers must be justified.


## Problem 1 Root finding

Let $f(x)=\cos (x)-x$. This function has a single root $x_{0}$ somewhere in $[0,1]$, and we wish to compute it.

## 1a (10 points)

Perform two steps with both the bisection method and Newton's method. Justify your choice of starting values.

## 1b (10 points)

Which of the two methods can we expect to be the most accurate after several iterations? Justify your answer.

## Solution:

1a
For the bisection method we choose $x_{0}=0, x_{1}=1$. Then $f\left(x_{0}\right)=1>0$ and $f\left(x_{1}\right)=\cos (1)-1<0$. Since $f$ is continuous, it has a zero in $\left(x_{0}, x_{1}\right)$, and the bisection method will be able to find it. We compute

$$
x_{2}=\frac{x_{0}+x_{1}}{2}=\frac{1}{2}
$$

and note that $f\left(x_{2}\right)=\cos (1 / 2)-1 / 2>0$, and therefore the new interval will be $\left(x_{2}, x_{1}\right)$. We finally get

$$
x_{3}=\frac{x_{1}+x_{2}}{2}=\frac{3}{4} .
$$

For Newton's method, we note that $f^{\prime}(x)=-\sin (x)-1$, which is nonzero in $[0,1]$. Hence, as long as the iteration stays in $[0,1]$, the method will converge. We set e.g. $x_{0}=0$ and get

$$
\begin{aligned}
& x_{1}=x_{0}-\frac{f\left(x_{0}\right)}{f^{\prime}\left(x_{0}\right)}=-\frac{1}{-1}=1 \\
& x_{2}=x_{1}-\frac{f\left(x_{1}\right)}{f^{\prime}\left(x_{1}\right)}=1-\frac{\cos (1)-1}{-\sin (1)-1} \approx 0.7504
\end{aligned}
$$

## 1b

The bisection method converges linearly, while Newton's method converges quadratically, so we can expect Newton's method to be the most accurate.

## Problem 2 Polynomial interpolation (10 points)

Let $f:[0,2] \rightarrow \mathbb{R}$ be a given function and let $n \in \mathbb{N}$. We wish to interpolate $f$ using an $n$-th order polynomial $p$.

- Explain how we should do this in order to minimize the maximal error $\|f-p\|_{C([0,2])}=\sup _{x \in[0,2]}|f(x)-p(x)|$.
- Give an estimate of $\|f-p\|_{C([0,2])}$.

Solution: Assume first that we are on the interval $[-1,1]$, and let $x_{0}, \ldots, x_{n} \in[-1,1]$ be distinct interpolation points. We let $p$ interpolate $f$ over these points:

$$
p(x)=\sum_{k=0}^{n} f\left(x_{k}\right) \prod_{\substack{l=0, \ldots, n \\ l \neq k}} \frac{x-x_{l}}{x_{k}-x_{l}}
$$

The basic error estimate is

$$
\|f-p\|_{C([-1,1])} \leqslant \frac{\left\|f^{(n+1)}\right\|_{C([-1,1])}}{(n+1)!}\left\|w_{n}\right\|_{C([-1,1])}
$$

where $w_{n}(x)=\prod_{k=0}^{n}\left(x-x_{k}\right)$. The term $\left\|w_{n}\right\|_{C([-1,1])}$ is the least possible when $x_{0}, \ldots, x_{n}$ are chosen as the Chebysheff points, yielding $\left\|w_{n}\right\|_{C([-1,1])}=2^{-n}$, and therefore

$$
\|f-p\|_{C([-1,1])} \leqslant \frac{\left\|f^{(n+1)}\right\|_{C([-1,1])}}{2^{n}(n+1)!}
$$

Any other choice of interpolation points will yield a larger right-hand side.

To transform this analysis to the interval $[0,2]$, it is enough to note that the two intervals are of the same length, and that translating a function does not change its norm, so the same results apply:

$$
\|f-p\|_{C([0,2])} \leqslant \frac{\left\|f^{(n+1)}\right\|_{C([0,2])}}{2^{n}(n+1)!}
$$

## Problem 3 Polynomial interpolation

Let $f:[0,1] \rightarrow \mathbb{R}$ be the function $f(x)=\cos (2 x)-e^{x}$. For some $n \in \mathbb{N}$, let $p$ be the $n$-th order polynomial which interpolates $f$ over the uniform grid $0,1 / n, \ldots, 1$.

## 3a (10 points)

Prove that $\|f-p\|_{C([0,1])} \rightarrow 0$ as $n \rightarrow \infty$.
$\left(\right.$ Here, $\|f-p\|_{C([0,1])}=\sup _{x \in[0,1]}|f(x)-p(x)|$.)

## 3b (10 points)

How large must $n$ be in order to guarantee that $\|f-p\|_{C([0,1])} \leqslant 10^{-10}$ ?
Hint: In this problem you might (or might not) need Stirling's approximation:

$$
m!\geqslant m^{m} e^{-m}
$$

Solution: The basic error estimate is: For every $x \in[0,1]$ there is some $\xi \in[0,1]$ such that

$$
|f(x)-p(x)| \leqslant \frac{\left|f^{(n+1)}(\xi)\right|}{(n+1)!} \prod_{k=0}^{n}\left|x-x_{k}\right|
$$

We have

$$
\left|f^{(m)}(\xi)\right| \leqslant\left|\frac{d^{m}}{d \xi^{m}} \cos (2 \xi)\right|+\left|\frac{d^{m}}{d \xi^{m}} e^{\xi}\right| \leqslant 2^{m}+e
$$

Moreover, $\left|x-x_{k}\right| \leqslant 1$, so we get

$$
|f(x)-p(x)| \leqslant \frac{2^{n+1}+1}{(n+1)!} \quad \forall x \in[0,1]
$$

Using Stirling's formula we get

$$
\|f-p\|_{C([0,1])} \leqslant \frac{2^{n+1}+1}{(n+1)!} \leqslant \frac{e^{n+1}\left(2^{n+1}+1\right)}{(n+1)^{n+1}}
$$

## 3a

It is clear that the expression above converges to zero as $n \rightarrow \infty$.

## 3b

Testing different values of $n$ shows that $n=29$ gives an upper bound of $\approx 4.3 \times 10^{-11}$.

Alternative solution: If, say, $x \in\left[x_{m}, x_{m+1}\right]$ then

$$
\prod_{k=0}^{n}\left|x-x_{k}\right|=\frac{1}{n^{n+1}} \prod_{k=0}^{n}|x n-k| \leqslant \frac{(m+1)!(n-m)!}{n^{n+1}} \leqslant \frac{(n+1)!}{n^{n+1}}
$$

where the last inequality follows from $1 \leqslant\binom{ n+1}{m+1}=\frac{(n+1)!}{(m+1)!(n-m)!}$. We get

$$
|f(x)-p(x)| \leqslant \frac{\left\|f^{(n+1)}\right\|}{(n+1)!} \frac{(n+1)!}{n^{n+1}} \leqslant \frac{2^{n+1}+1}{n^{n+1}}
$$

3a
It is clear that $\frac{2^{n+1}+1}{n^{n+1}} \rightarrow 0$ as $n \rightarrow \infty$.
3b
With this improved estimate we find that $n=12$ gives an upper bound $\sim 7.66 \times 10^{-11}$.

## Problem 4 QR factorization

Let $A, Q$ and $R$ be the matrices

$$
A=\left(\begin{array}{cc}
0 & 1 \\
\sqrt{2} & 3 \sqrt{2} \\
0 & 1
\end{array}\right), \quad R=\sqrt{2}\left(\begin{array}{cc}
1 & 3 \\
0 & 1 \\
0 & 0
\end{array}\right), \quad Q=\frac{1}{\sqrt{2}}\left(\begin{array}{ccc}
0 & 1 & 1 \\
\sqrt{2} & 0 & 0 \\
0 & 1 & -1
\end{array}\right)
$$

Note that $A=Q R$ (you don't have to show this).

## $4 \mathrm{a} \quad$ (5 points)

Explain what it means that $Q R$ is the QR factorization of $A$. Justify your answer.

## 4b (10 points)

Find the least squares solution of the equation

$$
A x=b, \quad \text { where } b=\left(\begin{array}{l}
1 \\
2 \\
3
\end{array}\right)
$$

## Solution:

## 4a

A QR factorization consists of an orthogonal matrix $Q$ and an upper triangular matrix $R$ (with 1's as its first nonzero element in each row, if the factorization is in normal form). It is straightforward to see that $Q^{\top} Q=I$, so $Q$ is orthogonal, and that $R$ is upper triangular.

## 4b

We wish to minimize $\|A x-b\|=\left\|R x-Q^{\top} b\right\|$. Write

$$
R=\binom{R_{1}}{0}, \quad Q^{\boldsymbol{\top}} b=\binom{c_{1}}{c_{2}}, \quad R_{1}=\sqrt{2}\left(\begin{array}{ll}
1 & 3 \\
0 & 1
\end{array}\right), \quad c_{1} \in \mathbb{R}^{2}, c_{2} \in \mathbb{R}
$$

Then $\|A x-b\|^{2}=\left\|R x-Q^{\top} b\right\|^{2}=\left\|R_{1} x-c_{1}\right\|^{2}+\left\|c_{2}\right\|^{2}$, so we need to minimize the first term; to this end, we solve $R_{1} x=c_{1}$. We compute

$$
c_{1}=\binom{2}{2 \sqrt{2}} \quad \Rightarrow \quad x=R_{1}^{-1} c_{1}=\binom{\sqrt{2}-6}{2} .
$$

## Problem 5 SVD (10 points)

Compute the singular value decomposition (SVD) of

$$
A=\left(\begin{array}{ll}
1 & 2 \\
3 & 6
\end{array}\right) .
$$

Hint: You may use the fact that one of the eigenpairs of the normal matrix $A^{\top} A$ is $\lambda_{1}=50, \mathbf{v}_{1}=\frac{1}{\sqrt{5}}\binom{1}{2}$.

Solution: We see that $A$ is non-invertible, so $A^{\top} A$ must also be non-invertible, whence the second eigenvalue is $\lambda_{2}=0$. The second eigenvector is chosen such that $V=\left(\begin{array}{ll}\mathbf{v}_{1} & \mathbf{v}_{2}\end{array}\right)$ is orthogonal; this is achieved by letting $\mathbf{v}_{2}=\frac{1}{\sqrt{5}}\binom{2}{-1}$. We get the two singular values

$$
\sigma_{1}=\sqrt{50}=5 \sqrt{2}, \quad \sigma_{2}=0
$$

Setting $S=\operatorname{diag}\left(\sigma_{1}, \sigma_{2}\right)$ we want to find an orthogonal matrix $U$ such that $A=U S V^{\top}$, or $U S=A V$. Writing $U=\left(\begin{array}{ll}\mathbf{u}_{1} & \mathbf{u}_{2}\end{array}\right)$, we have

$$
U S=\left(\begin{array}{ll}
\sigma_{1} \mathbf{u}_{1} & 0
\end{array}\right) \quad \Rightarrow \quad \mathbf{u}_{1}=\frac{1}{\sigma_{1}} A \mathbf{v}_{1}=\frac{1}{\sqrt{10}}\binom{1}{3} .
$$

Finally, we let $\mathbf{u}_{2}$ be such that $U$ is orthogonal: $\mathbf{u}_{2}=\frac{1}{\sqrt{10}}\binom{3}{-1}$. Thus, $A=U S V^{\top}$ with

$$
U=\frac{1}{\sqrt{10}}\left(\begin{array}{cc}
1 & 3 \\
3 & -1
\end{array}\right), \quad S=\left(\begin{array}{cc}
5 \sqrt{2} & \\
& 0
\end{array}\right), \quad V=\frac{1}{\sqrt{5}}\left(\begin{array}{cc}
1 & 2 \\
2 & -1
\end{array}\right) .
$$

## Problem 6

We wish to approximate the integral $I(f)=\int_{0}^{20} f(x) d x$ of a function $f$.

## 6 a (5 points)

If we wish to approximate $I(f)$ using an 5 -point quadrature rule, which quadrature rule should we choose to make the error as small as possible? Justify your answer.

## 6b (10 points)

Recall that the Gauss quadrature of order 3 on the interval $[-1,1]$ is

$$
\begin{equation*}
\int_{-1}^{1} g(x) d x \approx f(-\sqrt{1 / 3})+f(\sqrt{1 / 3}) \tag{1}
\end{equation*}
$$

Write down the composite integration rule over $N=2$ subintervals which approximates $I(f)$. Use the quadrature rule (1) in the composite method.

## Solution:

6a
The $n$-point Gauss quadrature rule has order $2 n-1$, which is the largest possible. We should therefore use the 3 -point Gauss quadrature rule.
$6 b$
Translating the interval $[-1,1]$ to $[0,10]$ gives quad. points and weights

$$
x_{0}=5-5 / \sqrt{3}, \quad x_{1}=5+5 / \sqrt{3}, \quad w_{0}=w_{1}=5
$$

and on the interval $[10,20]$

$$
x_{2}=15-5 / \sqrt{3}, \quad x_{3}=15+5 / \sqrt{3}, \quad w_{2}=w_{3}=5
$$

Thus, the composite quadrature rule is

$$
I(f) \approx 5(f(5-5 / \sqrt{3})+f(5+5 / \sqrt{3})+f(15+5 / \sqrt{3})+f(15+5 / \sqrt{3}))
$$

## Problem 7 Runge-Kutta method (10 points)

Consider the ODE

$$
\left\{\begin{array}{l}
x^{\prime}(t)=f(x(t), t) \\
x(0)=x_{0}
\end{array}\right.
$$

where $f$ is a given smooth function, and the Runge-Kutta method

$$
\begin{aligned}
k & =f\left(y_{n}+h k / 2, t_{n}+h / 2\right) \\
y_{n+1} & =y_{n}+h k .
\end{aligned}
$$

Set $f(x, t)=\lambda x$ for some $\lambda \in \mathbb{C}$ with $\operatorname{Re}(\lambda)<0$. Find the stability function of this method, and determine whether the method is unconditionally stable or not.

Hint: If you are unable to determine stability, it's enough to insert $h \lambda=$ 1, 10, 100 in the stability function and conclude based on that.

Solution: We insert $f(x, t)=\lambda x$ and get

$$
\begin{aligned}
& k=\lambda\left(y_{n}+h k / 2\right) \quad \Leftrightarrow \quad k=\frac{h \lambda}{1-h \lambda / 2} y_{n} \\
& y_{n+1}=y_{n}+h k=y_{n}\left(1+\frac{h \lambda}{1-h \lambda / 2}\right)=y_{n} R(h \lambda)
\end{aligned}
$$

where $R(z)=1+\frac{z}{1-z / 2}=\frac{1+z / 2}{1-z / 2}=\frac{2+z}{2-z}$. If $\operatorname{Re}(z)<0$ then $|2+z|<|2-z|$, and therefore $|R(z)|<1$ for all such $z$. We conclude that the method is unconditionally stable.

