

4.69.) M balls dist. among m urns.

• Stage i : Select one ball, place at random in one of $m-1$ other urns.

• $n_i := \#$ balls in urn i

• Markov chain with state: (n_1, \dots, n_m)

• Guess limiting prob. & verify. Show that the Markov chain is time-reversible.

Guess for limiting probabilities:

State space S is such that

$$i \in S \iff i = (n_1^i, n_2^i, \dots, n_m^i)$$

• where $\sum_{k=1}^m n_k^i = M$

Guess that

(**) ← (Ways of choosing n_1^i & ... & n_m^i out of M ; unordered)

$$\Pi_i = \frac{M!}{(n_1^i)! \cdots (n_m^i)!} \left[\frac{1}{m}\right]^{n_1^i} \cdots \left[\frac{1}{m}\right]^{n_m^i}$$

proportion of time spent in state i ; with n_1^i balls in urn 1, ..., n_m^i balls in urn m .

Probability of having n_1^i balls in 1, ..., n_m^i balls in m

↑
prob. of choosing urn 1 for n_1^i balls

↑
prob. of choosing urn m for n_m^i balls

$$= \frac{M!}{(n_1^i)! \cdots (n_m^i)!} \left[\frac{1}{m}\right]^M$$

Since $\sum_{j=1}^m n_j^i = M$

M ; (**) → Why is this true?

$$\Pi_i = \left(\begin{array}{l} \text{\# ways to choose} \\ n_1^i \text{ balls in 1,} \\ \dots \\ n_m^i \text{ balls in } m \end{array} \right) \cdot \left(\begin{array}{l} \text{prob. of} \\ \text{choosing} \\ \text{urn 1} \\ \text{for the} \\ n_1^i \text{ balls} \end{array} \right) \cdots \left(\begin{array}{l} \text{prob. of} \\ \text{choosing} \\ \text{urn } m \\ \text{for } n_m^i \\ \text{balls} \end{array} \right)$$

↓
= # of ways to group M into m groups of sizes n_1^i, \dots, n_m^i

$$= \frac{M!}{(n_1^i)! \dots (n_m^i)!}$$

(~)

because e.g.,

ways to group $M=10$ into $m=3$ groups of sizes $n_1=2, n_2=5, n_3=3$ is

$$\frac{10!}{2! \cdot 8!} \cdot \frac{8!}{5! \cdot 3!} \cdot 1 = \frac{10!}{2! \cdot 5! \cdot 3!}$$

ways to choose $n_1=2$ out of 10; binomial coefficient / combination formula

ways to choose $n_2=5$ out of the remaining $10-2=8$ balls

no need to choose the final group since we are to choose 3 & 3 are left

This generalizes (via induction) to (~).

We are told to check time-reversibility of the Markov chain & at the same time verify that

our guess of limiting probabilities is actually correct:

To check time reversibility, we need to check:

$$\pi_i P_{ij} \stackrel{?}{=} \pi_j P_{ji} \quad \forall i, j$$

Note that $P_{ij} \neq 0 \Leftrightarrow \exists r, s \in \{1, \dots, n\}$ s.t.

$$i = (n_1, \dots, n_r, \dots, n_s, \dots, n_m) \text{ \& } j = (n_1, \dots, n_{r-1}, \dots, n_{s+1}, \dots, n_m)$$

$$j = (n_1, \dots, n_{r-1}, \dots, n_{s+1}, \dots, n_m)$$

That is, $P_{ij} \neq 0$ iff. you take a ball from one urn (r here) and put it in another urn (s here)

$$\begin{array}{l}
 \underbrace{P_{ij}} = \underbrace{\frac{n_r}{M}} \cdot \underbrace{\frac{1}{m-1}} \\
 \text{prob. go from } i \text{ to } j = \text{prob. choose one of the } n_r \text{ balls in } r \text{ out of } M \text{ in total} \cdot \text{prob. choose } s \text{ out of remaining } m-1 \text{ urns} \\
 \text{\& } \underbrace{P_{ji}} = \underbrace{\frac{n_s+1}{M}} \cdot \underbrace{\frac{1}{m-1}} \\
 \text{prob. go from } j \text{ to } i = \text{prob. choose one of the } n_s+1 \text{ balls in } s \text{ out of } M \text{ in total} \cdot \text{choose } r \text{ out of } m-1 \text{ urns}
 \end{array}$$

Therefore,

$$\pi_i P_{ij} = \pi_j P_{ji} \quad \Leftrightarrow$$

$$\frac{n_r}{M} \frac{1}{m-1} \frac{M!}{(n_1)! \dots (n_m)!} = \frac{n_s+1}{M} \frac{1}{m-1} \frac{M!}{(n_1)! \dots (n_r-1)! \dots (n_s+1)! \dots (n_m)!}$$

\Leftrightarrow

$$\frac{n_r}{(n_s)! (n_r)!} = \frac{n_s+1}{(n_s+1)! (n_r-1)!}$$

\Leftrightarrow

$$\frac{1}{(n_s)! (n_r-1)!} = \frac{1}{(n_s)! (n_r-1)!}$$

On R.H.S. we have $(n_s+1)!$ & $(n_r-1)!$ instead of $(n_s)!$ & $(n_r)!$

Def of faculty

Which is clearly true.

Hence, by tracing the iff.

backwards, we find that

$$\prod_i P_{ij} = \prod_j P_{ji} \quad \forall i, j, \quad (*)$$

and hence we have time-reversibility.

From top of pg. 252: If we can find non-negative numbers summing to 1 s.t. (*) holds, these are the limiting probabilities.

We know (*) holds, also

$$\pi_i = \frac{M!}{(n_1^i)! \dots (n_m^i)!} \underbrace{\left(\frac{1}{m}\right)^M}_{\geq 0} \geq 0$$

Finally, if the numbers do not sum to one, we can make this OK by normalizing

(\star) $\alpha \geq 0$ still holds).

4.73.) • k players, $v_i > 0, i = 1, \dots, k$ is value of player i .

• Each period: 2 play a game.

• Winner plays against randomly chosen player among $k-1$ remaining.

• If i & j play, i wins with prob. $\frac{v_i}{v_i + v_j}$

• $X_n =$ winner game n .

∞

a) Transition prob. of Markov chain $\{X_n\}_{n \geq 1}$

$$\begin{aligned}
 \underbrace{P_{ij}}_{\substack{\text{prob. } i \text{ won} \\ \text{game, then} \\ \text{plays against} \\ j \text{ \& } j \text{ wins}}} &= P(\underbrace{j \text{ is chosen}}_{\text{player}}, \underbrace{j \text{ wins}}_{\text{against } i} \mid i \text{ won}) \\
 &= P(\text{choose } j \text{ out of } k-1) \cdot P(j \text{ wins over } i) \\
 &= \frac{1}{k-1} \cdot \frac{v_j}{v_i + v_j} \\
 &= \text{prob. go from state } i \text{ to } j
 \end{aligned}$$

Note: This is for i different from j . For $i=j$, one must sum over all $k-1$ potential other opponents in game nr. 2.

(6)

b) Stationarity equations:

$$\left\{ \begin{array}{l} \sum_{i=1}^k \pi_i = 1 \end{array} \right.$$

$$\vec{\pi} = \vec{\pi} \underline{P}$$

transition probability
matrix

Note: $\{X_n\}_{n \geq 1}$ is a finite state, irreducible

Markov chain (since $v_i > 0 \forall i$, this

means that $\frac{v_i}{v_j + v_i} > 0 \forall i, j$, so any player

can beat any other; Hence, all states communicate)

Hence, the stationarity eq. are given as above.

c) Time reversibility equations:

$$\pi_i P_{ij} = \pi_j P_{ji} \quad \forall i, j$$

If also $\pi_i \geq 0 \forall i$ & sum to 1, then from pg.

252 we know that these are actually the limiting probabilities.

d) Proportion of games won by j :

This is the same as finding π_j .

Can either solve stationarity eqns. from b) or time-reversibility eqns. from c).

Try to use stationarity eqns: Hopeless!

Instead: Aim to use c) & guess for solution.

$$\pi_i \frac{1}{k-1} \frac{v_j}{v_i + v_j} = \pi_j \frac{1}{k-1} \frac{v_i}{v_i + v_j} \quad \forall i \neq j$$

Note that the time reversibility is trivially true for $i=j$.

$$(*) \quad \pi_i v_j = \pi_j v_i \quad \forall i \neq j$$

$$\text{Guess: } \pi_i = \frac{v_i}{\sum_{j=1}^k v_j} > 0 \quad (\text{since } v_i > 0 \quad \forall i)$$

Also, (*) holds $\forall i, j$.

$$\text{Finally, } \sum_{i=1}^k \pi_i = \frac{1}{\sum_{i=1}^k v_i} \sum_{i=1}^k v_i = 1$$

Hence, from pg 252, this $\vec{\pi}$ is the limiting distribution.

$$\text{so } \pi_j = \frac{v_j}{\sum_{i=1}^k v_i}$$

e) Proportion of games involving j as player

$\{\vec{X}_j\}_{j \geq 1}$ is the Markov chain where $\vec{X}_j = (X_j, Y_j)$,
 & Y_j is the loser of game j .

Then,

$$P_{(i,j), (w,l)} = \begin{cases} \frac{1}{k-1} \frac{v_i}{v_i + v_l} & , \text{if } w=i, l \neq i \\ \frac{1}{k-1} \frac{v_w}{v_i + v_w} & , \text{if } l=i, w \neq i \\ 0 & \text{otherwise} \end{cases}$$

(same as a)

prob. go to w
 winner l loser
 given i won over
 j

states = $k \cdot (k-1)$

k people to choose from as winner
 one of remaining $k-1$ is the loser

$$\sum_{i=1}^k \sum_{j \neq i}^k v_i = \sum_{i=1}^k v_i (k-1) = (k-1) \sum_{i=1}^k v_i$$

Time-reversibility:

$$\pi_{i,j} P_{(i,j), (w,l)} = \pi_{w,l} P_{(w,l), (i,j)} \quad \forall (w,l), (i,j)$$

$w \neq i$:

$$\pi_{i,j} \frac{1}{k-1} \frac{v_w}{v_i + v_w} = \pi_{w,l} \frac{v_i}{v_i + v_w} \frac{1}{k-1}$$

$$\pi_{i,j} v_w = \pi_{w,l} v_i$$

Guess:

$$\pi_{i,j} = \frac{v_i}{\sum_{i=1}^k \sum_{j \neq i}^k v_i} = \frac{v_i}{(k-1) \sum_{i=1}^k v_i}$$

($k-1$ terms)

Can check, as in d), that this satisfies time-reversibility & is ≥ 0 & sums to 1.

$\pi_{i,j}$ is the proportion of time where i wins against j .

Hence, the proportion of time j loses is

$$\begin{aligned} \sum_{i \neq j} \pi_{i,j} &= \sum_{i \neq j} \frac{v_i}{(k-1) \sum_{i=1}^k v_i} \\ &= \frac{\sum_{i \neq j} v_i}{(k-1) \sum_{i=1}^k v_i} \end{aligned}$$

Hence, the proportion games j is in is the sum of the time she wins & the times she loses, so

proportion of games j is in

$$= \frac{v_j}{\sum_{i=1}^k v_i} + \frac{\sum_{i \neq j} v_i}{(k-1) \sum_{i=1}^k v_i}$$

From d) and above

$$= \frac{v_j (k-1) + \sum_{i \neq j} v_i}{(k-1) \sum_{i=1}^k v_i}$$

Exercise 2

a

$$P = \begin{bmatrix} 0 & p & 1-p \\ 1-p & 0 & p \\ p & 1-p & 0 \end{bmatrix}$$

b The chain is ergodic because all states communicate and p cannot assume values 0 and 1. By symmetry, $\pi_1 = 1/3, \pi_2 = 1/3, \pi_3 = 1/3$.

c

$$\begin{aligned} P[X_n = 1, X_{n+1} = 2] &= P[X_{n+1} = 2|X_n = 1]P[X_n = 1] \\ &= p_{12}P[X_n = 1]. \end{aligned}$$

For n large, $P[X_n = 1] = \pi_1 = 1/3$, therefore $P[X_n = 1, X_{n+1} = 2] = p/3$. Similarly, $P[X_n = 1, X_{n+1} = 2] = (1-p)/3$.

d To be time reversible, $\pi_i p_{ij} = \pi_j p_{ji}$ should hold. Since in this case $\pi_i = \pi_j \forall i, j$ then it must hold $p_{ij} = p_{ji}$, i.e., $p = 1-p$, so $p = 1/2$.

NB Note that p and $1-p$ may be inverted in the solution if the triangle was differently oriented.