

STK3405 - Exercise 6.3-6.6

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Exercise 6.3: Prove the upper bound of Corollary 6.2.8.

SOLUTION: The upper bound is

$$h(\mathbf{p}) \leq \prod_{j=1}^p \prod_{i \in P_j} \rho_i.$$

To prove this, note that for independent component state variables, the reliability of the minimal path series structures is

$$P(\rho_j(\mathbf{X}^{P_j}) = 1) = E[\prod_{i \in P_j} X_i] = \prod_{i \in P_j} \rho_i.$$

Hence,

$$h(\mathbf{p}) \leq \prod_{j=1}^p P(\rho_j(\mathbf{X}^{P_j}) = 1) = \prod_{j=1}^p \prod_{i \in P_j} \rho_i$$

where the inequality follows from Corollary 6.2.6 (or alternatively, by using the same proof technique as in this result).

Exercise 6.4: Prove the upper bound in Corollary 6.2.6 by applying the lower bound on the dual structure function ϕ^D .

SOLUTION: We know that

$$\max_{1 \leq j \leq p} \prod_{i \in P_j} p_i \leq h.$$

Applied to the dual structure function ϕ^D , we get:

$$\max_{1 \leq j \leq p^D} \prod_{i \in P_j^D} p_i^D \leq h^D.$$

From the definition of h^D, p_i^D and the fact that minimal path sets for ϕ are minimal cut sets for ϕ^D , this is equivalent with

$$\max_{1 \leq j \leq k} \prod_{i \in K_j} (1 - p_i) \leq 1 - h.$$

Hence, using that $\max\{\cdot\} = -\min\{-\cdot\}$ and the definition of the coproduct,

$$\begin{aligned} h &\leq 1 - \max_{1 \leq j \leq k} \prod_{i \in K_j} (1 - p_i) \\ &= \min_{1 \leq j \leq k} (1 - \prod_{i \in K_j} (1 - p_i)) \\ &= \min_{1 \leq j \leq k} \prod_{i \in K_j} p_i \end{aligned}$$

which is the upper bound of Corollary 6.2.6.

Exercise 6.5: Prove the upper bound in Corollary 6.2.8 by applying the lower bound on the dual structure function ϕ^D .

SOLUTION: We know that

$$\prod_{j=1}^k \prod_{i \in K_j} p_i \leq h(\mathbf{p}).$$

Applied to the dual structure,

$$\prod_{j=1}^{k^D} \prod_{i \in K_j^D} p_i^D \leq h^D(\mathbf{p}^D).$$

By using the definition of h^D , p_i^D and that minimal path sets of ϕ are minimal cut sets of ϕ^D , we find that this is equivalent to

$$\prod_{j=1}^p \prod_{i \in P_j} (1 - p_i) \leq 1 - h(\mathbf{p}).$$

That is,

$$\begin{aligned}h(\mathbf{p}) &\leq 1 - \prod_{j=1}^{\rho} \prod_{i \in P_j} (1 - p_i) \\&= 1 - \prod_{j=1}^{\rho} (1 - \prod_{i \in P_j} (1 - (1 - p_j))) \\&= 1 - \prod_{j=1}^{\rho} (1 - \prod_{i \in P_j} p_i) \\&= \prod_{j=1}^{\rho} \prod_{i \in P_j} p_i\end{aligned}$$

where we have used the definition of the coproduct and some algebra. Note that this is the upper bound we wanted to prove, so this concludes our solution.

Exercise 6.6: Consider the network in Figure 1 consisting of independent components with component reliabilities p .

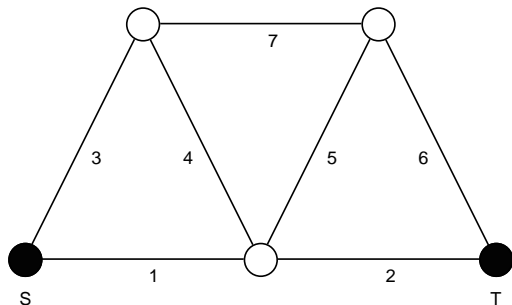


Figure: Illustration of the network for Exercise 6.6.

a) What is the reliability $h(p)$ for this system?

SOLUTION: We factor w.r.t. component 7:

$$h(\mathbf{p}) = ph(1_7, \mathbf{p}) + (1 - p)h(0_7, \mathbf{p}).$$

If component 7 works (make a drawing!): In this case, components 4 and 5 are in parallel. Hence, we can parallel reduce these two components to a new component 4' where $p_{4'} = p \amalg p$. The resulting structure is a bridge structure, so

$$\begin{aligned} h(1_7, \mathbf{p}) &= p_{4'}(p \amalg p)(p \amalg p) + (1 - p_{4'})(p^2 \amalg p^2) \\ &= \dots \\ &= (p \amalg p)^3 + (1 - p)^2(p^2 \amalg p^2). \end{aligned}$$

If component 7 doesn't work (make a drawing!): The resulting system is an s-p system. In this case, components 3 and 4 are in series, and so are 5 and 6. Hence, we can series reduce 3 and 4 to a new component 3' with reliability $p_{3'} = p^2$. Similarly, we can series reduce components 5 and 6 to a new component 5' with reliability $p_{5'} = p^2$. Then, we can parallel reduce 3' and 1 to a new component 1' with reliability $p_{1'} = p \parallel p^2$. Similarly, we can parallel reduce 5' and 2 to a new component 2' with reliability $p_{2'}$. Then,

$$\begin{aligned} h(0_7, \mathbf{p}) &= p_{1'} p_{2'} \\ &= (p \parallel p^2)^2. \end{aligned}$$

So,

$$h(\mathbf{p}) = p\{(p \prod p)^3 + (1 - p^2)(p^2 \prod p^2)\} + (1 - p)(p \prod p^2)^2.$$

b) For $p \in [0, 1]$, write a program to compute the reliability $h(p)$. Plot $h(p)$.

SOLUTION:

Please ask via e-mail if you have trouble with this. You should end up with an S-shaped curve starting at $(0,0)$ and ending in $(1,1)$.

c) In the same plot, illustrate the bounds from Corollary 6.2.6 and 6.2.8. Comment on the result.

SOLUTION:

Recall from Corollary 6.2.6 that

$$\max_{1 \leq j \leq p} \prod_{i \in P_j} p_i \leq h \leq \min_{1 \leq j \leq k} \prod_{i \in K_j} p_i.$$

The minimal path and cuts sets are:

Minimal path sets: $P_1 = \{1, 2\}$, $P_2 = \{2, 3, 4\}$, $P_3 = \{1, 5, 6\}$,
 $P_4 = \{3, 4, 5, 6\}$, $P_5 = \{3, 6, 7\}$, $P_6 = \{2, 3, 5, 7\}$, $P_7 = \{1, 4, 6, 7\}$.

Minimal cut sets: $K_1 = \{1, 3\}$, $K_2 = \{2, 6\}$, $K_3 = \{1, 4, 7\}$,
 $K_4 = \{2, 3, 4, 5\}$, $K_5 = \{2, 5, 7\}$, $K_6 = \{1, 4, 5, 6\}$.

So the lower bound is

$$\max_{1 \leq j \leq p} \prod_{i \in P_j} p = \max\{p^2, p^3, p^4\} = p^2$$

since $p \in [0, 1]$.

Similarly, the upper bound is

$$\min_{1 \leq j \leq k} \prod_{i \in K_j} p = \min\{1 - (1 - p)^2, 1 - (1 - p)^3, 1 - (1 - p)^4\} = 1 - (1 - p)^2$$

since $p \in [0, 1]$.

From Corollary 6.2.8,

$$\prod_{j=1}^k \prod_{i \in K_j} p_i \leq h(\mathbf{p}) \leq \prod_{j=1}^p \prod_{i \in P_j} p_i.$$

Hence, the lower bound is

$$\prod_{j=1}^k \prod_{i \in K_j} p = (1 - (1 - p)^2)^2 (1 - (1 - p)^3)^2 (1 - (1 - p)^4)^2.$$

And the upper bound is

$$\prod_{j=1}^p \prod_{i \in P_j} p = 1 - (1 - p^2)(1 - p^3)^3(1 - p^4)^3.$$

If you plot all of these bounds, you see that the bounds from Corollary 6.2.8 look overall better than those from Corollary 6.2.6. However, they are not always better: For small p 's, the lower bound from 6.2.6 is better than that from 6.2.8. For large p 's, the upper bound from 6.2.6 is better than that from 6.2.8.

d) Is it possible to improve these bounds further?

SOLUTION:

To improve the bounds further, one can take the maximum of the two lower bounds and similarly a minimum of the two upper bounds. This will still be lower and upper bounds, respectively, but these new bounds will be better than those from Corollary 6.2.6 and Corollary 6.2.8.