

Van Kampen's theorem (§ 1.2)

Motivation : suppose $X = A \cup B$, and choose

$x_0 \in A \cap B \rightsquigarrow$ how do we compute $\pi_1(X, x_0)$

from $\pi_1(A, x_0)$ and $\pi_1(B, x_0)$?

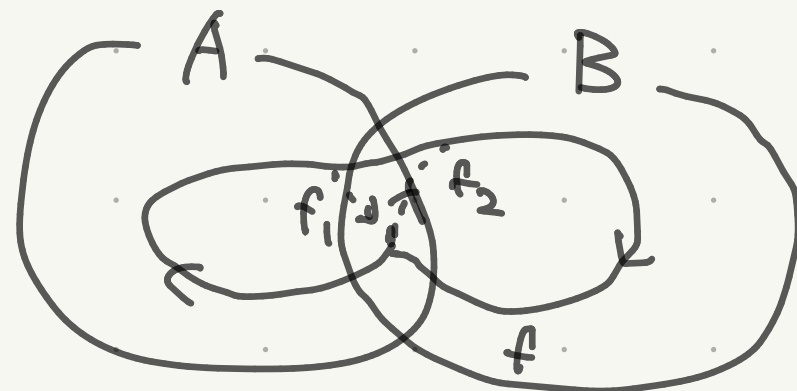
- if $A \cap B$ is path-connected, any $[f] \in \pi_1(X, x_0)$

can be written as $[f_1] \cdots [f_n]$, with

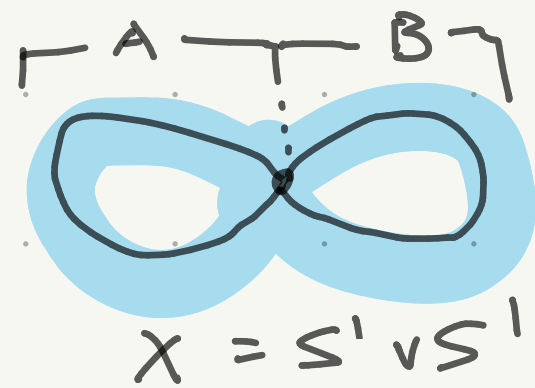
$f_i \in \pi_1(A, x_0)$ or $\pi_1(B, x_0)$

\rightsquigarrow we need to understand the relation between the images of

$\pi_1(A, x_0)$ and $\pi_1(B, x_0)$ in $\pi_1(X, x_0)$



"easy" case : $A \cap B \simeq \{x_0\} \rightsquigarrow$ there would be no (nontrivial) relation between them



\rightsquigarrow we should / can distinguish



$\rightsquigarrow \pi_1(X, x_0)$ is going to be "freely" generated

by a copy of $\pi_1(A, x_0)$ and a one of $\pi_1(B, x_0)$

more intricate : if $A \cap B$ has nontrivial π_1 ,

the images of $\pi_1(A \cap B, x_0)$ in $\pi_1(A, x_0)$ and

$\pi_1(B, x_0)$ should be identified in $\pi_1(X, x_0)$

$$\begin{array}{ccccc} A \cap B & \rightarrow & A & \rightarrow & X \\ & & \downarrow & & \uparrow \\ & & B & \rightarrow & X \end{array}$$

Free product of groups

goal: make sense of a group G with subgroups

G_1, G_2, \dots such that:

- they generate G (we have enough)
- no nontrivial relation among them (no redundancy)

"trivial rel": $(g_1 g_2) (g_2^{-1} g_1^{-1}) = g_1 g_1^{-1}$, etc. $g_1, g_1^{-1} \in G_1$
 $g_2 \in G_2$

observation: in particular we should have:

- any elem $g \in G$ can be written as

$$g = g_1 \cdots g_n, \quad g_i \in G_{j_i}$$

for some unique j_i

Def. $(G_\alpha)_{\alpha \in I}$ collection of groups;

- a word (for this collection) is an expression

$g_1 g_2 \dots g_n$ with $g_i \in G_{\alpha_i}$ for some $\alpha_i \in I$

(formally an elem. of $(\coprod_{\alpha \in I} G_\alpha)^{\{1, \dots, n\}}$ for some n)

- a reduced word is a word $g_1 \dots g_n$ ($g_i \in G_{\alpha_i}$)

such that $\alpha_i \neq \alpha_{i+1}$ for $i = 1, \dots, n-1$, $g_i \neq e$

- the free product of $(G_\alpha)_{\alpha \in I}$ is the set of

all reduced words (+ the "empty word" e of

length 0)

write

$\ast_{\alpha \in I} G_\alpha$,

$\ast (G_\alpha)_{\alpha \in I}$, etc.

Non-reduced words

$g g' h$

$g, g' \in G_\alpha, h \in G_\beta$

$g e h$

$g \in G_\alpha, e \in G_\beta$ (unit elem)

$h \in G_\gamma$

Reduced

$g h g'$

$g, g' \in G_\alpha, h \in G_\beta$

Group structure on $\ast_{\alpha \in I} G_{\alpha}$

want : - each G_{β} is a subgroup of $\ast_{\alpha} G_{\alpha}$

- interpret a reduced word $g_1 \dots g_n \in \ast_{\alpha} G_{\alpha}$

as the product of g_1, g_2, \dots, g_n

formal def of $(g_1 \dots g_m) \cdot (h_1 \dots h_n)$

$g_i \in G_{\alpha_i}$ $h_i \in G_{\beta_i}$

- if $\alpha_m \neq \beta_1$, take the reduced word $g_1 \dots g_m h_1 \dots h_n$ as the prod.

- if $\alpha_m = \beta_1$ and $g_m \neq h_1^{-1}$, take the reduced word $g_1 \dots g_{m-1} \underbrace{(g_m h_1)}_{\text{compute in } G_{\alpha_m}} h_2 \dots h_n$ as the prod.

(cont.) if $\alpha_m = \beta_1$ and $g_m = h_1^{-1}$, look at
 $(g_1 \dots g_{m-1}) (h_2 \dots h_n)$ and use the same
algorithm as above to compute the prod.

(induction on length; $g_i g_i^{-1} = e$)
the empty word e is the neutral element
empty word

$$(g_1 \dots g_n) e = g_1 \dots g_n = e (g_1 \dots g_n)$$

∴ the reduced word $g_n^{-1} \dots g_1^{-1}$ becomes the
inverse of $g_1 \dots g_n$

$$(g_1 \dots g_n) (g_n^{-1} \dots g_1^{-1}) = (g_1 \dots g_{n-1}) (g_{n-1}^{-1} \dots g_1^{-1}) = \dots = e$$

Example: free group F_S

S : set (will be the index set "I")

$G_s \cong \mathbb{Z}$ for each $s \in S$, $g_s \in G_s$ corresponding to $1 \in \mathbb{Z}$

$F_S = \ast_{s \in S} G_s$ free group on the set S

$$= \{ g_{s_1}^{n_1} g_{s_2}^{n_2} \cdots g_{s_k}^{n_k} : k=0,1,\dots, s_i \neq s_{i+1}, n_i \in \mathbb{Z} \setminus \{0\} \}$$

- generally: group G with generating set $S \subset G$

\rightsquigarrow surjective group hom $F_S \rightarrow G$, $g_s \mapsto s$

$$g_{s_1}^{n_1} g_{s_2}^{n_2} \cdots g_{s_k}^{n_k} \mapsto s_1^{n_1} \cdots s_k^{n_k}$$

$$G = (\Sigma \mid R)$$

↑ ↖ relation
generator set

R : some set of words in letters $s \in \Sigma$

e.g. $s_1 s_2, s_1 s_2 s_1, \dots$

$$F_\Sigma / \langle R \rangle$$

Basic property of free product:

$G_\alpha \xrightarrow{\varphi_\alpha} H$ group hom for each $\alpha \in I$

\equiv group hom $\ast_{\alpha \in I} G_\alpha \xrightarrow{\varphi} H$
equivalent

up to the correspondence $\varphi(g) = \varphi_\alpha(g)$ for $g \in G_\alpha$

$\ast_{\alpha} G_\alpha$ is the coproduct of $(G_\alpha)_{\alpha}$ in the category of groups; compare with direct product

$H \xrightarrow{\psi} \prod_{\alpha} G_\alpha \equiv$ group homs $H \xrightarrow{\varphi_\alpha} G_\alpha$

in particular $G \supset G_\alpha$, $(G_\alpha)_{\alpha \in I}$ generating G
subgrp.

\rightsquigarrow surjective group hom $\ast_{\alpha \in I} G_\alpha \rightarrow G$

$$\text{Hom} \left(\prod_{\alpha \in I} G_{\alpha}, H \right) = \left\{ \text{group hom } \prod_{\alpha \in I} G_{\alpha} \xrightarrow{\varphi} H \right\}$$

$$\updownarrow 1:1$$

$$\prod_{\alpha \in I} \text{Hom}(G_{\alpha}, H) = \left\{ (\varphi_{\alpha})_{\alpha \in I} : \varphi_{\alpha} : G_{\alpha} \rightarrow H \right\}$$

$$\varphi \rightsquigarrow (\varphi_{\alpha})_{\alpha}$$

$$\varphi_{\alpha} = \varphi|_{G_{\alpha}}$$

$$(\varphi_{\alpha})_{\alpha \in I} \rightsquigarrow \varphi(g_1, \dots, g_k) = \varphi_{\alpha_1}(g_1) \dots \varphi_{\alpha_k}(g_k)$$

$g_i \in G_{\alpha_i}$

$$\psi : H \rightarrow K \quad \text{hom} \quad \psi \circ \varphi \in \text{Hom} \left(\prod_{\alpha \in I} G_{\alpha}, K \right)$$

$$\updownarrow$$

$$(\psi \circ \varphi)_{\alpha \in I} \in \prod_{\alpha \in I} \text{Hom}(G_{\alpha}, K)$$

Examples

(infinite dihedral group)

$$1. \quad \mathbb{Z}_2 * \mathbb{Z}_2 \cong \mathbb{Z} \rtimes_{\pm 1} \mathbb{Z}_2 \quad \mathbb{Z}_2 = \mathbb{Z}/2\mathbb{Z}$$

$\left. \begin{array}{l} a^e \\ \downarrow \\ \mathbb{Z}_2 \end{array} \right\} \downarrow \left. \begin{array}{l} b^e \\ \downarrow \\ \mathbb{Z}_2 \end{array} \right\}$

 semidirect prod.

$$\langle a, b \mid a^2 = e = b^2 \rangle \quad \rightsquigarrow \{ (n, g) : n \in \mathbb{Z}, g \in \mathbb{Z}_2 \}$$

$$(m, g) \cdot (n, h) = (m + n^g, gh)$$

$$n^g = -n \quad \text{for the nontriv. } g$$

with nontriv. $g \in \mathbb{Z}_2$:

$$\begin{cases} (1, g)^2 = (1 - 1, g^2) = (0, e) = e \\ (0, g)^2 = (0, g^2) = e \end{cases}$$

$$\rightsquigarrow \mathbb{Z}_2 * \mathbb{Z}_2 \rightarrow \mathbb{Z} \rtimes_{\pm 1} \mathbb{Z}_2, \quad a \mapsto (1, g), \quad b \mapsto (0, g)$$

this is an isomorphism.

$$(b a b^{-1} = b a = (a b)^{-1})$$

$$\swarrow (0, g) (1, e) (0, g)^{-1} = (-1, e)$$

$$2 \quad \text{PSL}_2(\mathbb{Z}) = \text{SL}_2(\mathbb{Z}) / \{\pm I_2\} \quad (\text{modular group})$$

is isomorphic to $\mathbb{Z}_2 * \mathbb{Z}_3$

the image of $\begin{bmatrix} -1 & -1 \\ 0 & 1 \end{bmatrix}$ in $\text{PSL}_2(\mathbb{Z})$ has order 3

$\begin{bmatrix} 0 & -1 \\ 1 & 0 \end{bmatrix}$ has order 2

$\leadsto \mathbb{Z}_2 * \mathbb{Z}_3 \rightarrow \text{PSL}_2(\mathbb{Z})$, which turns out to be an isomorphism

Amalgamated free product

$(G_\alpha)_{\alpha \in I}$: collection of groups

A : group, $\varphi_\alpha: A \rightarrow G_\alpha$ homomorphism

(φ_α injective \leadsto "common subgroup of the G_α ")

Def. the amalgamated free product of the groups $(G_\alpha)_\alpha$ over A is

$$\ast_A (G_\alpha)_{\alpha \in I} = \left(\ast_{\alpha \in I} G_\alpha \right) / N,$$

N = the normal subgroup generated by

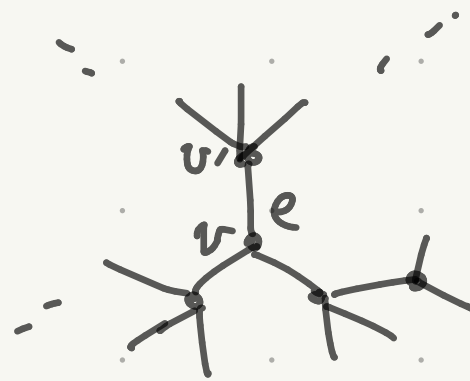
$$\varphi_\alpha(x) \varphi_\beta(x)^{-1} \quad x \in A, \alpha, \beta \in I$$

$$\leadsto \text{ in } \ast_A (G_\alpha)_\alpha \quad [\varphi_\alpha(x)] = [\varphi_\beta(x)]$$

So $\ast_A (G_\alpha)_{\alpha \in I}$: generated by $(G_\alpha)_{\alpha \in I}$ - with identification of $\varphi_\alpha(x)$ and $\varphi_\beta(x)$ for $x \in A$ $\alpha, \beta \in I$ but no other nontrivial relation

" Bass - Serre theory "

automorphism groups of trees



are built up using

a malgumated free product & another

operation (HNN extension)

(stabilizer of an edge e) \subset (stab. of a vertex v)

if v is an endpoint of e

Van Kampen's th'm ("usual" version)

Suppose $X = A \cup B$, A & B open,

$A \cap B$ path-connected, $x_0 \in A \cap B$

then $\pi_1(X, x_0) \cong \pi_1(A, x_0) *_{\pi_1(A \cap B, x_0)} \pi_1(B, x_0)$

with respect to homs

$\rightarrow \pi_1(A, x_0)$

$\pi_1(A \cap B, x_0) \rightarrow \pi_1(B, x_0)$

induced by inclusion $A \cap B \rightarrow A$ $\cong \mathbb{Z}$
 $\rightarrow B$

Consequence $\pi_1(S^1 \vee S^1, x_0) \cong F_2 (= \mathbb{Z} * \mathbb{Z})$



$\begin{array}{c} \text{---} \text{---} \text{---} \text{---} \text{---} \\ | \\ A \\ | \\ \text{---} \text{---} \text{---} \text{---} \text{---} \\ | \\ B \end{array}$

$F_{\{a, b\}}$
 and $\pi_1(\underbrace{S^1 \vee \dots \vee S^1}_{n \times}, x_0) \cong F_n$

$\pi_1(\bigvee_{\alpha \in I} X_{\alpha}, x_0) \cong *_{\alpha \in I} (\pi_1(X_{\alpha}, x_0))_{\alpha \in I}$

Van Kampen's th'm (general version)

$$X = \bigcup_{\alpha \in I} A_{\alpha}, \quad A_{\alpha} \text{ open (path-conn)}$$

$$A_{\alpha} \cap A_{\beta} \cap A_{\gamma} \text{ path-conn. for all } \alpha, \beta, \gamma \in I$$

$$x_0 \in \bigcap_{\alpha \in I} A_{\alpha}$$

$$\text{then } \pi_1(X, x_0) \cong \left(\ast_{\alpha \in I} \pi_1(X_{\alpha}, x_0) \right) / N$$

N : the normal subgroup of $\ast_{\alpha \in I} \pi_1(X_{\alpha}, x_0)$

generated by $i_{\alpha\beta}(\omega) i_{\beta\alpha}(\omega)^{-1}$ for $\omega \in \pi_1(A_{\alpha} \cap A_{\beta}, x_0)$

$$i_{\alpha\beta} : \pi_1(A_{\alpha} \cap A_{\beta}, x_0) \rightarrow \pi_1(A_{\alpha}, x_0)$$

$$i_{\beta\alpha} : \pi_1(A_{\beta} \cap A_{\alpha}, x_0) \rightarrow \pi_1(A_{\beta}, x_0)$$

$$\text{i.e. } [i_{\alpha\beta}(\omega)] = [i_{\beta\alpha}(\omega)] \quad (\omega \in \pi_1(A_{\alpha} \cap A_{\beta})) \text{ in } \pi_1(X)$$

