

Summary

- Highest weight theory (cont'd)
 - Verma module
- Categorical structure

• Highest weight theory

(Highest wght vecs, Verma modules)

(π^w, V^w) simple quot. of Verma mod.

Prop. $w \in P_{++} \Rightarrow \dim V^w < \infty$

Step 1 $V^w = U W$ ($\Rightarrow \pi_{\mathfrak{E}\alpha}^w, \pi_{\mathfrak{F}\alpha}^w$ loc. nilpot.)
 $\mathfrak{S}\alpha$ -inv fin dim. \hookrightarrow RHS is \mathfrak{g} -inv

$\therefore \mathfrak{g} \otimes W \rightarrow V^w$ is $\mathfrak{K}\alpha$ -equivar.

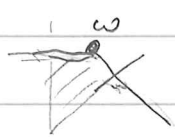
$w \in P_{++} \Rightarrow W = \langle \mathfrak{F}\alpha^k v : k=0, \dots, w(H\alpha) \rangle$
 is $\mathfrak{S}\alpha$ -inv.

Step 2 $\text{Supp}(V^w) = \{ \lambda : V_\lambda^w \neq 0 \}$ is W -inv.

$\therefore \mathfrak{S}\alpha$ -inv. by Step 1

Step 3 $\text{Supp}(V^w)$ is finite.

$\therefore \text{Supp}(V^w) \cap P_{++} \subset \{ w - \sum_{i=1}^n k_i \alpha_i : k_i \geq 0 \}$

E.g.  A_2 is $\begin{cases} \text{finite} \\ \text{fund. dom for } W \end{cases}$

• Categorical structure

Overall motivation

G (cpt) grp \rightsquigarrow "monoidal category" & "fiber functor"
 \mathfrak{e} \mathfrak{F}

s.t. $(\mathfrak{e}, \mathfrak{F})$ can recover G
 (abstract characterization of $(\mathfrak{e}, \mathfrak{F})$)

Ex. $G \rightsquigarrow \mathcal{C} : \begin{cases} \text{category of finite } G\text{-sets} \\ (X, \alpha: G \curvearrowright X), (X \times Y, \alpha \times \beta) \\ F(X, \alpha) = X \end{cases}$

Ex. $G \rightsquigarrow \mathcal{C} : \begin{cases} \text{category of fin. dim unitary} \\ \text{reps } (H, \pi: G \rightarrow U(H)), (H \otimes H', \pi \otimes \pi') \\ F(H, \pi) = H \end{cases}$

How does "abstract characterization" work?

Ex. Galois theory

K field, \bar{K} alg. closure of K .

$\mathcal{C} =$ category of finite étale algs / K .

(A étale $\equiv A \simeq L_1 \times \dots \times L_m, L_i$ fin. sep / K)

$F(A) = \text{Hom}_{K\text{-alg}}(A, \bar{K}), F(A \otimes B) = F(A) \times F(B)$

$F(A)$ has left action of $\text{Gal}(\bar{K}/K) = \text{Aut}_{K\text{-alg}}(\bar{K})$

in fact $\text{Gal}(\bar{K}/K) = \text{"Aut}(F)$ ", ($\mathcal{C} \simeq \text{Gal-Sets}$)

Ex. \rightarrow 3-dim QFT $A = (A(0))_{0 \in \mathbb{R}^3}$ fam. of algs.

$\rightsquigarrow \mathcal{C} : \text{"rep. of } A \text{"}$

(minor cond. on A) $\rightsquigarrow \mathcal{C} \simeq \text{Rep } G$

Doplicher-Roberts for some G .

Ex. étale fund. grp. of scheme (Grothendieck)

$X : (\text{conn, loc. Noetherian}) \quad x : \text{Spec}(\bar{K}) \rightarrow X$

$\mathcal{C} : \text{cat. of finite étale schemes / } X$

$F(Y) = \text{"fiber of } Y \rightarrow X \text{ over } x \text{"}$

$$= \left\{ \begin{array}{ccc} \text{Spec}(\bar{K}) & \rightarrow & Y \\ \downarrow & & \downarrow \\ x & \rightarrow & X \end{array} \right\}$$

Category \mathcal{C} is given by

- objects $X, Y \dots \in \text{Ob}(\mathcal{C})$
 - morphisms $S, T \dots \in \text{Mor}(X, Y) \quad (X, Y \in \text{Ob}(\mathcal{C}))$
- $\text{Id}_X \in \text{End}(X) = \text{Mor}(X, X)$
 $(S_1, S_2)S_3 = S_1(S_2S_3), \quad \text{Id}_Y S = S = S \text{Id}_X.$

Ex. $\mathcal{C} = (G\text{-Sets})$

$\text{Mor}((X, \alpha), (Y, \beta)) = \{f : X \rightarrow Y, \beta \circ f = f \circ \alpha\}$

Ex. $\mathcal{C} = \text{Rep } G$

$\text{Mor}((H^1, \pi^1), (H^2, \pi^2)) = \{T : H^1 \rightarrow H^2 \text{ lin.} \\ \pi^2 \circ T = T \circ \pi^1\}$

Functor $F : \mathcal{C} \rightarrow \mathcal{C}'$ is given by

- maps $F : \text{Ob}(\mathcal{C}) \rightarrow \text{Ob}(\mathcal{C}')$
- $\text{Mor}(X, Y) \rightarrow \text{Mor}(F(X), F(Y))$
- s.t. $F(\text{Id}_X) = \text{Id}_{F(X)}, \quad F(S \circ T) = F(S) \circ F(T)$

Natural transform $\phi : F \rightarrow F'$ is given by

morphisms $\phi_X : F(X) \rightarrow F'(X)$ s.t.

$F, F' : \mathcal{C} \rightarrow \mathcal{C}'$

$$\begin{array}{ccc} F(X) & \xrightarrow{\phi_X} & F'(X) \\ F(T) \downarrow & \circlearrowleft & \downarrow F'(T) \\ F(Y) & \xrightarrow{\phi_Y} & F'(Y) \end{array} \quad \forall T \in \text{Mor}(X, Y)$$

Nat. iso. : all ϕ_X are invertible.

Ex.

$\mathcal{C} = (G\text{-Sets}_f), \quad \mathcal{C}' = (\text{Sets}_f) \quad F(X, \alpha) = X.$

$g \in G \mapsto \phi^g \in \text{Aut}(F)$ by $\phi^g_{(X, \alpha)} = \alpha_g$ as

map $F(X, \alpha) \rightarrow F(X, \alpha)$

Ex. $\mathcal{C} = \text{Vec}_K = \mathcal{C}' \quad F(V) = V, \quad F'(V) = V^{**}$

$\phi : F \rightarrow F'$ by $\phi_V(x) (\zeta) = \zeta(x)$ for $\zeta \in V^*$

