

## Entropies & entropic inequalities

## Recap:

• VN entropy:  $H(\rho) = -\text{Tr}[\rho \log(\rho)]$

↪ operational interpretation: Compression of  $q$ -states

• Relative entropy:  $D(\rho \parallel \sigma) = \begin{cases} \text{Tr}[\rho(\log(\rho) - \log(\sigma))] & \text{if } \ker(\sigma) \subseteq \ker(\rho) \\ +\infty & \text{else.} \end{cases}$

Lemma:  $D(\rho \parallel \sigma) = \frac{1}{\ln(2)} \lim_{\epsilon \searrow 0} \frac{1 - \text{Tr}[\rho^{1-\epsilon} \sigma^\epsilon]}{\epsilon}$

Central inequality: Data-processing inequality

$$D(T(\rho) \parallel T(\sigma)) \leq D(\rho \parallel \sigma)$$

$\forall q$ -channels  $T$ .

## Immediate consequences of data-processing:

(1) Joint convexity of relative entropy

$$\rho_1, \rho_2, \sigma_1, \sigma_2 \in \mathcal{D}(H) \quad , \quad \lambda \in [0, 1]$$

Define:  $\rho = (1-\lambda)\rho_1 \otimes |1\rangle\langle 1| + \lambda\rho_2 \otimes |2\rangle\langle 2| \in \mathcal{D}(H \otimes \mathbb{C}^2)$

$$\sigma = (1-\lambda)\sigma_1 \otimes |1\rangle\langle 1| + \lambda\sigma_2 \otimes |2\rangle\langle 2| \in \mathcal{D}(H \otimes \mathbb{C}^2)$$

Then:  $\text{Tr}_2[\rho] = (1-\lambda)\rho_1 + \lambda\rho_2$

$$\text{Tr}_2[\sigma] = (1-\lambda)\sigma_1 + \lambda\sigma_2$$

$$D((1-\lambda)\rho_1 + \lambda\rho_2 \parallel (1-\lambda)\sigma_1 + \lambda\sigma_2) = D(\text{Tr}_2[\rho] \parallel \text{Tr}_2[\sigma])$$

Data-processing

$$\leq D(\rho \parallel \sigma) = (1-\lambda)D(\rho_1 \parallel \sigma_1) + \lambda D(\rho_2 \parallel \sigma_2)$$

uses that

$$\log((1-\lambda)\rho_1 \otimes |1\rangle\langle 1| + \lambda\rho_2 \otimes |2\rangle\langle 2|)$$

$$= \log(1-\lambda) \mathbb{1}_H \otimes |1\rangle\langle 1| + \log(\lambda) \mathbb{1}_H \otimes |2\rangle\langle 2| \\ + \log(\rho_1) \otimes |1\rangle\langle 1| + \log(\rho_2) \otimes |2\rangle\langle 2|$$

(2) Klein's inequality

$$D(\rho \parallel \sigma) \geq D(\text{Tr}[\rho] \parallel \text{Tr}[\sigma]) = 0.$$

↑  
Data-processing

~> Tomorrow: Strengthening of this inequality.

### (3) Concavity of vN entropy

Note that  $H(\rho) = \log(d) - D(\rho \| \frac{1}{d})$

$$\begin{aligned} \text{Therefore: } H((1-\lambda)\rho + \lambda\sigma) &= \log(d) - D((1-\lambda)\rho + \lambda\sigma \| (1-\lambda)\frac{1}{d} + \lambda\frac{1}{d}) \\ &\geq \log(d) - (1-\lambda)D(\rho \| \frac{1}{d}) - \lambda D(\sigma \| \frac{1}{d}) \\ &= (1-\lambda)H(\rho) + \lambda H(\sigma). \end{aligned}$$

□

(4) Subadditivity of vN entropy

$$\rho_{AB} \in \mathcal{D}(H_A \otimes H_B) \Rightarrow \ker(\rho_A \otimes \rho_B) \stackrel{\text{spectral theorem}}{=} \text{span} \{ |v\rangle \otimes |w\rangle \in H_A \otimes H_B : |v\rangle \in \ker(\rho_A) \text{ or } |w\rangle \in \ker(\rho_B) \}$$

$$\subseteq \ker(\rho_{AB})$$

if  $|v\rangle \in \ker(\rho_A)$ , then  $\langle v | \rho_A | v \rangle = \sum_{i=1}^{\dim(H_B)} \underbrace{\langle v | \otimes \langle i | }_{\geq 0} \rho_{AB} (|v\rangle \otimes |i\rangle) = 0$

$$\Rightarrow \rho_{AB} |v\rangle \otimes |i\rangle = 0 \quad \forall i \in \{1, \dots, \dim(H_B)\}$$

$$\Rightarrow \rho_{AB} |v\rangle \otimes |w\rangle = 0 \quad \forall |w\rangle \in H_B$$

$$\Rightarrow |v\rangle \otimes |w\rangle \in \ker(\rho_{AB}) \quad \forall |w\rangle \in H_B$$

Same argument works if  $|w\rangle \in \ker(\rho_B)$ .

$$D(\rho_{AB} \parallel \rho_A \otimes \rho_B) = \text{Tr} \left[ \rho_{AB} \left( \log(\rho_{AB}) - \log(\rho_A \otimes \rho_B) \right) \right]$$

$$= \log(\rho_A) \otimes \Pi_{\text{Im}(\rho_B)} + \Pi_{\text{Im}(\rho_A)} \otimes \log(\rho_B)$$

$$= -H(\rho_{AB}) + H(\rho_A) + H(\rho_B)$$

$$\geq 0 \quad \text{by Klein's inequality.}$$

□

(5) Strong subadditivity of vN entropy

As before, compute

$$\begin{aligned} D(\rho_{ABC} \parallel \frac{1_{H_A}}{d_A} \otimes \rho_{BC}) &= \text{Tr} \left[ \rho_{ABC} \left( \log(\rho_{ABC}) - \log \left( \frac{1_{H_A}}{d_A} \otimes \rho_{BC} \right) \right) \right] \\ &= -\log(d_A) 1_{H_A} \otimes \text{Tr}_{\text{im}(\rho_{BC})} \\ &\quad + 1_{H_A} \otimes \log(\rho_{BC}) \\ &= -H(\rho_{ABC}) + \log(d_A) + H(\rho_{BC}) \end{aligned}$$

$$D(\rho_{AB} \parallel \frac{1_{H_A}}{d_A} \otimes \rho_B) = -H(\rho_{AB}) + \log(d_A) + H(\rho_B)$$

Data-processing inequality

$$\Rightarrow D(\rho_{AB} \parallel \frac{1_{H_A}}{d_A} \otimes \rho_B) \leq D(\rho_{ABC} \parallel \frac{1_{H_A}}{d_A} \otimes \rho_{BC})$$

$\Rightarrow$

~~$H(\rho_{ABC})$~~

$$H(\rho_{ABC}) + H(\rho_B) \leq H(\rho_{AB}) + H(\rho_{BC})$$



Exercise 3: (conditional entropy)

For  $\rho_{AB} \in \mathcal{D}(\mathcal{H}_A \otimes \mathcal{H}_B)$  define

$$H(A|B)_{\rho_{AB}} = H(\rho_{AB}) - H(\rho_B)$$

(1) "Classically, discarding systems does not increase the entropy"

Recall product rule:  $P_{AB}(x,y) = P_B(y) P_{A|B}(x|y)$

$$\begin{aligned} H(\rho_{AB}) &= - \sum_{x,y} P_B(y) P_{A|B}(x|y) \log(P_B(y) P_{A|B}(x|y)) \\ &= - \sum_{x,y} P_B(y) P_{A|B}(x|y) (\log(P_B(y)) + \log(P_{A|B}(x|y))) \\ &= H(\rho_B) + \underbrace{\sum_y P_B(y) H(P_{A|B}(\cdot|y))}_{\geq 0} \geq H(\rho_B) \end{aligned}$$

$$\Rightarrow H(A|B)_{\rho_{AB}} \geq 0.$$

□



(2) "Quantumly, discarding systems does increase the entropy in general"

$$\rho_{AB} = \omega_d \in \mathcal{D}(\mathbb{C}^d \otimes \mathbb{C}^d) \quad \text{normalized maximally ent. state.}$$

$$\omega_d \text{ is pure} \implies H(\omega_d) = 0$$

$$\rho_B = \text{Tr}_A[\omega_d] = \frac{1}{d} \implies H(\rho_B) = \log(d)$$

$$H(A|B)_{\rho_{AB}} = -\log(d) < 0.$$

□

(3) "Conditioning does not increase the vN entropy"

$$H(A|B, C)_{\mathcal{P}_{ABC}} = H(\mathcal{P}_{ABC}) - H(\mathcal{P}_{BC})$$

$$H(A|B)_{\mathcal{P}_{ABC}} = H(\mathcal{P}_{AB}) - H(\mathcal{P}_B)$$

$$\text{SSA} \Rightarrow H(\mathcal{P}_{ABC}) + H(\mathcal{P}_B) \leq H(\mathcal{P}_{AB}) + H(\mathcal{P}_{BC})$$

$$\Rightarrow \boxed{H(A|B, C)_{\mathcal{P}_{ABC}} \leq H(A|B)_{\mathcal{P}_{ABC}}}$$

## (4) Cool application of conditional entropy

Thm: (Quantum Shearer's inequality)

For  $t \in \mathbb{N}$  consider a family  $\mathcal{F} = \{F_1, \dots, F_K\}$  of subsets of  $\{1, \dots, n\}$  s.t.

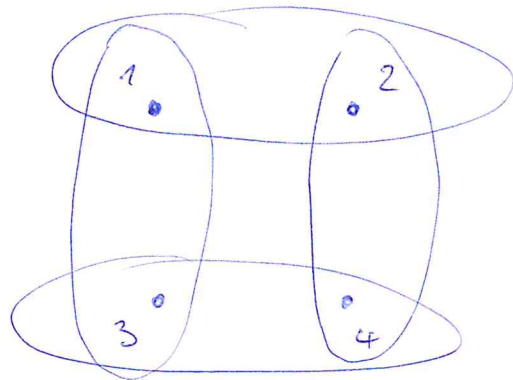
$$|\{k \in \{1, \dots, K\} : i \in F_k\}| = t \quad \forall i \in \{1, \dots, n\}.$$

Then:

$$H(\rho_{1,2,\dots,n}) \leq \frac{1}{t} \sum_{F \in \mathcal{F}} H(\rho_F)$$

for any  $\rho_{1,2,\dots,n} \in \mathcal{D}(H_1 \otimes \dots \otimes H_n)$ .

Example:



$$H(\rho_{1,2,3,4}) \leq \frac{1}{2} \left( H(\rho_{1,2}) + H(\rho_{2,4}) \right. \\ \left. + H(\rho_{3,4}) + H(\rho_{1,3}) \right)$$

Proof: Consider  $F \subseteq \{1, 2, \dots, n\}$  and denote its elements by  $(i_1, \dots, i_\ell)$   
 s.t.  $i_1 < i_2 < \dots < i_\ell$ .

$$\begin{aligned} \sum_{j=1}^{|F|} H(i_j | i_1, \dots, i_{j-1})_{\mathcal{P}_{1,2,\dots,n}} &= H(i_1)_{\mathcal{P}_{1,2,\dots,n}} + H(i_2 | i_1)_{\mathcal{P}_{1,2,\dots,n}} + H(i_3 | i_1, i_2)_{\mathcal{P}_{1,2,\dots,n}} + \dots \\ &= \cancel{H(i_1)} + \cancel{H(i_1, i_2)} - \cancel{H(i_1)} + \cancel{H(i_1, i_2, i_3)} - \cancel{H(i_1, i_2)} \\ &\quad + \dots \\ &= H(i_1, \dots, i_{|F|})_{\mathcal{P}_{1,2,\dots,n}} = H(\mathcal{P}_F). \end{aligned}$$

$$\frac{1}{t} \sum_{F \in \mathcal{F}} H(\mathcal{P}_F) = \frac{1}{t} \sum_{F \in \mathcal{F}} \sum_{j=1}^{|F|} H(i_j^F | i_1^F, \dots, i_{j-1}^F)_{\mathcal{P}_{1,2,\dots,n}}$$

Conditioning decreases entropy

$$\geq \frac{1}{t} \sum_{F \in \mathcal{F}} \sum_{j=1}^{|F|} H(i_j^F | 1, \dots, i_{j-1}^F)_{\mathcal{P}_{1,2,\dots,n}}$$

Every  $i_j^F$  occurs  $t$  times

$$= \frac{1}{t} \cdot t \sum_{i=1}^n H(i | 1, \dots, i-1)_{\mathcal{P}_{1,2,\dots,n}} = H(\mathcal{P}_{1,2,\dots,n})$$

□

Another example:  $\mathcal{F} = \mathcal{F}_k = \{ F \subseteq \{1, \dots, n\} : |F| = k \}$

$\leadsto$  all  $k$ -element subsets  $\implies t = \binom{n-1}{k-1} = \frac{k}{n} \binom{n}{k}$

Quantum Shearer's inequality  $\implies \frac{k}{n} H(\rho_{1, \dots, n}) \leq \frac{1}{\binom{n}{k}} \sum_{F \in \mathcal{F}_k} H(\rho_F)$

Application: Entropy production

Depolarizing channel:  $T_p = (1-p) \text{id} + p \text{Tr}[\cdot] \frac{1}{d}$

$$T_p^{\otimes n}(\rho) = \sum_{k=0}^n \sum_{F \in \mathcal{F}_k} (1-p)^k p^{n-k} \left( \bigotimes_{\ell \in F^c} \left( \frac{1}{d} \right)_\ell \otimes \rho_F \right)$$

How much entropy does  $T_p^{\otimes n}$  produce?

$$\underline{H(T_p^{\otimes n}(\rho))} \stackrel{\text{Concavity } n}{\geq} \sum_{k=0}^n \sum_{F \in \mathcal{F}_k} (1-p)^k p^{n-k} \underbrace{H\left(\bigotimes_{\ell \in F^c} \left(\frac{1}{d}\right)_\ell \otimes \rho_F\right)}_{=(n-k) \log(d) + H(\rho_F)}$$

Use

$$\begin{aligned} & \sum_{k=0}^n \binom{n}{k} (n-k) (1-p)^k p^{n-k} \\ &= \sum_{k=0}^n \binom{n-1}{k-1} n (1-p)^k p^{n-k} \\ &= n(1-p)^n + \left( \sum_{k=0}^{n-1} \binom{n-1}{k} (1-p)^k p^{n-k-1} \right) np \\ &\geq np \end{aligned}$$

$$\geq \sum_{k=0}^n \binom{n}{k} (1-p)^k p^{n-k} (n-k) \log(d) + \sum_{k=0}^n (1-p)^k p^{n-k} \underbrace{\sum_{F \in \mathcal{F}_k} H(\rho_F)}_{\geq \binom{n-1}{k-1} H(\rho)} \geq \underline{pn \log(d) + (1-p)H(\rho)}$$

$\geq \binom{n}{k} \frac{k}{n} H(\rho)$   
Q-Shearer

### Exercise 4: (Fannes' inequality)

Want to quantify the continuity of  $\rho \mapsto H(\rho)$ .

Thm: For  $\rho, \sigma \in \mathcal{D}(H)$  s.t.  $\|\rho - \sigma\|_1 \leq \frac{1}{2}$  we have

$$|H(\rho) - H(\sigma)| \leq \|\rho - \sigma\|_1 \log(\dim(H)) + \eta(\|\rho - \sigma\|_1) \quad (*)$$

where  $\eta(x) = -x \log x$ .

Two remarks:

(1)  $\eta(x) \rightarrow 0$  as  $x \rightarrow 0 \Rightarrow (*)$  quantifies the continuity of  $\sqrt{N}$  entropy

(2) RHS only depends on logarithm of  $\dim(H)$ .

$\Rightarrow$  Even if  $d \sim$  exponentially large, then  $(*)$  is still useful.



Proof: Set  $d = \dim(H)$

$$\rho = \sum_{i=1}^d \lambda_i |v_i\rangle\langle v_i| \quad \& \quad \sigma = \sum_{j=1}^d \mu_j |w_j\rangle\langle w_j| \quad \text{spectral decompositions}$$

s.t.  $\lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_d$  and  $\mu_1 \geq \mu_2 \geq \dots \geq \mu_d$ .

Step 1: Show that  $\|\rho - \sigma\|_1 \geq \sum_i |\lambda_i - \mu_i|$ .

Write  $\rho - \sigma = P - Q$  for  $P, Q \in \mathcal{B}(H)^+$  s.t.  $PQ = 0$  (Jordan-Hahn decomposition)

$$\Rightarrow \|\rho - \sigma\|_1 = \text{Tr}[(\rho - \sigma)^+] = \text{Tr}[P] + \text{Tr}[Q].$$

Define  $X := P + \rho = Q + \sigma \Rightarrow P + Q = 2X - \rho - \sigma$

$\leadsto$  Spectral decomposition:  $X = \sum_{k=1}^d \tau_k |u_k\rangle\langle u_k|$

s.t.  $\tau_1 \geq \dots \geq \tau_d$ .

$$\text{Show: } \left. \begin{array}{l} \tau_i \geq \max(\lambda_i, \mu_i) \quad \forall i \\ \Downarrow \\ 2\tau_i \geq \lambda_i + \mu_i + |\lambda_i - \mu_i| \quad \forall i \end{array} \right\} \Rightarrow \begin{aligned} \|\rho - \sigma\|_1 &= \text{Tr}[P + Q] \\ &= \text{Tr}[2X] - \text{Tr}[\rho] - \text{Tr}[\sigma] \\ &= \sum_i 2\tau_i - \lambda_i - \mu_i \geq \sum_i |\lambda_i - \mu_i|. \end{aligned}$$

Left to show:  $\tau_i \geq \max(\lambda_i, \mu_i) \quad \forall i$

Let's show  $\tau_i \geq \lambda_i \rightsquigarrow \tau_i \geq \mu_i$  follows in the same way.

Setup:  $X = \sum_{\alpha} \tau_{\alpha} |u_{\alpha}\rangle\langle u_{\alpha}| \quad \tau_1 \geq \tau_2 \geq \dots \geq \tau_d$

$$S = \sum_i \lambda_i |v_i\rangle\langle v_i| \quad \lambda_1 \geq \lambda_2 \geq \dots \geq \lambda_d$$

$$X = P + S \geq S. \quad (*)$$

Consider two subspaces of  $H$

$$S_{\ell} = \text{span}\{|v_1\rangle, \dots, |v_{\ell}\rangle\} \rightsquigarrow \dim(S_{\ell}) = \ell$$

$$M_{\ell} = \text{span}\{|u_{\ell}\rangle, \dots, |u_d\rangle\} \rightsquigarrow \dim(M_{\ell}) = d - \ell + 1.$$

$$\dim(M_{\ell}) + \dim(S_{\ell}) = d - \ell + 1 + \ell = d + 1 > d$$

$$\Rightarrow \exists |\gamma\rangle \in S_{\ell} \cap M_{\ell}, \quad \langle \gamma | \gamma \rangle = 1$$

Two expansions:  $|\gamma\rangle = \sum_{j=\ell}^d \alpha_j |u_j\rangle = \sum_{i=1}^{\ell} \beta_i |v_i\rangle \quad \sum_j |\alpha_j|^2 = \sum_i |\beta_i|^2 = 1.$

$$\Rightarrow \tau_{\ell} \geq \sum_{j=\ell}^d |\alpha_j|^2 \tau_j = \langle \gamma | X | \gamma \rangle \stackrel{(*)}{\geq} \langle \gamma | S | \gamma \rangle = \sum_{i=1}^{\ell} |\beta_i|^2 \lambda_i \geq \lambda_{\ell}.$$

□



Step 2: Show that  $|H(\rho) - H(\sigma)| \leq \sum_i \eta(|\lambda_i - \mu_i|)$

$$\text{Calculus} \Rightarrow |\eta(\lambda_i) - \eta(\mu_i)| \leq \eta(|\lambda_i - \mu_i|) \quad (*)$$

With (\*) we have

$$|H(\rho) - H(\sigma)| = \left| \sum_i \eta(\lambda_i) - \eta(\mu_i) \right| \leq \sum_i |\eta(\lambda_i) - \eta(\mu_i)| \leq \sum_i \eta(|\lambda_i - \mu_i|)$$

Step 3: Show that  $|H(\rho) - H(\sigma)| \leq \|\rho - \sigma\|_1 \log(d) + \eta(\|\rho - \sigma\|_1)$

$$\eta(x) = -x \log(x)$$

$$S = \sum_{i=1}^d |\lambda_i - \mu_i| \stackrel{\text{Step 1}}{\leq} \|\rho - \sigma\|_1$$

$$\eta(|\lambda_i - \mu_i|) = S \eta\left(\frac{|\lambda_i - \mu_i|}{S}\right) - |\lambda_i - \mu_i| \log(S)$$

$$\leadsto S \eta\left(\frac{|\lambda_i - \mu_i|}{S}\right) = -S \cdot \frac{|\lambda_i - \mu_i|}{S} (\log(|\lambda_i - \mu_i|) - \log(S)) = \eta(|\lambda_i - \mu_i|) + |\lambda_i - \mu_i| \log(S)$$

$$|H(\rho) - H(\sigma)| \stackrel{\text{Step 2}}{\leq} \sum_i \eta(|\lambda_i - \mu_i|) = S \underbrace{\sum_{i=1}^d \eta\left(\frac{|\lambda_i - \mu_i|}{S}\right)}_{= H(P) \leq \log(d)} - \underbrace{\sum_{i=1}^d |\lambda_i - \mu_i| \log(S)}_{= S \log(S) = \eta(S)}$$

$P_i = \frac{|\lambda_i - \mu_i|}{S}$

$$\leq S \log(d) + \eta(S)$$

$$\|\rho - \sigma\|_1 \leq \frac{1}{2}$$

$$\leq \|\rho - \sigma\|_1 \log(d) + \eta(\|\rho - \sigma\|_1)$$

$\eta$  monotone increasing

□