

STK3405 – Lecture 12

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Associated random variables

Definition (Associated random variables)

Let T_1, \dots, T_n be random variables, and let $\mathbf{T} = (T_1, \dots, T_n)$. We say that T_1, \dots, T_n are associated if

$$\text{Cov}(\Gamma(\mathbf{T}), \Delta(\mathbf{T})) \geq 0,$$

for all binary non-decreasing functions Γ and Δ .

NOTE: We only require $\text{Cov}(\Gamma(\mathbf{T}), \Delta(\mathbf{T})) \geq 0$ for all *binary* non-decreasing functions.



Associated random variables (cont.)

Theorem (Properties of Associated variables)

Associated random variables have the following properties:

- (i) Any subset of a set of associated random variables also consists of associated random variables.*
- (ii) A single random variable is always associated.*
- (iii) Non-decreasing functions of associated random variables are associated.*
- (iv) If two sets of associated random variables are independent, then their union is a set of associated random variables.*



Associated random variables (cont.)

Theorem (Association of paths and cuts)

Let X_1, \dots, X_n be the associated or independent component state variables of a monotone system (C, ϕ) . Moreover, let the minimal path series structures of the system be $(P_1, \rho_1), \dots, (P_p, \rho_p)$, where:

$$\rho_j(\mathbf{X}^{P_j}) = \prod_{i \in P_j} X_i, \quad j = 1, \dots, p.$$

Then, ρ_1, \dots, ρ_p are associated.

Similarly, let the minimal cut parallel structures of the system be $(K_1, \kappa_1), \dots, (K_k, \kappa_k)$, where:

$$\kappa_j(\mathbf{X}^{K_j}) = \prod_{i \in K_j} X_i, \quad j = 1, \dots, k.$$

Then, $\kappa_1, \dots, \kappa_k$ are associated.

Bounds for the system reliability

Theorem (6.2.1)

Let T_1, \dots, T_n be associated random variables such that $0 \leq T_i \leq 1$, $i = 1, \dots, n$. We then have:

$$E\left[\prod_{i=1}^n T_i\right] \geq \prod_{i=1}^n E[T_i]$$

$$E\left[\prod_{i=1}^n T_i\right] \leq \prod_{i=1}^n E[T_i]$$



Bounds for the system reliability (cont.)

In the next results we interpret $T_i, i = 1, \dots, n$ as the lifetimes of components in a binary monotone system.

Theorem (6.2.2)

Let T_1, \dots, T_n be non-negative associated random variables. Then, for all $t_i, i = 1, \dots, n$:

$$P\left[\bigcap_{i=1}^n (T_i > t_i)\right] \geq \prod_{i=1}^n P[T_i > t_i],$$

$$P\left[\bigcap_{i=1}^n (T_i \leq t_i)\right] \geq \prod_{i=1}^n P[T_i \leq t_i].$$



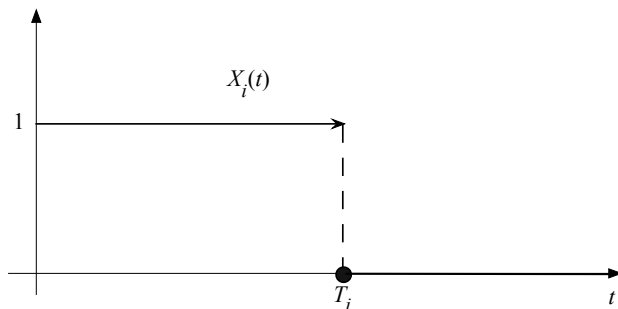
Bounds for the system reliability (cont.)

PROOF: Interpreting T_i as the lifetime of the i th component, we may introduce the binary component state process $\{X_i(t)\}$, $i = 1, \dots, n$, where:

$$X_i(t) = \begin{cases} 1, & t < T_i \\ 0, & t \geq T_i. \end{cases}$$



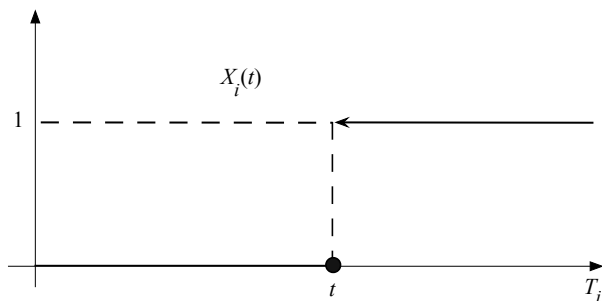
Bounds for the system reliability (cont.)



In this we have plotted $X_i(t)$ as a function of $t \geq 0$ for a given value of T_i . The plot shows that $X_i(t)$ is a non-increasing function of t .



Bounds for the system reliability (cont.)



Alternatively, $X_i(t)$ can be plotted as a function of the lifetime T_i for a *fixed point of time* t . We see that $X_i(t)$ is a non-decreasing function of T_i .

Hence, since T_1, \dots, T_n are assumed to be associated random variables, it follows that $X_1(t), \dots, X_n(t)$ are associated for all $t \geq 0$.



Bounds for the system reliability (cont.)

Applying the first part of Theorem 6.2.1, we get:

$$\begin{aligned}P\left[\bigcap_{i=1}^n (T_i > t_i)\right] &= P\left[\bigcap_{i=1}^n (X_i(t_i) = 1)\right] \\&= P\left[\prod_{i=1}^n X_i(t_i) = 1\right] = E\left[\prod_{i=1}^n X_i(t_i)\right] \\&\geq \prod_{i=1}^n E[X_i(t_i)] = \prod_{i=1}^n P[X_i(t_i) = 1] \\&= \prod_{i=1}^n P[T_i > t_i],\end{aligned}$$

and thus the first inequality is proved.



Bounds for the system reliability (cont.)

Applying the second part of Theorem 6.2.1, we get:

$$\begin{aligned}P\left[\bigcap_{i=1}^n (T_i \leq t_i)\right] &= P\left[\bigcap_{i=1}^n (X_i(t_i) = 0)\right] = 1 - P\left[\bigcup_{i=1}^n (X_i(t_i) = 1)\right] \\&= 1 - P\left[\prod_{i=1}^n X_i(t_i) = 1\right] = 1 - E\left[\prod_{i=1}^n X_i(t_i)\right] \\&\geq 1 - \prod_{i=1}^n E[X_i(t_i)] = 1 - \left[1 - \prod_{i=1}^n (1 - E[X_i(t_i)])\right] \\&= \prod_{i=1}^n (1 - P[X_i(t_i) = 1]) = \prod_{i=1}^n P[T_i \leq t_i],\end{aligned}$$

and thus the second inequality is proved.



Bounds for the system reliability (cont.)

Corollary (6.2.3)

Let T_1, \dots, T_n be non-negative associated random variables. Then:

$$P(\min_{1 \leq i \leq n} T_i > t) = P(\bigcap_{i=1}^n T_i > t) \geq \prod_{i=1}^n P(T_i > t)$$
$$P(\max_{1 \leq i \leq n} T_i > t) = 1 - P(\bigcap_{i=1}^n T_i \leq t) \leq \prod_{i=1}^n P(T_i > t).$$

NOTE: If (C, ϕ) is a series system with component lifetimes T_1, \dots, T_n , then the system lifetime is equal to $\min_{1 \leq i \leq n} T_i$.

If (C, ϕ) is a parallel system with component lifetimes T_1, \dots, T_n , then the system lifetime is equal to $\max_{1 \leq i \leq n} T_i$.



Bounds for the system reliability (cont.)

We skip the following results:

- Theorem 6.2.4 (Very crude upper and lower bounds)
- Theorem 6.2.5 (We incorporate this result in Corollary 6.2.6)
- Theorem 6.2.7 (We incorporate this result in Corollary 6.2.8)



Bounds for the system reliability (cont.)

Corollary (6.2.6)

Consider a monotone system (C, ϕ) , where $C = \{1, \dots, n\}$ and with minimal path sets P_1, \dots, P_p and minimal cut sets K_1, \dots, K_k .

Moreover, assume that the component state variables, X_1, \dots, X_n are associated with component reliabilities p_1, \dots, p_n . Then we have:

$$\max_{1 \leq j \leq p} \prod_{i \in P_j} p_i \leq h \leq \min_{1 \leq j \leq k} \prod_{i \in K_j} p_i.$$



Bounds for the system reliability (cont.)

PROOF: We have that:

$$\min_{i \in P_r} X_i \leq \max_{1 \leq r \leq p} \min_{i \in P_r} X_i = \phi(\mathbf{X}) = \min_{1 \leq s \leq k} \max_{i \in K_s} X_i \leq \max_{i \in K_s} X_i,$$

for all $r = 1, \dots, p$ and all $s = 1, \dots, k$.

This implies that:

$$P(\min_{i \in P_r} X_i = 1) \leq h \leq P(\max_{i \in K_s} X_i = 1)$$

for all $r = 1, \dots, p$ and all $s = 1, \dots, k$.

Hence, we must have:

$$\max_{1 \leq j \leq p} P[\min_{i \in P_j} X_i = 1] \leq h \leq \min_{1 \leq j \leq k} P[\max_{i \in K_j} X_i = 1].$$



Bounds for the system reliability (cont.)

Furthermore, since X_1, \dots, X_n are associated, we may use Theorem 6.2.1 and get:

$$P[\min_{i \in P_j} X_i = 1] = E[\prod_{i \in P_j} X_i] \geq \prod_{i \in P_j} E[X_i] = \prod_{i \in P_j} p_i$$

$$P[\max_{i \in K_j} X_i = 1] = E[\prod_{i \in K_j} X_i] \leq \prod_{i \in K_j} E[X_i] = \prod_{i \in K_j} p_i$$

Inserting these inequalities into the bounds on the previous slide we get:

$$\max_{1 \leq j \leq p} \prod_{i \in P_j} p_i \leq h \leq \min_{1 \leq j \leq k} \prod_{i \in K_j} p_i.$$



Bounds for the system reliability (cont.)

Corollary (6.2.8)

Let (C, ϕ) be a binary monotone system where $C = \{1, \dots, n\}$, and assume that the component state variables, X_1, \dots, X_n are independent with component reliabilities p_1, \dots, p_n .

Moreover, let P_1, \dots, P_p and K_1, \dots, K_k be respectively the minimal path and cut sets of the system.

Then we have:

$$\prod_{j=1}^k \prod_{i \in K_j} p_i \leq h(\mathbf{p}) \leq \prod_{j=1}^p \prod_{i \in P_j} p_i.$$



Bounds for the system reliability (cont.)

PROOF: We introduce:

$$\rho_j(\mathbf{X}) = \prod_{i \in P_j} X_i, \quad j = 1, \dots, p,$$

$$\kappa_j(\mathbf{X}) = \prod_{i \in K_j} X_i, \quad j = 1, \dots, k.$$

Since ρ_1, \dots, ρ_p and $\kappa_1, \dots, \kappa_k$ are non-decreasing functions of \mathbf{X} , they are associated. Hence, by Theorem 6.2.1 we have:

$$h(\mathbf{p}) = E\left[\prod_{j=1}^p \prod_{i \in P_j} X_i\right] = E\left[\prod_{j=1}^p \rho_j(\mathbf{X})\right] \leq \prod_{j=1}^p E[\rho_j(\mathbf{X})]$$

$$h(\mathbf{p}) = E\left[\prod_{j=1}^k \prod_{i \in K_j} X_i\right] = E\left[\prod_{j=1}^k \kappa_j(\mathbf{X})\right] \geq \prod_{j=1}^k E[\kappa_j(\mathbf{X})]$$



Bounds for the system reliability (cont.)

Moreover, since the component state variables, X_1, \dots, X_n , are *independent*, we have:

$$E[\rho_j(\mathbf{X})] = E\left[\prod_{i \in P_j} X_i\right] = \prod_{i \in P_j} p_i,$$

$$E[\kappa_j(\mathbf{X})] = E\left[\prod_{i \in K_j} X_i\right] = \prod_{i \in K_j} p_i.$$

Inserting this into the bounds on the previous slide, i.e.:

$$\prod_{j=1}^k E[\kappa_j(\mathbf{X})] \leq h(\mathbf{p}) \leq \prod_{j=1}^p E[\rho_j(\mathbf{X})],$$

we get:

$$\prod_{j=1}^k \prod_{i \in K_j} p_i \leq h(\mathbf{p}) \leq \prod_{j=1}^p \prod_{i \in P_j} p_i.$$



Bounds for the system reliability (cont.)

In the coming examples we shall compare the bounds from Corollary 6.2.6 to those from Corollary 6.2.8.

Let $h_n(p)$ denote the reliability of a parallel system of n components where all components have the same reliability p . We then have:

$$h_2(p) = p \text{ II } p = 1 - (1 - p)(1 - p)$$

$$= 1 - (1 - 2p + p^2) = 2p - p^2,$$

$$h_3(p) = p \text{ II } p \text{ II } p = 1 - (1 - p)(1 - p)(1 - p)$$

$$= 1 - (1 - 3p + 3p^2 - p^3) = 3p - 3p^2 + p^3.$$



Bounds for the system reliability (cont.)

EXAMPLE 1: A 3-out-of-4 system with $p_i = p$, $i = 1, 2, 3, 4$ where all the component state variables are independent.

The minimal path sets for the 3-out-of-4 system are:

$$P_1 = \{1, 2, 3\}, P_2 = \{1, 2, 4\}, P_3 = \{1, 3, 4\}, P_4 = \{2, 3, 4\},$$

and the minimal cut sets are:

$$K_1 = \{1, 2\}, K_2 = \{1, 3\}, K_3 = \{1, 4\}, K_4 = \{2, 3\}, K_5 = \{2, 4\}, K_6 = \{3, 4\}.$$



Bounds for the system reliability (cont.)

The lower and upper bounds in Corollary 6.2.6 are denoted by $l_1(p)$ and $u_1(p)$ respectively, and are given by:

$$l_1(p) = \max_{1 \leq j \leq 4} \prod_{i \in P_j} p = \max_{1 \leq j \leq 4} p^3 = p^3,$$

$$u_1(p) = \min_{1 \leq j \leq 6} \prod_{i \in K_j} p = \min_{1 \leq j \leq 6} h_2(p) = 2p - p^2.$$

The lower and upper bounds in Corollary 6.2.8 are denoted by $l_2(p)$ and $u_2(p)$ respectively, and are given by:

$$l_2(p) = \prod_{j=1}^6 \prod_{i \in K_j} p = \prod_{j=1}^6 h_2(p) = (2p - p^2)^6,$$

$$u_2(p) = \prod_{j=1}^4 \prod_{i \in P_j} p = h_2(p^3) \Pi h_2(p^3) = 2(2p^3 - p^6) - (2p^3 - p^6)^2.$$



Bounds for the system reliability (cont.)

The *true* reliability of the 3-out-of-4 system is given by:

$$h(p) = \sum_{i=3}^4 \binom{4}{i} p^i (1-p)^{4-i} = 4p^3(1-p) + p^4.$$



Bounds for the system reliability (cont.)

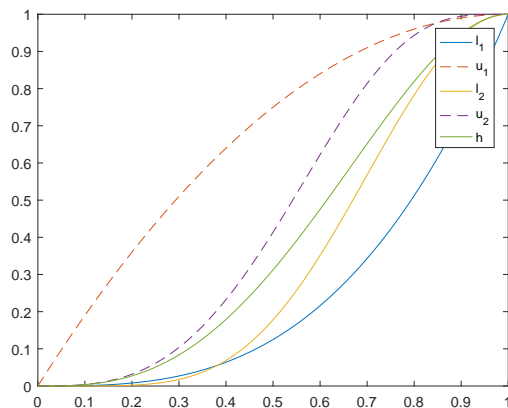


Figure: The true reliability function h as well as the bounds l_1 , u_1 , l_2 , u_2 .



Bounds for the system reliability (cont.)

EXAMPLE 2: A bridge system with $p_i = p$, $i = 1, 2, 3, 4, 5$ where all the component state variables are independent.

The minimal path sets for the bridge system are:

$$P_1 = \{1, 4\}, P_2 = \{1, 3, 5\}, P_3 = \{2, 3, 4\}, P_4 = \{2, 5\},$$

and the minimal cut sets are:

$$K_1 = \{1, 2\}, K_2 = \{1, 3, 5\}, K_3 = \{2, 3, 4\}, K_4 = \{4, 5\}.$$



Bounds for the system reliability (cont.)

The lower and upper bounds in Corollary 6.2.6 are denoted by $l_1(p)$ and $u_1(p)$ respectively, and are given by:

$$l_1(p) = \max_{1 \leq j \leq 4} \prod_{i \in P_j} p = \max(p^2, p^3, p^3, p^2) = p^2,$$

$$u_1(p) = \min_{1 \leq j \leq 4} \prod_{i \in K_j} p = \min(h_2(p), h_3(p), h_3(p), h_2(p)) = 2p - p^2.$$

The lower and upper bounds in Corollary 6.2.8 are denoted by $l_2(p)$ and $u_2(p)$ respectively, and are given by:

$$l_2(p) = \prod_{j=1}^4 \prod_{i \in K_j} p = h_2(p)^2 \cdot h_3(p)^2,$$

$$u_2(p) = \prod_{j=1}^4 \prod_{i \in P_j} p = h_2(p^2) \cdot h_2(p^3).$$



Bounds for the system reliability (cont.)

The *true* reliability of the bridge system is given by:

$$h(p) = p \cdot h_2(p)^2 + (1 - p) \cdot h_2(p^2).$$



Bounds for the system reliability (cont.)

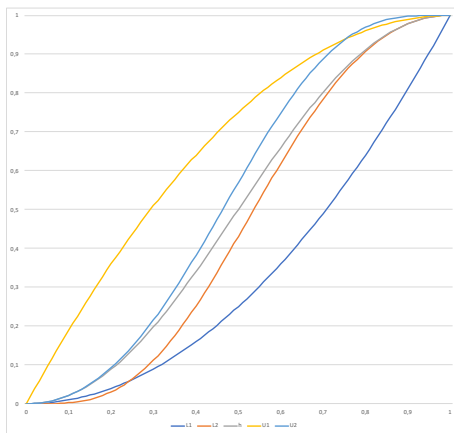


Figure: The true reliability function h as well as the bounds l_1 , u_1 , l_2 , u_2 .



Bounds for the system reliability (cont.)

We see that in both examples the bounds from Corollary 6.2.8 are better than those from Corollary 6.2.6 for *most* of the p -values.

NOTE:

- The lower bound l_1 from Corollary 6.2.6 is better than l_2 from Corollary 6.2.8 for small values of p .
- The upper bound u_1 from Corollary 6.2.6 is better than u_2 from Corollary 6.2.8 for large p -values.

In order to always get the best bounds, we may introduce l^* and u^* defined as follows:

$$l^* = \max(l_1, l_2),$$

$$u^* = \min(u_1, u_2)$$

