STK2130 - Week 4

A. B. Huseby

Department of Mathematics University of Oslo, Norway

First-passage probabilities

Consider a Markov chain $\{X_n\}$ with state space S, and let A be a non-empty proper subset of S.

We want to calculate the following probability:

$$P(X_k \in \mathcal{A} \text{ for some } 1 \leq k \leq m | X_0 = i)$$

In order to analyze this we introduce:

$$N = \min\{n : X_n \in A\},\$$

where we let $N = \infty$ is $X_n \notin A$ for all n.

NOTE: *N* represents the first time the Markov chain enters A.

$$W_n = \left\{ \begin{array}{ll} X_n & \text{if } n < N \\ A & \text{if } n \ge N \end{array} \right.$$

NOTE: When $\{X_n\}$ enters A, $\{W_n\}$ is absorbed in state A.

The transition probabilities of $\{W_n\}$, denoted $Q_{i,j}$, are given by:

$$Q_{i,j} = P_{i,j}, \qquad i \notin \mathcal{A}, j \notin \mathcal{A},$$
 $Q_{i,A} = \sum_{j \in \mathcal{A}} P_{i,j}, \quad i \notin \mathcal{A}, j \in \mathcal{A},$ $Q_{A,A} = 1.$

We now have:

$$P(X_k \in \mathcal{A} \text{ for some } 1 \le k \le m | X_0 = i)$$

$$= P(W_m = A | X_0 = i)$$

$$= P(W_m = A | W_0 = i) = Q_{i,A}^m.$$

We now consider the following probability:

$$\alpha = P(X_k \notin \mathcal{A} \text{ for all } 1 \leq k \leq m-1, X_m = j | X_0 = i)$$

CASE 1. $i, j \notin A$

In this case the event $\{X_k \notin A \text{ for all } 1 \leq k \leq m-1, X_m = j\}$ is equivalent to the event $\{W_m = j\}$.

Hence, it follows that:

$$\alpha = P(W_m = j | W_0 = i) = Q_{i,j}^m.$$

CASE 2. $i \notin A$ and $j \in A$

In this case we have:

$$\alpha = \sum_{r \notin \mathcal{A}} P(X_k \notin \mathcal{A} \text{ for all } 1 \le k \le m-2, X_{m-1} = r, X_m = j | X_0 = i)$$

$$=\sum_{r\notin\mathcal{A}}P(X_k\notin\mathcal{A}\text{ for all }1\leq k\leq m-2,X_{m-1}=r|X_0=i)$$

$$P(X_m = j | X_0 = i, X_k \notin A \text{ for all } 1 \le k \le m-2, X_{m-1} = r)$$

$$= \sum_{r \notin \mathcal{A}} P(X_k \notin \mathcal{A} \text{ for all } 1 \le k \le m-2, X_{m-1} = r | X_0 = i) \cdot P(X_m = j | X_{m-1} = r)$$

$$= \sum_{r \notin \mathcal{A}} Q_{i,r}^{m-1} \cdot P_{r,j}$$



CASE 3. $i \in A$ and $j \notin A$

In this case we have:

$$\alpha = \sum_{r \notin \mathcal{A}} P(X_1 = r, X_k \notin \mathcal{A} \text{ for all } 2 \le k \le m - 1, X_m = j | X_0 = i)$$

$$= \sum_{r \notin \mathcal{A}} P(X_1 = r | X_0 = i)$$

$$\cdot P(X_k \notin \mathcal{A} \text{ for all } 2 \le k \le m - 1, X_m = j | X_0 = i, X_1 = r)$$

$$= \sum_{r \notin \mathcal{A}} P(X_1 = r | X_0 = i)$$

$$\cdot P(X_k \notin \mathcal{A} \text{ for all } 2 \le k \le m - 1, X_m = j | X_1 = r)$$

$$= \sum_{r \notin \mathcal{A}} P_{i,r} \cdot Q_{r,j}^{m-1}$$

CASE 4. $i \in A$ and $j \in A$

By combining the previous arguments we get:

$$\alpha = \sum_{r,s \notin \mathcal{A}} P_{i,r} \cdot Q_{r,s}^{m-2} \cdot P_{s,j}$$

Unconditional probabilities

We introduce the probabilities:

$$\pi_i^{(n)} = P\{X_n = i\}, \quad i \in \mathcal{S}, \quad n = 0, 1, 2, \dots$$

We then have:

$$\pi_j^{(n+m)} = \sum_{i \in \mathcal{S}} P\{X_{n+m} = j \cap X_n = i\}$$

$$= \sum_{i \in \mathcal{S}} P\{X_n = i\} \cdot P\{X_{n+m} = j | X_n = i\}$$

$$= \sum_{i \in \mathcal{S}} \pi_i^{(n)} \cdot P_{ij}^m.$$

In particular:

$$\pi_j^{(1)} = \sum_{i \in \mathcal{S}} \pi_i^{(0)} \cdot P_{ij}, \qquad \pi_j^{(n+1)} = \sum_{i \in \mathcal{S}} \pi_i^{(n)} \cdot P_{ij}$$

Unconditional probabilities (cont.)

Assume that:

$$\lim_{n\to\infty}\pi_i^{(n)}=\pi_i,\quad i\in\mathcal{S}.$$

Then we obviously also have:

$$\lim_{n\to\infty}\pi_j^{(n+1)}=\pi_j,\quad j\in\mathcal{S}.$$

In particular, if $S = \{1, ..., k\}$ and $\pi = (\pi_1, ..., \pi_k)$, then:

$$\pi_{j} = \lim_{n \to \infty} \pi_{j}^{(n+1)} = \lim_{n \to \infty} \sum_{i=1}^{k} \pi_{i}^{(n)} \cdot P_{ij} = \sum_{i=1}^{k} \lim_{n \to \infty} \pi_{i}^{(n)} \cdot P_{ij} = \sum_{i=1}^{k} \pi_{i} \cdot P_{ij}$$

Thus, π must satisfy:

$$\pi = \pi \cdot P$$
.



Example 4.8 revisited

 $P\{\text{Rain tomorrow}|\text{Rain today}\} = 0.75$ $P\{\text{Rain tomorrow}|\text{No rain today}\} = 0.35$

$$\mathbf{P} = \left[\begin{array}{cc} 0.75 & 0.25 \\ 0.35 & 0.65 \end{array} \right]$$

$$\mathbf{P}^{(2)} = \left[\begin{array}{cc} 0.65 & 0.35 \\ 0.49 & 0.51 \end{array} \right]$$

$$\mathbf{P}^{(4)} = \left[\begin{array}{cc} 0.5940 & 0.4060 \\ 0.5684 & 0.4316 \end{array} \right]$$

$$\mathbf{P}^{(12)} = \left[\begin{array}{cc} 0.5833 & 0.4167 \\ 0.5833 & 0.4167 \end{array} \right]$$

Example 4.8 revisited (cont.)

We now let $\pi = (\pi_1, \pi_2)$, and consider the equation:

$$\pi = \pi \cdot P$$
.

which in this case becomes:

$$\pi_1 = 0.75\pi_1 + 0.35\pi_2$$

$$\pi_2 = 0.25\pi_1 + 0.65\pi_2$$

By inserting $\pi_2 = 1 - \pi_1$ into the first equation, we get:

$$\pi_1 = 0.75\pi_1 + 0.35(1 - \pi_1) = 0.40\pi_1 + 0.35$$

From this it follows that:

$$\pi_1 = 0.35/(1 - 0.40) = 0.5833$$

 $\pi_2 = 1 - \pi_1 = 0.4167$

Chapter 4.3. Classification of States

Let $\{X_n\}$ be a Markov chain with state space S and transition probability matrix P.

State j is said to be accessible from state i, denoted as $i \to j$, if $P_{ij}^n > 0$ for some n > 0.

Note that we have:

$$\max_{n} P_{ij}^{n} \leq P(\bigcup_{n=1}^{\infty} \{X_{n} = j\} | X_{0} = i)$$

$$\leq \sum_{n=0}^{\infty} P\{X_{n} = j | X_{0} = i\} = \sum_{n=0}^{\infty} P_{ij}^{n}.$$

Hence, $i \rightarrow j$ if and only if:

$$P(\bigcup_{n=1}^{\infty} \{X_n = j\} | X_0 = i) > 0.$$

Chapter 4.3. Classification of States (cont.)

A state diagram for a Markov chain is a directed graph where the nodes represent the states and the edges represent possible one-step transitions. More precisely, the state diagram contains an edge from node i to node j if and only if $P_{ij} > 0$.

If $i, j \in S$, then $i \to j$ if and only if the state diagram contains at least one directed path from i to j.



If such a path exists, we have:

$$P_{ij}^n \geq P_{i,k_1} \cdot P_{k_1,k_2} \cdots P_{k_{n-2},k_{n-1}} \cdot P_{k_{n-1},j} > 0.$$

Communicating states

States *i* and *j* communicate, denoted as $i \leftrightarrow j$, if $i \rightarrow j$ and $j \rightarrow i$.

The relation \leftrightarrow is an equivalence relation. That is \leftrightarrow satisfies the following properties:

- Reflexivity: $i \leftrightarrow i$.
- Symmetry: $i \leftrightarrow j$ if and only if $j \leftrightarrow i$.
- Transitivity: $i \leftrightarrow j$ and $j \leftrightarrow k$ implies $i \leftrightarrow k$.

Reflexivity follows since we always have $P_{ii}^0 = 1 > 0$. Symmetry follows directly from the definition.

Communicating states (cont.)

To prove transitivity we assume that $i \leftrightarrow j$ and $j \leftrightarrow k$.

Hence, in particular $i \to j$ and $j \to k$, implying that there exists $m, n \ge 0$ such that $P_{ij}^m > 0$ and $P_{jk}^n > 0$.

By the Chapman-Kolmogorov equations, we have:

$$P_{ik}^{m+n} = \sum_{r \in \mathcal{S}} P_{ir}^m P_{rk}^n \ge P_{ij}^m \cdot P_{jk}^n > 0.$$

Hence, by definition $i \to k$.

By a similar argument we can show that $k \to i$ as well.

Hence, we conclude that $i \leftrightarrow k$.

Communicating states (cont.)

Two states that communicate are said to be in the same (equivalence) class.

Two classes of states are either identical or disjoint.

PROOF: Assume that $A, B \subseteq S$ represent two equivalence classes, and assume that $A \cap B \neq \emptyset$. That is, there exists a state i such that $i \in A \cap B$.

Then choose $j \in A$ and $k \in B$ arbitrarily.

Now, $i, j \in \mathcal{A}$ implies that $i \leftrightarrow j$ and $i, k \in \mathcal{B}$ implies that $i \leftrightarrow k$.

Hence, by transitivity we also have $j \leftrightarrow k$. That is, j and k belong to the same equivalence class.

Since this holds for any $j \in A$ and $k \in B$, this implies that A = B

The equivalence classes partition the state space \mathcal{S} into a number of disjoint sets. A Markov chain is called <u>irreducible</u> if the number of equivalence classes is <u>one</u>.

Example 4.15

Consider a Markov chain with state space $S = \{0, 1, 2\}$ and transition probability matrix:

$$\mathbf{P} = \left[\begin{array}{ccc} \frac{1}{2} & \frac{1}{2} & 0 \\ \frac{1}{2} & \frac{1}{4} & \frac{1}{4} \\ 0 & \frac{1}{3} & \frac{2}{3} \end{array} \right]$$

We then observe:

Since
$$P_{01} = \frac{1}{2} > 0$$
, it follows that $0 \to 1$

Since
$$P_{10} = \frac{1}{2} > 0$$
, it follows that $1 \to 0$

Since
$$P_{12} = \frac{1}{4} > 0$$
, it follows that $1 \rightarrow 2$

Since
$$P_{21} = \frac{1}{3} > 0$$
, it follows that $2 \rightarrow 1$

Hence, $0 \leftrightarrow 1$ and $1 \leftrightarrow 2$, and by transitivity $0 \leftrightarrow 2$ as well. Thus, the Markov chain is irreducible.

Example 4.15 (cont.)

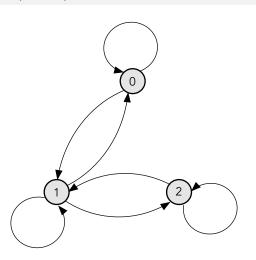


Figure: State diagram of an irreducible Markov chain with one class $\{0,1,2\}$

Example 4.16

A Markov chain with state space $S = \{0, 1, 2, 3\}$ and matrix:

$$\mathbf{P} = \begin{bmatrix} \frac{1}{2} & \frac{1}{2} & 0 & 0 \\ \frac{1}{2} & \frac{1}{2} & 0 & 0 \\ \frac{1}{4} & \frac{1}{4} & \frac{1}{4} & \frac{1}{4} \\ 0 & 0 & 0 & 1 \end{bmatrix}$$

$$P_{01}=P_{10}=\frac{1}{2}, \text{ implying that } 0\leftrightarrow 1$$

 $P_{ij}=0, \text{ implying that } i\not\rightarrow j, \quad i=0,1,\quad j=2,3$
 $P_{2i}=\frac{1}{4}, \text{ implying that } 2\rightarrow i, \quad i=0,1,2,3$
 $P_{3i}=0, \text{ implying that } 3\not\rightarrow i, \quad i=0,1,2$

The Markov chain has classes {0,1}, {2} and {3}, and is not irreducible.

Example 4.16 (cont.)

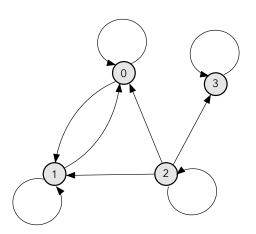


Figure: State diagram of a Markov chain with three classes $\{0,1\}$, $\{2\}$ and $\{3\}$.

Recurrent and transient states

We consider the probabilities:

$$f_i = P(\bigcup_{r=1}^{\infty} \{X_r = i\} | X_0 = i\}, \quad i \in \mathcal{S}.$$

- State *i* is recurrent if $f_i = 1$.
- State *i* is transient if $f_i < 1$.

Assume that $X_0 = i$, and let N_i denote the number of times $X_n = i$.

- If *i* is recurrent, then $P(N_i = \infty) = 1$.
- If *i* is transient, then $P(N_i = n) = f_i^{n-1}(1 f_i), n = 1, 2, ...$

If *i* is transient and $X_0 = i$, then N_i has a geometric distribution with $E[N_i] = 1/(1 - f_i)$.

Proposition 4.1

Let $I_i^{(n)} = I(X_n = i), n = 0, 1, ...$ We can then write:

$$N_i = \sum_{n=0}^{\infty} I_i^{(n)}$$

Hence, we have:

$$E[N_i|X_0 = i] = \sum_{n=0}^{\infty} E[I_i^{(n)}|X_0 = i]$$

$$= \sum_{n=0}^{\infty} P[X_n = i|X_0 = i] = \sum_{n=0}^{\infty} P_{ii}^n$$

- State *i* is recurrent, if $\sum_{n=1}^{\infty} P_{ii}^n = \infty$.
- State *i* is transient, if $\sum_{n=1}^{\infty} P_{ii}^n < \infty$.

Corollary 4.2

If state i is recurrent, and $i \leftrightarrow j$, then state j is recurrent as well. Thus, recurrence is a class property.

PROOF: Since $i \leftrightarrow j$, there exists k and m such that $P_{ij}^k > 0$ and $P_{ji}^m > 0$.

Hence, for any n = 1, 2, ... we have:

$$P_{jj}^{m+n+k} \geq P_{ji}^m \cdot P_{ii}^n \cdot P_{ij}^k.$$

Summing over all n, and using that i is recurrent, $P_{ij}^k > 0$ and $P_{ji}^m > 0$ we get:

$$\sum_{n=1}^{\infty} P_{jj}^{m+n+k} \ge P_{ji}^{m} \cdot P_{ij}^{k} \cdot \sum_{n=1}^{\infty} P_{ii}^{n} = \infty$$

Hence, we conclude that *j* is recurrent as well

Corollary 4.2 (cont.)

- Corollary 4.2 also implies that transience is a class property. For if state i is transient and $i \leftrightarrow j$, then state j must also be transient. For if j were recurrent then, by Corollary 4.2, i would also be recurrent and hence could not be transient.
- Corollary 4.2 along with the fact that not all states in a finite Markov chain can be transient leads to the conclusion that all states of a finite irreducible Markov chain are recurrent.

Example 4.17

Consider a Markov chain with state space $\mathcal{S} = \{0,1,2,3\}$ and transition probability matrix:

$$\mathbf{P} = \left[\begin{array}{cccc} 0 & 0 & \frac{1}{2} & \frac{1}{2} \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{array} \right]$$

It is easy to verify that $i \leftrightarrow j$ for all $i, j \in S$. Hence, the Markov chain is irreducible and thus all states must be recurrent

Example 4.17 (cont.)

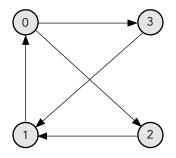


Figure: State diagram of an irreducible Markov chain with one class $\{0,1,2,3\}$

Example 4.18

Consider a Markov chain with state space $\mathcal{S} = \{0, 1, 2, 3, 4\}$ and transition probability matrix:

$$m{P} = \left[egin{array}{ccccc} rac{1}{2} & rac{1}{2} & 0 & 0 & 0 \\ rac{1}{2} & rac{1}{2} & 0 & 0 & 0 \\ 0 & 0 & rac{1}{2} & rac{1}{2} & 0 \\ 0 & 0 & rac{1}{2} & rac{1}{2} & 0 \\ rac{1}{4} & rac{1}{4} & 0 & 0 & rac{1}{2} \end{array}
ight]$$

This chain has classes $\{0,1\}$, $\{2,3\}$ and $\{4\}$.

The first two classes are recurrent and the third transient

Example 4.18 (cont.)

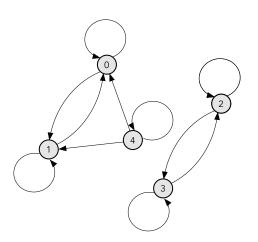


Figure: State diagram of a Markov chain with classes $\{0,1\}$, $\{2,3\}$ and $\{4\}$

Example 4.19 - Random walk

Consider a Markov chain with state space $\mathcal{S} = \{\dots, -2, -1, 0, 1, 2, \dots\}$ and where 0 and:

$$P_{i,i+1} = p, \quad P_{i,i-1} = (1-p), \quad i \in S.$$

It is obvious that $i \leftrightarrow j$ for all $i, j \in \mathcal{S}$. Hence, according to Corollary 4.2 all states are either recurrent or transient.

In order to check for recurrence, it is sufficient to check if $\sum_{n=1}^{\infty} P_{00}^n = \infty$.

We then observe that X_n is odd if n is odd, and X_n is even if n is even. Hence, since 0 is even, we have:

$$P_{00}^{2n-1}=0, \quad n=1,2,\ldots$$

$$P_{00}^{2n} = {2n \choose n} p^n (1-p)^n = \frac{(2n)!}{n! \, n!} [p(1-p)]^n, \quad n = 1, 2, \dots$$



Example 4.19 - Random walk (cont.)

We then use Stirling's formula for n!:

$$n! \approx n^{n+1/2}e^{-n}\sqrt{2\pi}$$

From this we get:

$$\frac{(2n)!}{n!n!} \approx \frac{(2n)^{2n+1/2}e^{-2n}\sqrt{2\pi}}{(n^{n+1/2}e^{-n}\sqrt{2\pi})^2} = \frac{(2n)^{2n+1/2}e^{-2n}\sqrt{2\pi}}{n^{2n+1}e^{-2n}(2\pi)} = \frac{2^{2n}}{\sqrt{n\pi}} = \frac{4^n}{\sqrt{n\pi}}$$

Hence:

$$P_{00}^{2n} = \frac{(2n)!}{n! \, n!} [p(1-p)]^n \approx \frac{(4p(1-p))^n}{\sqrt{n\pi}}$$

Example 4.19 - Random walk (cont.)

This implies that:

$$\sum_{n=1}^{\infty} P_{00}^{2n} \approx \sum_{n=1}^{\infty} \frac{(4p(1-p))^n}{\sqrt{n\pi}}$$

This series is divergent if and only if $p = \frac{1}{2}$.

Hence, the states are recurrent if and only if $p = \frac{1}{2}$.